Abstract

Fold expressions, which syntactically look deceptively like conjunctions/subjections for the purpose of constraint ordering are in fact atomic constraints. We propose rules for the normalization and ordering of fold expressions over && and ||.

Revisions

R1

- Wording improvements: The previous version of this paper incorrectly looked at the size of the packs involved in the fold expressions. This was unnecessary and was removed. The current design does not look at the template argument/parameter mapping to establish subsumption of fold expressions.

- A complete implementation of this proposal is available on Compiler Explorer. The implementation section was expanded.

- Add an additional example.

Motivation

This paper is an offshoot of P2841R0 [2] which described the issue with lack of subsumption for fold expressions. This was first observed in a Concept TS issue.

This question comes up ever so often on online boards and various chats.

- [StackOverflow] How are fold expressions used in the partial ordering of constraints?
- [StackOverflow] How to implement the generalized form of std::same_as?

In Urbana, core observed “We can’t constrain variadic templates without fold-expressions” and almost folded (!) fold expressions into the concept TS. The expectation that these features should interoperate well then appear long-standing.
Subsumption and fold expressions over && and ||

Consider:

```cpp
template <class T> concept A = std::is_move_constructible_v<T>;
template <class T> concept B = std::is_copy_constructible_v<T>;
template <class T> concept C = A<T> && B<T>;
```

```cpp
template <class... T>
requires (A<T> && ...)
void g(T...);
```

```cpp
template <class... T>
requires (C<T> && ...)
void g(T...);
```

We want to apply the subsumption rule to the normalized form of the requires clause (and its arguments). As of C++23, the above `g` is ambiguous.

This is useful when dealing with algebraic-type classes. Consider a concept constraining a (simplified) environment implementation via a type-indexed `std::tuple`. (In real code, the environment is a type-tag indexed map.)

```cpp
template <typename X, typename... T>
concept environment_of = (... && requires (X& x) { { get<T>(x) } -> std::same_as<T&>; } );
```

```cpp
auto f(sender auto&& s, environment_of<std::stop_token> auto env); // uses std::allocator
auto f(sender auto&& s, environment_of<std::stop_token, std::pmr::allocator> auto env); // uses given allocator
```

Without the subsumption fixes to fold expressions, the above two overloads conflict, even though they should be partially ordered.

A similar example courtesy of Barry Revzin:

```cpp
template <std::ranges::bidirectional_range R> void f(R&); // #1
template <std::ranges::random_access_range R> void f(R&); // #2
```

```cpp
template <std::ranges::bidirectional_range... R> void g(R&...); // #3
template <std::ranges::random_access_range... R> void g(R&...); // #4
```

<table>
<thead>
<tr>
<th>C++23</th>
<th>This Paper</th>
</tr>
</thead>
<tbody>
<tr>
<td>f(std::vector{1, 2, 3}); // Ok</td>
<td>f(std::vector{1, 2, 3}); // Ok, calls #2</td>
</tr>
<tr>
<td>g(std::vector{1, 2, 3}); // Error: call to 'g' is ambiguous</td>
<td>g(std::vector{1, 2, 3}); // Ok, calls #4</td>
</tr>
</tbody>
</table>

[Compiler Explorer Demo]
Impact on the standard

This change makes ambiguous overload valid and should not break existing valid code.

Implementability

This was implemented in Clang. Importantly, what we propose does not affect compilers’ ability to partially order functions by constraints without instantiating them, nor does it affect the caching of subsumption, which is important to minimize the cost of concepts on compile time: The template arguments of the constraint expressions do not need to be observed to establish subsumption.

An implementation does need to track whether an atomic constraint that contains an un-expanded pack was originally part of a and/or fold expression to properly implement the subsumption rules (&& subsumes || && & & || subsumes | |).

What this paper is not

When the pattern of the fold-expressions is a ‘concept’ template parameter, this paper does not apply. In that case, we need different rules which are covered in P2841R0 [2] along with the rest of the “concept template parameter” feature (specifically, for concept patterns we need to decompose each concept into its constituent atomic constraints and produce a fully decomposed sequence of conjunction/disjunction).

Design and wording strategy

To simplify the wording, we first normalize fold expressions to extract the non-pack expression of binary folds into its own normalized form, and transform (... && A) into (A && ...) as they are semantically identical for the purpose of subsumption. We then are left with either (A && ...) or (A || ...), and for packs of the same size, the rules of subsumptions are the same as for that of atomic constraints.

Wording

Constraint normalization [temp.constr.normal]

The normal form of an expression E is a constraint [temp.constr.constr] that is defined as follows:

- The normal form of an expression ( E ) is the normal form of E.
- The normal form of an expression E1 || E2 is the disjunction [temp.constr.op] of the normal forms of E1 and E2.
- The normal form of an expression E1 & & E2 is the conjunction of the normal forms of E1 and E2.
The normal form of a concept-id $C<A_1, A_2, \ldots, A_n>$ is the normal form of the constraint-expression of $C$, after substituting $A_1, A_2, \ldots, A_n$ for $C$’s respective template parameters in the parameter mappings in each atomic constraint. If any such substitution results in an invalid type or expression, the program is ill-formed; no diagnostic is required.

**[Example]**

```cpp
template<typename T> concept A = T::value || true;
template<typename U> concept B = A<U*>;
template<typename V> concept C = B<V&>;
```

Normalization of $B$’s constraint-expression is valid and results in $T::value$ (with the mapping $T \mapsto U*$) $\lor$ $true$ (with an empty mapping), despite the expression $T::value$ being ill-formed for a pointer type $T$. Normalization of $C$’s constraint-expression results in the program being ill-formed, because it would form the invalid type $V&*$ in the parameter mapping. — end example}

For a fold-operator [expr.prim.fold] that is either $\&\&$ or $||$, the normal form of an expression $(\ldots$ fold-operator $E\ldots)$ is $(E$ fold-operator $\ldots)$.

The normal form of an expression $(E_1 \&\& \ldots \&\& E_2)$ or $(E_2 \&\& \ldots \&\& E_1)$ where $E_1$ contains an unexpanded pack is the conjunction of the normal forms of $E_2$ and $(E_1 \&\& \ldots)$.

**[Editor’s note: With P2841 [?], we want to treat fold expression differently when the pattern of the fold expression denotes a concept template parameter.]**

The normal form of an expression $(E_1 || \ldots || E_2)$ or $(E_2 || \ldots || E_1)$ where $E_1$ is an unexpanded pack is the disjunction of the normal forms of $E_2$ and $(E_1 || \ldots)$.

In an expression of the form $(E$ op $\ldots)$ where $op$ is $||$ or $\&\&$, $E$ is replaced by its normal form.

The normal form of any other expression $E$ is the atomic constraint whose expression is $E$ and whose parameter mapping is the identity mapping.

**Partial ordering by constraints**

A constraint $P$ subsumes a constraint $Q$ if and only if, for every disjunctive clause $P_i$ in the disjunctive normal form of $P$, $P_i$ subsumes every conjunctive clause $Q_j$ in the conjunctive normal form of $Q$, where

- a disjunctive clause $P_i$ subsumes a conjunctive clause $Q_j$ if and only if there exists an atomic constraint $P_{ia}$ in $P_i$ for which there exists an atomic constraint $Q_{ja}$ in $Q_j$ such that $P_{ia}$ subsumes $Q_{ja}$, and

- an atomic constraint $A$ subsumes another atomic constraint $B$ if and only if $A$ and $B$ are identical using the rules described in [temp.constr.atomic].

- $A$ is a fold-expression of the form $(P$ $\&\&\ldots)$, $B$ is a fold-expression of the form $(Q$ $\&\&\ldots)$ or $(Q$ $||\ldots)$ and $P$ subsumes $Q$. 
- Otherwise, \( A \) is a *fold-expression* of the form \((P \mid \ldots)\) or \((P \&\ldots)\) and \( B \) is a *fold-expression* of the form \((Q \mid \ldots)\) and \( P \) subsumes \( Q \).

- Otherwise, \( A \) and \( B \) are identical using the rules described in [temp.constr.atomic].

*Example:* Let \( A \) and \( B \) be atomic constraints [temp.constr.atomic]. The constraint \( A \land B \) subsumes \( A \), but \( A \) does not subsume \( A \land B \). The constraint \( A \) subsumes \( A \lor B \), but \( A \lor B \) does not subsume \( A \). Also note that every constraint subsumes itself. — *end example*

### Acknowledgments

Thanks to Robert Haberlach for creating the original Concept TS issue.

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### References

