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Atomic Compare

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History

R1

R1 was reviewed by SG1 in Brno with the following feedback:

Poll: SG1 likes the P4221R1 design direction and would like to see this again with following changes:

- `compare` either removed, or include an extended discussion on the value of `compare` and its potential pitfalls on provenance due to expected not being updated; potentially update `compare` to take expected by reference and update its value instead.
- discussion on `memcmp` to address whether this is adding a new capability to the standard that was not possible before, or just adding new APIs to make something that's already possible with `memcmp` more ergonomic.
- discussion on padding for the new APIs and the impacted examples
- add C free functions and `atomic_ref` support

R2

R2 addresses SG1 feedback by:

- Removing `compare`.
- Clarifying motivation.
- Discussing equality comparison in the standard, including the effect of padding and `memcmp` capabilities.
- Stating the new capabilities provided by the proposed function.
- Completing the wording, including free functions, `atomic_ref`, and specializations.

Introduction

This paper proposes adding a `compare_load` member function to `std::atomic<T>` and `std::atomic_ref<T>`, alongside corresponding free functions. This function performs an atomic comparison of the atomic object's value with an expected value, following the same value representation comparison semantics as `compare_exchange_strong`. It updates the expected argument on failure, operating entirely as a pure read without writing a new value to the atomic object.

Motivation

Standard C++ currently lacks a mechanism to perform a consistent, read-only, padding-independent value representation equality check on an atomic object.

A secondary motivation is to provide programmers with distinct memory orders for comparison success and failure. Currently, programmers must either commit upfront to a single memory order regardless of the evaluation outcome, or navigate the complex semantics of standalone fences.

The proposed `compare_load` function natively resolves both issues, providing a read-only, padding-independent value representation equality check with built-in support for outcome-dependent memory ordering.

Equality Comparison in the Standard

To clarify the exact semantics evaluated by `compare_load`, it is necessary to distinguish the three forms of equality comparison in C++ (see [\[basic.types.general\] p4](#)):

- **Semantic Equality (`operator==`):** Evaluates logical equivalence. It inherently ignores padding bits and can be arbitrarily customized via user-defined overloads.
- **Object Representation (`memcmp`):** Evaluates bitwise equality of the entire memory footprint of an object, including padding bits.
- **Value Representation:** Evaluates bitwise equality of only the bits that participate in representing the value, excluding any padding bits.

These distinct definitions result in behavioral differences when writing concurrent code:

- `operator==` can yield different results from both `memcmp` and value representation, even when a type has zero padding.
- `memcmp` is functionally identical to a value representation comparison only for types *without* padding. For types *with* padding, `memcmp` diverges and becomes an unsafe proxy for value equality due to the potential for indeterminate padding bits.
- Currently, `compare_exchange` is the only mechanism in the standard library capable of consistently evaluating value representation equality in the presence of padding, but it couples this evaluation to a mutating memory write.

New Capabilities

The introduction of **compare_load** provides fundamental concurrency capabilities that cannot be achieved by combining existing standard library facilities:

- **Read-only padding-neutral value representation comparison:** Unlike **operator==** (which relies on semantic logic), **memcmp** (which unsafely evaluates indeterminate padding bits), and **compare_exchange** (which incurs a hardware write), **compare_load** provides a consistent, read-only mechanism to evaluate value representation equality.
- **Outcome-dependent memory ordering:** Unlike a standard atomic **load** followed by a nonatomic **memcmp** comparison, which requires a programmer to commit to a single memory order upfront or navigate complicated standalone fences, **compare_load** allows programmers to express distinct memory orders for success and failure, optimizing performance based on the evaluation outcome.

Relationship to **compare_exchange**

The existing **compare_exchange** operations, **compare_exchange_weak** and **compare_exchange_strong**, defined in [\[atomics.types.operations\] p21–28](#), provide the semantic basis for the proposed **compare_load** function. Specifically, they establish:

- **Value representation** comparison (p23, p28).
- Handling of **padding bits** (p25, p28).
- Update of the **expected** argument on comparison failure (p23).
- Categorizing the operation as an atomic load upon comparison failure (p23).

Note that the rationale for a weak variant (p27) is inapplicable to **compare_load**, as it operates strictly as a read-only evaluation without a corresponding conditional store.

Furthermore, **compare_load** strictly follows the pointer provenance semantics of **compare_exchange**. The design intent is that any future standard modifications regarding pointer provenance or "zap" fixes applied to the specification of **compare_exchange** will apply equally and systematically to **compare_load**.

Preview

We propose adding a **compare_load** member function to `std::atomic<T>` and `std::atomic_ref<T>`, along with corresponding non-member functions.

compare_load

```
bool compare_load(T& expected, memory_order success, memory_order failure)
    const volatile noexcept;
constexpr bool compare_load(T& expected, memory_order success, memory_order failure)
    const noexcept;
bool compare_load(T& expected, memory_order order = memory_order::seq_cst)
    const volatile noexcept;
constexpr bool compare_load(T& expected, memory_order order = memory_order::seq_cst)
    const noexcept;
```

Note that **compare_load**:

- Performs value representation equality check between the value of the atomic object and the value of **expected**.
- Supports the same strong comparison semantics as **compare_exchange_strong**.
- Supports the same memory order constraints as a **load** operation but with optional distinct orders for comparison success and failure.
- Updates the **expected** argument with the value read from the atomic object if the equality comparison fails, as with the failure case of **compare_exchange**.

Usage Examples

The following table shows code snippets using the current standard C++ (on the left) and using the proposed `atomic_compare_load` function (on the right).

Current Standard C++	With Proposed Function
<pre> /// EX1 Load-compute-validate loop T value = a.load(); while (true) { // Calculation based on read value R result = calc(value); // Manually reload current value T current = a.load(); // Check validity of used value if (current == value) return result; // Update value manually value = current; } </pre>	<pre> /// EX1 Load-compute-validate loop T value = a.load(); while (true) { // Calculation based on read value R result = calc(value); // Check validity of used value if (a.compare_load(value)) return result; // Failed compare updates value // automatically. } </pre>
<pre> /// EX2 HP TRY_PROTECT bool try_protect(T*& ptr, atomic<T*>& src) { T* p = ptr; hp_.reset_protection(ptr); /* Light asymmetric memory barrier */ ptr = src.load(acquire); if (p == ptr) return true; else { hp_.reset_protection(); return false; } } </pre>	<pre> /// EX2 HP TRY_PROTECT bool try_protect(T*& ptr, atomic<T*>& src) { hp_.reset_protection(ptr); /* Light asymmetric memory barrier */ if (src.compare_load(ptr, acquire)) return true; else { hp_.reset_protection(); return false; } } </pre>
<pre> /// EX3 Outcome-dependent memory ordering T observed = a.load(relaxed); if (observed == value) { std::atomic_thread_fence(seq_cst); /* success path */ } else { std::atomic_thread_fence(acquire); value = observed; /* failure path */ } </pre>	<pre> /// EX3 Outcome-dependent memory ordering if (a.compare_load(value, seq_cst, acquire)) { // No need for explicit fence /* success path */ } else { // No need for explicit fence /* failure path */ } </pre>

Proposed Wording

In 17.3.2 [\[version.syn\]](#) p2, add the following macro to the list in alphabetical order:

```
__cpp_lib_atomic_compare_load 202XXXXXL // freestanding, also in <atomic>
```

In 32.5.2 [\[atomics.syn\]](#), add the following to namespace `std` after the `atomic_compare_exchange` declarations:

```
template<class T>
    bool atomic_compare_load(
        volatile atomic<T>*, typename atomic<T>::value_type*) noexcept;
template<class T>
    constexpr bool atomic_compare_load(
        atomic<T>*, typename atomic<T>::value_type*) noexcept;
template<class T>
    bool atomic_compare_load_explicit(volatile atomic<T>*,
        typename atomic<T>::value_type*, memory_order, memory_order) noexcept;
template<class T>
    constexpr bool atomic_compare_load_explicit(atomic<T>*,
        typename atomic<T>::value_type*, memory_order, memory_order) noexcept;
```

In 32.5.7.1 [\[atomics.ref.generic.general\]](#), add the following to the `atomic_ref<T>` class template synopsis after `compare_exchange`:

```
constexpr bool compare_load(value_type&, memory_order, memory_order) const noexcept;
constexpr bool compare_load(value_type&, memory_order = memory_order::seq_cst)
    const noexcept;
```

In 32.5.7.2 [\[atomics.ref.ops\]](#), add the following after `compare_exchange`:

```
constexpr bool compare_load(T& expected, memory_order success, memory_order failure)
    const noexcept;
constexpr bool compare_load(T& expected, memory_order order = memory_order::seq_cst)
    const noexcept;
```

Preconditions: `success` and `failure` are each one of `memory_order::relaxed`, `memory_order::acquire`, or `memory_order::seq_cst`.

Effects: Retrieves the value in `expected`. It then atomically compares the value representation of the value pointed to by **ptr* for equality with that previously retrieved from `expected`. If and only if the comparison is true, memory is affected according to the value of `success`, and if the comparison is false, memory is affected

according to the value of **failure**. When only one **memory_order** argument is supplied, the value of **success** is **order**, and the value of **failure** is **order**. If and only if the comparison is false then, after the atomic operation, the value in **expected** is replaced by the value pointed to by **ptr* during the atomic comparison. These operations are atomic load operations on the memory pointed to by **ptr*.

Returns: The result of the comparison.

In 32.5.7.3 [\[atomics.ref.int\]](#), add the following after **compare_exchange**:

```
constexpr bool compare_load(value_type&, memory_order, memory_order) const noexcept;
constexpr bool compare_load(value_type&, memory_order = memory_order::seq_cst)
    const noexcept;
```

In 32.5.7.4 [\[atomics.ref.float\]](#), add the following after **compare_exchange**:

```
constexpr bool compare_load(value_type&, memory_order, memory_order) const noexcept;
constexpr bool compare_load(value_type&, memory_order = memory_order::seq_cst)
    const noexcept;
```

In 32.5.7.5 [\[atomics.ref.pointer\]](#), add the following after **compare_exchange**:

```
constexpr bool compare_load(value_type&, memory_order, memory_order) const noexcept;
constexpr bool compare_load(value_type&, memory_order = memory_order::seq_cst)
    const noexcept;
```

In 32.5.8.1 [\[atomics.types.generic.general\]](#), add the following to the **atomic<T>** class template synopsis after **compare_exchange**:

```
bool compare_load(T&, memory_order, memory_order) const volatile noexcept;
constexpr bool compare_load(T&, memory_order, memory_order) const noexcept;
bool compare_load(T&, memory_order = memory_order::seq_cst) const volatile noexcept;
constexpr bool compare_load(T&, memory_order = memory_order::seq_cst) const noexcept;
```

In 32.5.8.2 [\[atomics.types.operations\]](#), add the following after **compare_exchange**:

```
bool compare_load(T& expected, memory_order success, memory_order failure)
    const volatile noexcept;
constexpr bool compare_load(T& expected, memory_order success, memory_order failure)
    const noexcept;
bool compare_load(T& expected, memory_order order = memory_order::seq_cst)
    const volatile noexcept;
constexpr bool compare_load(T& expected, memory_order order = memory_order::seq_cst)
```



```
const noexcept;
```

Constraints: For the volatile overload of this function, `is_always_lock_free` is true.

Preconditions: `success` and `failure` are each one of `memory_order::relaxed`, `memory_order::acquire`, or `memory_order::seq_cst`.

Effects: Retrieves the value in `expected`. It then atomically compares the value representation of the value pointed to by `this` for equality with that previously retrieved from `expected`. If and only if the comparison is true, memory is affected according to the value of `success`, and if the comparison is false, memory is affected according to the value of `failure`. When only one `memory_order` argument is supplied, the value of `success` is `order`, and the value of `failure` is `order`. If and only if the comparison is false then, after the atomic operation, the value in `expected` is replaced by the value pointed to by `this` during the atomic comparison. These operations are atomic load operations on the memory pointed to by `this`.

Returns: The result of the comparison.

In 32.5.8.2 [atomics.types.operations] p28, modify Note 6, Note 7, and Note 8 as follows:

[Note 6: Under cases where the memcpy and memcmp semantics of the compare-and-exchange [and compare_load](#) operations apply, the comparisons can fail for values that compare equal with `operator==` if the value representation has trap bits or alternate representations of the same value. ... -- end note]

[Note 7: Because compare-and-exchange [and compare_load](#) acts on an object's value representation, padding bits that never participate in the object's value representation are ignored. ... -- end note]

[Note 8: For a union with bits that participate in the value representation of some members but not others, compare-and-exchange [and compare_load](#) might always fail. This is because such padding bits have an indeterminate value when they do not participate in the value representation of the active member. ... -- end note]

In 32.5.8.3 [atomics.types.int], add the following after `compare_exchange`:

```
bool compare_load(integral-type&, memory_order, memory_order) const volatile noexcept;
constexpr bool compare_load(integral-type&, memory_order, memory_order) const noexcept;
bool compare_load(integral-type&, memory_order = memory_order::seq_cst) const volatile
    noexcept;
constexpr bool compare_load(integral-type&, memory_order = memory_order::seq_cst) const
    noexcept;
```

In 32.5.8.4 [atomics.types.float], add the following after `compare_exchange`:

```
bool compare_load(floating-point-type&, memory_order, memory_order) const volatile
```

```
    noexcept;
constexpr bool compare_load(floating-point-type&, memory_order, memory_order) const
    noexcept;
bool compare_load(floating-point-type&, memory_order = memory_order::seq_cst) const
    volatile noexcept;
constexpr bool compare_load(floating-point-type&, memory_order = memory_order::seq_cst)
    const noexcept;
```

In 32.5.8.5 [\[atomics.types.pointer\]](#), add the following after `compare_exchange`:

```
bool compare_load(T*&, memory_order, memory_order) const volatile noexcept;
constexpr bool compare_load(T*&, memory_order, memory_order) const noexcept;
bool compare_load(T*&, memory_order = memory_order::seq_cst) const volatile noexcept;
constexpr bool compare_load(T*&, memory_order = memory_order::seq_cst) const noexcept;
```