

std::generator: Synchronous Coroutine Generator for Ranges

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Abstract

We propose a standard library type `std::generator` which implements a coroutine generator that models `std::ranges::input_range`.

Acknowledgements

I'd like to thank Lewis Baker and Corentin Jabot, whose P2168 [3] I looted shamelessly for this proposal. This paper is presented as a new revision of their unpublished D2168R4 for continuity of design.

Revisions

P2502R2

- Rebase onto N4910.
- Change directions to the Editor from blue "Drafting Note" to magenta "Editor's note" in square brackets. Hopefully this makes them distinctive so as not to be confused with the actual wording changes.
- Changes to wording for `elements_of`:
 - Don't present as a subclause that holds a single subclause; flatten.
 - Don't reuse the same input range in the example, which could result in undefined behavior with some input ranges. Instead add a `bool` template parameter that selects between the two different uses with `if constexpr`.
 - Annotate member `range` with `[[no_unique_address]]` because a range type *could* be empty even if they typically are not.
- Editorial clarifications in `[coroutine.generator.overview]`:
 - Change "it was returned from" to "from which it was returned."

- Strike the sentence that redundantly describes how `co_yield` works.
 - Change “element of the generator” to “element of the sequence.”
 - Rename the `iota` generator in the example to `ints` to avoid confusion with `std::iota`, `std::ranges::iota_view`, and `std::ranges::views::iota`.
- `<generator>` no longer explicitly includes `<coroutine>` and `<ranges>`.
 - Rename `generator`’s first template parameter from `R` to `Ref` to avoid confusion with range parameters which are usually `R`.
 - `generator` publicly derives from `ranges::view_interface` instead of specializing `ranges::enable_view`.
 - Unify all “Mandates” conditions on `generator`’s template parameters into a single bulleted list.
 - Express the type requirement on `generator`’s `Allocator` template parameter as a precondition.
 - Allocators must use real pointers; both `generator`’s template parameter `Allocator` and any allocator used in an operator `new`.
 - Replace all occurrences (all instances?) of “instance” in the wording (primarily “generator instance” becomes “generator object”).
 - `generator` moves don’t invalidate iterators.
 - Simplify the specification of `generator`’s assignment operator: it now takes a `generator` argument by-value and swaps `coroutine` handles between `*this` and the argument. (Usage should be equivalent to the prior formulation.)
 - “Initial suspend point” has no hyphens, and is defined in `[dcl.fct.def.coroutine]`.
 - We don’t need to emphasize that the stack is empty in `begin`.
 - Strike a couple of `friend sometype` declarations.
 - Strike exposition-only constructors.
 - `Coroutine` handles do not “denote” `coroutines`, they “refer to” `coroutines` (`[expr.await]/3.5` is the first of many usages).
 - `get_return_object` returns `generator(...)` instead of `generator{...}`.
 - Clarify behavior of `promise_type::final_suspend`, especially when the stack of associated `coroutine` handles becomes empty.
 - Rename template parameter `Alloc2` of `promise`’s `yield_value(elements_of)` member function template to `Alloc`.
 - Massively simplify and un-template the `yield_value` overload that takes `const remove_reference_t<yielded>&`. (We don’t need to artificially template to ensure it’s distinct from

`yield_value(yielded)`, the added constraints make it distinct even when `yielded` is a `const lvalue reference`.)

- In the description of the `yield_value` overload that takes `elements_of<generator<...>`, remove the metasyntactic variable `r` that refers to “an lvalue denoting the generator referenced by `g.range`”. `g.range` is already such an lvalue, just use it directly. Use the language “stored exception” instead of “captured exception” to avoid any possible confusion with lambda captures. Strike the “Variables with automatic storage duration” note as extraneous.
- In the description of the `yield_value` overload that takes `elements_of<Rng>`, similarly use `r.range` directly instead of defining `rng` as an identical lvalue. In the “effects equivalent to” code, define a local name for `range_reference_t<Rng>` and use that instead of `auto` and `decltype(e)`.
- Specify that *yield-expression* s that call one of the `promise::yield_value` overloads have type `void`.
- Corrected spelling of `__STDCPP_DEFAULT_NEW_ALIGNMENT__`.
- Remove extraneous `static` from all declarations of class-member `operator new` and `operator delete`; class-member `operator new` and `operator delete` are static members regardless.
- Merge the specifications of the `operator new` overloads into a single specification.
- Fix precondition on `iterator::operator*()` and `iterator::operator++()`.
- Remove precondition from `iterator::operator++(int)` which is extraneous thanks to “Effects: equivalent to”.
- Make `operator==` a hidden friend instead of a member for consistency with the rest of the library. Also use “Effects: equivalent to” for `operator==` to avoid explicit preconditions.
- Add “top of stack coroutine is suspended” preconditions to `iterator::operator*` (explicitly) and `iterator::operator++` (implicit via effects-equivalent-to).
- Since all “parameter preview” arguments to `operator new` are lvalues, `operator new` overloads now take the allocator, implicit object, and “ignored” parameters by `const&`.
- Make the “stack of associated coroutines” an exposition-only `unique_ptr<stack<coroutine_handle<>>>` member of `generator`.
- Don’t use “Preconditions” for a template parameter type requirement.
- (Not requested by LWG) Change implementation of generic `yield_value` per Lewis’s suggestion to encourage HALO. (Yes, this overload is overspecified.)
- Add missing *Returns*: `element` to `iterator::operator=`.
- Fix “Returns: `return *this;`” for `iterator::operator++()`.
- Define the nested classes within namespace `std`.

- Update implementation.

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- The type `generator::yielded` is now publicly-visible and not exposition-only per 2021-12-13 LEWG discussion.
- Feature-test macro insertion is now ordered properly with the other wording.
- Coroutines for `generator<T&&>` can now yield lvalues; the generated element is an xvalue denoting a copy of the lvalue. This saves folks the trouble of writing `co_yield auto(lvalue)` instead of `co_yield lvalue`.
- Reorder `generator`'s template parameters (again, for the last last time). Coroutine Task Force discussions convinced all participants (including the author) that it was too big a break with the rest of the Standard Library to have the `Allocator` not be the final template parameter.
- Simplify `elements_of` into a vanilla two-element struct.
- Update implementations accordingly.
- Per 2022-01-025 LEWG direction, `generator<T>` now behaves like `generator<T&&>` instead of like `generator<const T&>` per the design in P2529R0 [9].

P2502R0

- Reorder `generator`'s template parameters. This allows the reference type to be more easily defaulted to a true reference, while still respecting requests for differing value and reference types. This preserves the previous design's ease-of-use, while providing full generality.
- Remove concerns about the $\mathcal{O}(1)$ destruction requirement for `view`, which has been relaxed by P2415R2 "What is a view?" [7].

D2168R4

- Wording improvements

P2168R3

- Wording improvements

P2168R2

- Some wording fixes
- Improve the section on allocator support
- Updated implementation

P2168R1

- Add benchmarks results and discussion about performance
- Introduce `elements_of` to avoid ambiguities when a generator is convertible to the reference type of the parent generator.
- Add allocator support
- Symmetric transfer works with generators of different value / allocator types
- Remove `iterator::operator->`
- Put generator in a new `<generator>` header.
- Add an other example to motivate the `Value` template parameter

Example

```
std::generator<int> fib() {
    auto a = 0, b = 1;
    while (true) {
        co_yield std::exchange(a, std::exchange(b, a + b));
    }
}

int answer_to_the_universe() {
    auto rng = fib() | std::views::drop(6) | std::views::take(3);
    return std::ranges::fold_left(std::move(rng), 0, std::plus{});
}
```

Motivation

C++ 20 had very minimalist library support for coroutines. Synchronous generators are an important use case for coroutines, one that cannot be supported without the machinery presented in this paper. Writing an efficient recursive generator is non-trivial; the standard should provide one.

Design

While the proposed `std::generator` interface is fairly straightforward, a few decisions are worth pointing out.

`input_view`

`std::generator` is a move-only view which models `input_range` and has move-only iterators. This is because the coroutine state is a unique resource (even if the coroutine *handle* is copyable).

Header

Multiple options are available as to where to put the generator class.

- `<coroutine>`, but `<coroutine>` is a low level header, and `generator` depends on bits of `<type_traits>` and `<iterator>`.
- `<ranges>`
- A new `<generator>`

This paper uses a new `<generator>` header since P2168R3 did so, and LEWG has provided no guidance to do otherwise. We do note on our MSVC STL branch implementation that `#include<ranges>` includes 52.6k lines of code, and `#include<generator>` 53.3k lines. (Note that `<generator>` is specified to include both `<ranges>` and `<coroutines>`.) Defining `generator` in `<ranges>` together with a `#include<coroutine>` would penalize people who want `<ranges>` but not `generator` by about 610 LoC.

Reference type

`generator` has 3 template parameters: `generator<R, V = void, Allocator = void>`

From `R` and `V`, we derive types:

```
using Value = conditional_t<is_void_v<V>, remove_cvref_t<R>, V>;  
using Reference = conditional_t<is_void_v<V>, R&&, R>;  
using Yielded = conditional_t<is_reference_v<Reference>, Reference, const Reference&>;
```

- `Value` is a cv-unqualified object type that specifies the value type of the generator's range and iterators,
- `Reference` specifies the reference type (not necessarily a core language reference type) of the generator's range and iterators, and
- `Allocator` is the type of allocator used for the coroutine state, which can be `void` to type-erase any allocator specified as a coroutine argument, defaulting to `allocator<byte>` when none is specified.
- `Yielded` (necessarily a reference type) is the type of the parameter to the primary overload of `yield_value` in the generator's associated promise type.

generator<meow>

Our expectation is that 98% of use cases will need to specify only one parameter. The resulting generator:

- has a value type of `remove_cvref_t<meow>`
- has a reference type of `meow`, if it is a reference type, or `meow&&` otherwise,
- expects `co_yield` to appear in the body of the generator with operands that are convertible to the reference type, and

- can use any allocator (via type-erasure) defaulting to `allocator<byte>`.

This avoids the performance pitfall from earlier revisions of the proposal that used the first argument type directly as reference type; users who naively chose `generator<std::string>` got an iterator that produces independent copies of the yielded value on every dereference, when they may have been satisfied by yielding a reference to the same constant value.

As a nod to ease of use, and to more readily reflect the semantics of prvalue-returning functions, we allow yielding lvalues from generators with an rvalue reference type when the lvalue is copyable. The result is that the yielded lvalue is copied into temporary storage, and an iterator for the generator will return an xvalue denoting that copy when dereferenced.

generator<meow, woof>

For the rare user who needs `generator` to step outside the box and use a proxy reference type, or who needs to generate a range whose iterators yield prvalues for whatever reason, we have two-argument `generator`. If `woof` is `void`, this is `generator<meow>`. Otherwise, the resulting `generator`:

- has a value type of `woof`,
- has a reference type of `meow`,
- expects `co_yield` to appear in the body of the generator with operands that are convertible to `meow`, if it is reference type, and otherwise `const meow&`, and
- can use any allocator (via type-erasure) if `woof` is `void`, or otherwise can use any allocator convertible to `woof`.

Your iterators can yield a prvalue, but it must be a prvalue of `copy_constructible` type so a copy of the operand of a single `co_yield` can be returned multiple times from repeated dereferences of the same iterator value.

generator<meow, woof, quack>

For use cases that want to specify an allocator type statically so they need not constantly pass pairs of `allocator_arg`, `my_allocator` arguments to every coroutine, we have three-argument `generator`. The resulting `generator` can use any allocator convertible to `quack`, defaulting to a default-constructed `woof` if it is `default_initializable`. The value, reference, and yielded types are as described for `generator<meow>` if `woof` is `void`, and as described for `generator<meow, woof>` otherwise.

Our expectation is that libraries that wish to declare many functions with the same statically-specified allocator type will define a template alias like

```
template <class R, class V = void>
using my_generator = generator<R, V, my_allocator>;
```

to ease declarations. In practice, this should mean `generator<R, void, my_allocator>` (with an explicit second argument of `void`) appears quite rarely in real code.

Allocator = void

In the 2022-01-13 Coroutine Task Force meeting, Lewis Baker brought up that LEWG feedback on some revision of P2168 had expressed dissatisfaction with the use of `void` to indicate “type-erase any allocator.” We discussed some alternatives:

- Devise a new type for this specific purpose. We felt that single-use tag types have low return-on-investment; every such tag type consumes a few more bytes of memory in nearly every compilation, and a few more neurons of memory in every programmer.
- Reuse an existing type for this purpose. `void` has the advantage that it can never be confused for a conforming allocator type, even if it someday does become a regular type. `void` requires no predeclaration, and no memory is expended to store its definition.

`monostate` was another suggestion that was universally reviled; that type already has a very clear purpose in the Standard Library that we didn’t feel inclined to muddle.

No one suggested another type that we found to be an improvement.

Our conclusion was that we didn’t find the use of `void` for this purpose confusing, we were happy to keep it given that more-recent LEWG discussion had resulted in no further feedback on this topic.

Obsolete discussion about reference specification

[*Note:* Before P2502R0, generator’s first parameter `Type` denoted the reference type of the range / iterators, and the value type was defaulted to `remove_cvref_t<Type>`. The following sections of design discussion from that era are preserved here. — *end note*]

In earlier versions of this paper, the reference type was exactly the first template parameter. This had the advantage of being simple. But it was a terrible performance trap:

Consider the behavior of the following code assuming the reference type is exactly the first template argument:

```
std::generator<std::string> f() {
    std::string hello = "hello";
    co_yield hello; // 0 or 1 copy depending on implementation
    co_yield "Hello"; // 1 copy (conversion from const char* to std::string)
}

for (auto&& str : f()) {} // 1 copy (*it returns std::string)
```

Of course the solution, which we advocated for, is for the user to manually specify an explicit reference type:

```
generator<const std::string&> f() {
    std::string hello = "hello";
    co_yield hello; // 0 or 1 copy depending on implementation
    co_yield "Hello"; // 1 copy (conversion from const char* to std::string)
}
```



```
for (auto&& str : f()) {} // 0 copy
```

This works, can be explained, and is even logical. You get what you asked for. It is nonetheless surprising for non-experts that using the simple `generator<string>` would create 2 copies per `co_yield`.

To hope users would not routinely forget to use a reference type when using `std::generator` calls for a heaping barrel of optimism.

We later proposed that for a `generator<T>`, its reference type be `conditional_t<is_reference_v<T>, T, const T&>`.

First parameter	reference type	default value	can yield mutable lvalue ref?
<code>int</code>	<code>const int&</code>	<code>int</code>	No
<code>const int&</code>	<code>const int&</code>	<code>int</code>	No
<code>int&</code>	<code>int&</code>	<code>int</code>	Yes
<code>int&&</code>	<code>int&&</code>	<code>int</code>	No
<code>const int&&</code>	<code>const int&&</code>	<code>int</code>	No

Attempts have been made to characterize the exact relations between reference, value, storage, and `co_yield` exception types and categories. Ultimately, a simpler mental model is to characterize what expressions can be yielded for a given reference type and how many copies are made for each scenario.

First parameter	<code>co_yield const T&</code>	<code>co_yield T&</code>	<code>co_yield T&&</code>	<code>co_yield U&&</code>
<code>T</code>	0	0	0	1
<code>const T&</code>	0	0	0	1
<code>T&</code>	Ill-formed	0	Ill-formed	Ill-formed
<code>T&&</code>	Ill-formed	Ill-formed	0	1
<code>const T&&</code>	Ill-formed	Ill-formed	0	1

In this table, we see that only `co_yield` that requires conversion incurs copy, which is expected. Coroutines guarantee that the yielded expression exceeds the lifetime of the `co_yield` expression, so `generator` can usefully store a pointer to the object denoted by a yielded xvalue.

`co_yield` expressions involving conversion can store the yielded value in an awaiter. The type of the stored expression is the reference type with its reference qualifiers stripped, but that is an implementation detail that is not observable and is therefore of limited interest. Of course, that type needs to be constructible from yielded values.

Besides the `T` case, this behaves very much like returning from a function that is intended.

Move-only and immovable types

LEWG was interested in how this works with `generator` of move-only and immovable types.

First parameter	co_yield const T&	co_yield T&	co_yield T&&
move_only	0	0	0
const move_only&	0	0	0
move_only&	Ill-formed	0	Ill-formed
move_only&&	Ill-formed	Ill-formed	0
const move_only&&	Ill-formed	Ill-formed	0
immovable	0	0	0
const immovable&	0	0	0
immovable&	Ill-formed	0	Ill-formed
immovable&&	Ill-formed	Ill-formed	0
const immovable&&	Ill-formed	Ill-formed	0

As that table shows, these types work exactly like other types. However, to be able to move from a move only reference type, the coroutine has to explicitly state so:

```
auto f = []() -> std::generator<move_only> { co_yield move_only{}; }();
for (auto&& x : f) {
    move_only mo = std::move(x); // ill-formed, decltype(x) is const move_only&
}
```

```
auto f = []() -> std::generator<move_only&&> { co_yield move_only{}; }();
for (auto&& x : f) {
    move_only mo = x; // ok
}
```

```
auto f = []() -> std::generator<move_only&> { move_only m; co_yield m; }();
for (auto&& x : f) {
    move_only mo = std::move(x); // dicey but okay
}
```

Potential downsides

```
auto f = []() -> std::generator<MyType> {
    MyType t;
    co_yield std::move(t);
}();
```

In the example above `std::move` doesn't move. Arguably more than usual. Indeed the code expands to something similar to:

```
auto&& __temp = std::move(t);
yield_value(__temp); // <=> promise.value = std::addressof(__temp); // no move
```

Of course, a move would not have occurred for a `std::generator<const MyType&>` either as these things are identical. It might be surprising? The only way to avoid that is to create temporary value for rvalue reference, which would force a move to actually occurs, at the cost

of performance.

Alternatives considered

Mandating a reference as the first parameter We could make `generator<int>` ill-formed and force people to specify a reference type like `generator<const int&>`. We do not think this is very user-friendly, given that we can provide a reasonable default.

We rejected this option.

Using T& as the default There are two issues with mutable references:

- They are mutable (They allow mutating the coroutine frame), which would be an *interesting* default.
- They are very restrictive as to the set of `co_yield` expression allowed with them.

We rejected this option.

Using T&& as the default This avoids a copy when doing `auto object = *it` (where it is a `std::generator::iterator`), but it is easy to misuse, consider:

```
auto f = [] -> std::generator<std::string> { co_yield "footgun"; }();
for (auto&& x : f | std::views::filter([](std::string s) { return s.size() > 0; })) {
    std::cout << x << '\n'; // outputs a single empty line
}
```

We rejected this option. [Note: Nevertheless, we came back to this option after much discussion of pros and cons in both LEWG and several additional meetings if a Coroutine Task Force formed by LEWG Chair specifically to build consensus. — end note]

Doing something clever for move-only types We considered returning T& for `move_only` types so that they can be moved from by default. We realized this was too clever and inconsistent. Notably, adding a copy constructor to T would change the meaning of the code.

We rejected this option.

Doing something clever for reference types By default `generator<reference_wrapper<T>>` could yield `reference_wrapper<T>` has that is already a “reference-like” type. However, no other view does that, “reference-like” is fuzzily defined, and this would probably cause more trouble than it’s worth.

We rejected this option.

Keeping the D2168R4 design Returning values has the potential to severely impact performance, is inconsistent with other views, and is not necessary. It also did not work with move-only types.

The change, along with an implementation strategy described in the “How to store the yielded value in the promise type?” guarantees that no copy needs to be made if the reference and yielded types are the same (with qualifiers stripped).

We think this new approach keeps the simplicity of the original design, improves performance, and works with more types.

Thank you LEWG, and in particular Mathias, for highlighting these concerns!

Separately specifyable Value Type

This proposal supports specifying both the “yielded” type, which is the iterator’s reference type (not required to be a reference) and its corresponding value type. This allow ranges to handle proxy types and wrapped reference, like this implementation of zip:

```
namespace ranges = std::ranges;

template<ranges::input_range Rng1, ranges::input_range Rng2>
std::generator<
    std::tuple<ranges::range_reference_t<Rng1>, ranges::range_reference_t<Rng2>>
    std::tuple<ranges::range_value_t<Rng1>, ranges::range_value_t<Rng2>>>
zip(Rng1 r1, Rng2 r2) {
    auto it1 = ranges::begin(r1);
    auto it2 = ranges::begin(r2);
    auto end1 = ranges::end(r1);
    auto end2 = ranges::end(r2);
    for (; it1 != end1 && it2 != end2; ++it1, ++it2) {
        co_yield {*it1, *it2};
    }
}
```

In this second example, using `string` as value type ensures that calling code can take the necessary steps to make sure iterating over a generator would not invalidate any of the yielded values.

```
// Yielding string literals : always fine
std::generator<std::string_view, std::string_view> string_views() {
    co_yield "foo";
    co_yield "bar";
}

std::generator<std::string_view, std::string> strings() {
    co_yield "start";
    std::string s;
    for (auto sv : string_views()) {
        s = sv;
        s.push_back('!');
        co_yield s;
    }
    co_yield "end";
}
```

```

// conversion to a vector of strings
// If the value_type was string_view, it would convert to a vector of string_view,
// which would lead to undefined behavior operating on elements of v that were
// invalidated while iterating through the generator.
auto v = std::ranges::to<vector>(strings()); // (P1206R3 [4])

```

How to store the yielded value in the promise type?

There are multiple implementation strategies possible to store the value in the generator. An early revision of this paper always stored a copy of the yielded value, leading to an extra copy. Later revisions supported storing the yielded value in an awaitable object returned from the promise's `yield_value` function.

However, the object denoted by a glvalue yield expression is guaranteed to live until the coroutine resumes. We can take advantage of that fact by storing only a pointer to the denoted object in the promise, if the result of dereferencing that pointer is convertible to the generator's reference type. We guarantee this is the case by providing a `yield_value` whose parameter type is always a reference type (`conditional_t<is_reference_v<Reference>, Reference, const Reference&>`). This forces any conversions to happen inside the coroutine itself, yielding a temporary glvalue that can later be dereferenced to an lvalue which is trivially `static_casted` to `Reference` in the iterator's `operator*`.

A drawback of this solution is that the yielded value is only destroyed at the end of the full expression in which `co_yield` appears, so given

```
(co_yield x, co_yield y); // x is destroyed after y is yielded.
```

We think this is a reasonable tradeoff given that this approach minimizes the number of copies must be made of the yielded value. We force the coroutine to materialize the element to be yielded, but after doing so can cleanly pass a reference to that element through the coroutine and iterator machinery and directly to consuming code.

Recursive generator

A "recursive generator" is a coroutine that supports the ability to directly `co_yield` a generator of the same type as a way of emitting the elements of that generator as elements of the current generator.

Example: A generator can `co_yield` other generators of the same type

```

std::generator<const std::string&> delete_rows(std::string table, std::vector<int> ids) {
    for (int id : ids) {
        co_yield std::format("DELETE FROM {0} WHERE id = {1};", table, id);
    }
}

std::generator<const std::string&> all_queries() {
    co_yield std::ranges::elements_of(delete_rows("user", {4, 7, 9 10}));
}

```

```

    co_yield std::ranges::elements_of(delete_rows("order", {11, 19}));
}

```

Example: A generator can also be used recursively

```

using namespace std;

struct Tree {
    Tree* left;
    Tree* right;
    int value;
};

generator<int> visit(Tree& tree) {
    if (tree.left) co_yield ranges::elements_of(visit(*tree.left));
    co_yield tree.value;
    if (tree.right) co_yield ranges::elements_of(visit(*tree.right));
}

```

In addition to being more concise, the ability to directly yield a nested generator has some performance benefits compared to iterating over the contents of the nested generator and manually yielding each of its elements.

Yielding a nested generator allows the consumer of the top-level coroutine to directly resume the current leaf generator when incrementing the iterator, whereas a solution that has each generator manually iterating over elements of the child generator requires $O(\text{depth})$ coroutine resumptions/suspensions per element of the sequence.

Example: Non-recursive form incurs $O(\text{depth})$ resumptions/suspensions per element and is more cumbersome to write:

```

using namespace std;

generator<int> slow_visit(Tree& tree) {
    if (tree.left) {
        for (int x : ranges::elements_of(visit(*tree.left)))
            co_yield x;
    }
    co_yield tree.value;
    if (tree.right) {
        for (int x : ranges::elements_of(visit(*tree.right)))
            co_yield x;
    }
}

```

Exceptions that propagate out of the body of nested generator coroutines are rethrown into the parent coroutine from the `co_yield` expression rather than propagating out of the top-level iterator's `operator++()`. This follows the mental model that `co_yield someGenerator` is semantically equivalent to manually iterating over the elements and yielding each element.

For example: `nested_ints()` is semantically equivalent to `manual_ints()`

```

std::generator<int> might_throw() {
    co_yield 0;
    throw some_error{};
}

std::generator<int> nested_ints() {
    try {
        co_yield std::ranges::elements_of(might_throw());
    } catch (const some_error&) {}
    co_yield 1;
}

// nested_ints() is semantically equivalent to the following:
std::generator<int> manual_ints() {
    try {
        for (int x : might_throw()) {
            co_yield x;
        }
    } catch (const some_error&) {}
    co_yield 1;
}

void consumer() {
    for (int x : nested_ints()) {
        std::cout << x << " "; // outputs 0 1
    }

    for (int x : manual_ints()) {
        std::cout << x << " "; // also outputs 0 1
    }
}

```

std::ranges::elements_of

`ranges::elements_of` is a utility function that prevents ambiguity when a nested generator type is convertible to the value type of the present generator

```

generator<int> f()
{
    co_yield 42;
}

generator<any> g()
{
    co_yield f(); // should we yield 42 or generator<int> ?
}

```

To avoid this issue, we propose that:

- `co_yield <expression>` yields the value directly, and

- `co_yield elements_of(<expression>)` yields successive elements from the nested generator.

For convenience, we further propose that `co_yield elements_of(x)` be extended to support yielding the values of arbitrary ranges beyond generators, ie

```
std::generator<int> f()
{
    std::vector<int> v = /*... */;
    co_yield std::ranges::elements_of(v);
}
```

Symmetric transfer

The recursive form can be implemented efficiently with symmetric transfer. Earlier works in [CppCoro] implemented this feature in a distinct `recursive_generator` type.

However, it appears that a single type is reasonably efficient thanks to HALO optimizations and symmetric transfer. The memory cost of that feature is two extra pointers per generator¹. It is difficult to evaluate the runtime cost of our design given the current coroutine support in compilers. However our tests show no noticeable difference between a generator and a `recursive_generator` which is called non-recursively. It is worth noting that the proposed design makes sure that HALO [8] optimizations are possible.

While we think a single generator type is sufficient and offers a better API, there are three options:

- A single `generator` type supporting recursive calls (this proposal).
- A separate type `recursive_generator` that can yield values from either a `recursive_generator` or a `generator`. That may offer very negligible performance benefits, same memory usage.
- A separate `recursive_generator` type which can only yield values from other `recursive_generator`s.

That third option would make the following ill-formed:

```
generator<int> f();
recursive_generator<int> g() {
    co_yield f(); // incompatible types
}
```

Instead you would need to write:

```
recursive_generator<int> g() {
    for (int x : f()) co_yield x;
}
```

¹The two pointers in our implementation have non-overlapping active times; we believe the pair can be optimized into a single pointer's space with some bit hacking to store a discriminator in the unused lower bits.

Such a limitation can make it difficult to decide at the time of writing a generator coroutine whether or not you should return a `generator` or `recursive_generator` as you may not know at the time whether or not this particular generator will be used within `recursive_generator` or not.

If you choose the `generator` return-type and then later someone wants to yield its elements from a `recursive_generator` then you either need to manually yield its elements one-by-one or use a helper function that adapts the generator into a `recursive_generator`. Both of these options can add runtime cost compared to the case where the generator was originally written to return a `recursive_generator`, as it requires two coroutine resumptions per element instead of a single coroutine resumption.

Because of these limitations, we are not recommending this approach.

Symmetric transfer is possible for different generator types as long as the reference type is the same, aka, different value type or allocator type does not preclude symmetric transfer (see the section on allocators).

Allocator support

In line with the design exploration done in section 2 of [P1681R0 \[6\]](#), `std::generator` supports both stateless and stateful allocators and strives to minimize the interface verbosity for stateless allocators by templating both the generator itself and the `promise_type`'s `new` operator on the allocator type. Details for this interface are found in [P1681R0 \[6\]](#).

`coroutine_parameter_preview_t` such as discussed in section 3 of [P1681R0 \[6\]](#) has not been explored in this paper.

```
std::generator<int> stateless_example() {
    co_yield 42;
}

template <class Allocator>
std::generator<int> allocator_example(std::allocator_arg_t, Allocator alloc) {
    co_yield 42;
}

my_allocator<std::byte> alloc;
input_range auto rng = allocator_example(std::allocator_arg, alloc);
```

The proposed interface requires that, if an allocator is provided, it is the second argument to the coroutine function, immediately preceded by an argument of type `std::allocator_arg_t`. This approach is necessary to distinguish the allocator desired to allocate the coroutine state from allocators whose purpose is to be used in the body of the coroutine function. The required argument order might be a limitation if any other argument is required to be the first. However, we cannot think of any scenario where that would be the case.

We think it is important that all standard and user coroutine types can accommodate similar interfaces for allocator support. In fact, the implementation for that allocator support can be

shared amongst `generator`, `lazy`, and other standard types.

By default `std::generator` type erases the allocator type, and uses `std::allocator` unless an allocator is provided to the coroutine function. Then:

Type erased allocator(default)

```
template <class Allocator>
std::generator<int> f(std::allocator_arg_t, Allocator alloc) {}

f(std::allocator_arg, my_alloc{});
```

Returns a generator of type `std::generator<int, void, void>` where the final `void` denotes that the allocator is type erased. The allocator is stored in the same allocation as the coroutine state if it is stateful or not default constructible; a pointer is always stored so that the `deallocate` method of the type erased allocator can be called.

No allocator

```
std::generator<int> f() {}

f();
```

Again, returns a generator of type `std::generator<int, void, void>` where the final `void` denotes that the allocator is type erased. A pointer is stored so that the `deallocate` method of the type-erased allocator can be called, but the default allocator (`std::allocator`) need not be stored since it is stateless.

Explicit stateless allocator

```
std::generator<int, void, std::stateless_allocator<int>> f() {}

f();
```

No extra storage is used for the allocator because it is stateless.

Explicit stateful allocator

```
std::generator<int, void, some_stateful_allocator<int>>
    f(std::allocator_arg_t, some_stateful_allocator<int> alloc) {}
f(std::allocator_arg, some_allocator); // must be convertible to some_stateful_allocator
```

The allocator is copied in the coroutine state.

Can we postpone adding allocator support?

A case can be made that allocator support could be added to `std::generator` later. However, because the proposed design has the allocator as a template parameter, adding allocator after `std::generator` ships would represent an ABI break. We recommend that we add allocator support as proposed in this paper now and make sure that the design remains consistent as work on `std::lazy` is made in this cycle. However, it would be possible to extend support for different mechanisms (such as presented in section 3 of [P1681R0 \[6\]](#) later.

Interaction of symmetric transfer and allocator support

The allocator must necessarily be part of a coroutine's promise type since implementations query the promise for allocation functions. Nonetheless, it would seem silly for a generator to be unable to nest another generator with identical element type but differing allocator. For that matter, even differing value types shouldn't be problematic: the only interface between the generator and the coroutine it wraps that differs depending on the type arguments to generator is `yield_value`. Ideally, generators would be able to recurse into other generators whose `yield_value` has the same parameter type even if all three template arguments to generator differ.

Our implementation uses a base class to implement the non-allocation behaviors for generator's promise type so that generators with different allocator types can yield each other. Doing so, however, requires that we partially erase the type of a `coroutine_handle` so we can resume it later knowing only that its promise type derives from a particular base.

There are at least two ways to implement this partial type erasure:

- Storing a pointer in the common base to a component with full type knowledge, which can then resume the targeted coroutine,
- Relax the preconditions on some of the `coroutine_handle` functions to allow conversion from `coroutine_handle<void>` to `coroutine_handle<T>` when the source's corresponding `address()` value was obtained from a `coroutine_handle` referring to a coroutine whose promise object is pointer-interconvertible with an object of type `T`.

Our current plan is to standardize the intent to allow yielding nested generators with different allocator and value types, leaving the details of the implementation unspecified, and to later separately propose the changes to `coroutine_handle` that enable that implementation to be maximally efficient.

Implementation and experience

`generator` has been provided as part of `cppcoro` and `folly`. However, `cppcoro` offers a separate `recursive_generator` type, which is different than the proposed design.

Folly uses a single `generator` type, which can be recursive but doesn't implement symmetric transfer. Despite that, Folly users found the use of `Folly::Generator` to be a lot more efficient than the eager algorithm they replaced with it.

`ranges-v3` also implements a `generator` type, which is never recursive and predates the work on move-only views and iterators [1], [2] which forces this implementation to ref-count the coroutine handler.

Our implementation [Implementation] consists of a single type that takes advantage of symmetric transfer to implement recursion - it notably works well with three different major standard libraries.

Performance & benchmarks

[*Note*: These benchmark results are fairly dated now - roughly a year old - and should be taken with a grain of salt. — *end note*]

Because implementations are still being perfected, and because performance is extremely dependant on whether HALO optimization (see [P0981R0](#) [8]) occurs, it is difficult at this time to make definitive statements about the performance of the proposed design.

At the time of the writing of this paper, Clang is able to inline non-nested coroutines whether the implementation supports nested coroutines or not, while GCC never performs HALO optimization.

When the coroutine is not inlined, support for recursion does not noticeably impact performance. And, when the coroutine yields another generator, the performance of the recursive version is noticeably faster than yielding each element of the range. This is especially noticeable with deep recursion.

	Clang	Clang ST ¹	GCC	GCC ST ¹	MSVC	MSVC ST ¹
Single value	(1) 0.235	(2) 2.36	12.4	13.4	61.9	63.7
Single value, noinline (3)	13.5	13.7	14.1	15.2	63.8	64.4
Deep nesting	43670266.0	(4) 427955.0	58801348	338736	224052033	4760914

¹ Symmetric transfer.

The values are expressed in nanoseconds. However, please note that the comparison of the same result across compiler is not meaningful, notably because the MSVC results were obtained on different hardware. That being said, we observe:

- Only Clang can perform constant folding of values yielded by simple coroutine (1)
- When the generator supports symmetric transfer, clang is not able to fully inline the generator construction, but HALO is still performed (2).
- When HALO is not performed, the relative performance of both approaches is similar (3).
- Supporting recursion is greatly beneficial to nested/recursive algorithms (4).

The code for these benchmarks, as well as more detailed results, can be found on [Github](#).

Wording

[Editor's note: Wording is relative to Working Draft N4910 [5].

Insert in lexicographical order in [version.syn] (updating YYYYXXL to the date of merge):]

```
#define __cpp_lib_generator                YYYYXXL // also in <generator>
```

[Editor's note: Modify [ranges.general] as follows:]

◆ **General** [ranges.general]

This Clause describes components for dealing with ranges of elements.

The following subclauses describe range and view requirements, and components for range primitives [and range generators](#) as summarized in Table [tab:range.summary].

[Editor's note: Add a new row at the end of [tab:range.summary] "Range generators" with header <generator> referring to the new subclause [coroutine.generator] added below.]

[Editor's note: Add the declaration of `ranges::elements_of` to the <ranges> synopsis:]

◆ **Header <ranges> synopsis** [ranges.syn]

```
namespace std::ranges {  
[...]  
  
    template<input_or_output_iterator I, sentinel_for<I> S, subrange_kind K>  
    inline constexpr bool enable_borrowed_range<subrange<I, S, K>> = true;  
  
    // [range.dangling], dangling iterator handling  
    struct dangling;  
  
    // [elementsof.overview], class template elements_of  
    template<range R, class Allocator = allocator<byte>>  
    struct elements_of;  
  
    template<range R>  
    using borrowed_iterator_t = conditional_t<borrowed_range<R>, iterator_t<R>, dangling>;  
  
[...]  
}
```

[Editor's note: Insert the following new subclause immediately after [range.dangling]:]

◆ **class template elements_of** [ranges.elementsof]

Specializations of `elements_of` encapsulate a range and act as a tag in overload sets to disambiguate when a range should be treated as a sequence rather than a single value.

[Example:

```

template <bool YieldElements>
std::generator<any> f(std::ranges::input_range auto&& rng) {
    if constexpr (YieldElements)
        co_yield std::ranges::elements_of(rng); // yield each element of rng
    else
        co_yield rng; // yield rng as a single value
}

```

— *end example*]

```

namespace std::ranges {
    template<range R, class Allocator = allocator<byte>>
    struct elements_of {
        [[no_unique_address]] R range;
        [[no_unique_address]] Allocator allocator = Allocator();
    };

    template<class R, class Allocator = allocator<byte>>
    elements_of(R&&, Allocator = Allocator()) -> elements_of<R&&, Allocator>;
}

```

[Editor's note: Add the following subclause to the end of [ranges]:]

 **Range Generators** [coroutine.generator]

 **Overview** [coroutine.generator.overview]

Class template `generator` presents a view of the elements yielded by the evaluation of a coroutine.

A generator generates a sequence of elements by repeatedly resuming the coroutine from which it was returned. Elements of the sequence are produced by the coroutine each time a `co_yield` statement is evaluated. When the `co_yield` statement is of the form `co_yield elements_of(rng)`, each element of the range `rng` is successively produced as an element of the sequence.

[*Example:*

```

std::generator<int> ints(int start = 0) {
    while (true)
        co_yield start++;
}

void f() {
    for (auto i : ints() | std::views::take(3))
        std::cout << i << ' '; // prints 0 1 2
}

```

— *end example*]

◆ Header <generator> synopsis

[generator.syn]

```
namespace std {
    // [coroutine.generator.class], class template generator
    template<class Ref, class V = void, class Allocator = void>
    class generator;
}
```

◆ Class template generator

[coroutine.generator.class]

```
namespace std {
    template<class Ref, class V = void, class Allocator = void>
    class generator : public ranges::view_interface<generator<Ref, V, Allocator>> {
    private:
        using value =          // exposition only
            conditional_t<is_void_v<V>, remove_cvref_t<Ref>, V>;
        using reference =      // exposition only
            conditional_t<is_void_v<V>, Ref&&, Ref>;
        class iterator;        // exposition only

    public:
        using yielded =
            conditional_t<is_reference_v<reference>, reference, const reference&>;

        class promise_type;

        generator(const generator&) = delete;
        generator(generator&& other) noexcept;

        ~generator();

        generator& operator=(generator other) noexcept;

        iterator begin();
        default_sentinel_t end() const noexcept;

    private:
        coroutine_handle<promise_type> coroutine_ = nullptr; // exposition only
        unique_ptr<stack<coroutine_handle<>>> active_; // exposition only
    };
}
```

Mandates:

- If Allocator is not void, `allocator_traits<Allocator>::pointer` is a pointer type.
- `value` is a cv-unqualified object type.
- `reference` is either a reference type, or a cv-unqualified object type that models copy-constructible.
- Let RRef denote `remove_reference_t<reference>&&` if `reference` is a reference type, or `reference` otherwise. Each of:

- `common_reference_with<reference&&, value&>`,
- `common_reference_with<reference&&, RRef&&>`, and
- `common_reference_with<RRef&&, const value&>`

is modeled. [*Note*: These requirements ensure the exposition-only *iterator* type can model `indirectly_readable` and thus `input_iterator`. — *end note*]

If `Allocator` is not void, it shall meet the `Cpp17Allocator` requirements.

Specializations of `generator` model `view` and `input_range`.

The behavior of a program that adds a specialization for `generator` is undefined.

❖ **Members** [generator.members]

```
generator(generator&& other) noexcept;
```

Effects: Initializes `coroutine_` with `exchange(other.coroutine_, {})` and `active_` with `exchange(other.active_, nullptr)`.

[*Note*: Iterators previously obtained from `other` are not invalidated; they become iterators into `*this`. — *end note*]

```
~generator();
```

Effects: Equivalent to:

```
if (coroutine_) {
    coroutine_.destroy();
}
```

[*Note*: Ownership of recursively yielded generators is held in awaitable objects in the `coroutine` frame of the yielding generator, so destroying the root generator effectively destroys the entire stack of yielded generators. — *end note*]

```
generator& operator=(generator other) noexcept;
```

Effects: Equivalent to:

```
swap(coroutine_, other.coroutine_);
swap(active_, other.active_);
```

Returns: `*this`.

[*Note*: Iterators previously obtained from `other` are not invalidated; they become iterators into `*this`. — *end note*]

```
iterator begin();
```

Preconditions: `coroutine_` refers to a `coroutine` suspended at its initial suspend point ([`dcl.fct.def.coroutine`]).

Effects: Pushes *coroutine_* into **active_*, then evaluates *coroutine_.resume()*.

Returns: An *iterator* object whose member *coroutine_* refers to the same coroutine as does *coroutine_*.

[*Note:* A program that calls `begin` more than once on the same generator has undefined behavior. — *end note*]

```
default_sentinel_t end() const noexcept;
```

Returns: `default_sentinel`.

class `generator::promise_type` **[`coroutine.generator.promise`]**

```
namespace std {
    template<class Ref, class V, class Allocator>
    class generator<Ref, V, Allocator>::promise_type {
    public:
        generator get_return_object() noexcept;

        suspend_always initial_suspend() const noexcept { return {}; }
        auto final_suspend() noexcept;

        suspend_always yield_value(yielded val) noexcept;

        auto yield_value(const remove_reference_t<yielded>& lval)
            requires is_rvalue_reference_v<yielded> &&
                constructible_from<remove_cvref_t<yielded>, const remove_reference_t<yielded>&>;

        template<class R2, class V2, class Alloc2, class Unused>
            requires same_as<typename generator<T2, V2, Alloc2>::yielded, yielded>
                auto yield_value(ranges::elements_of<generator<T2, V2, Alloc2>&&, Unused> g) noexcept;

        template<ranges::input_range Rng, class Alloc>
            requires convertible_to<ranges::range_reference_t<Rng>, yielded>
                auto yield_value(ranges::elements_of<Rng, Alloc> r) noexcept;

        void await_transform() = delete;

        void return_void() const noexcept {}

        void unhandled_exception();

        void* operator new(size_t size)
            requires same_as<Allocator, void> || default_initializable<Allocator>;

        template<class Alloc, class... Args>
            requires same_as<Allocator, void> || convertible_to<const Alloc&, Allocator>
                void* operator new(size_t size, allocator_arg_t, const Alloc& alloc, const Args&...);

        template<class This, class Alloc, class... Args>
            requires same_as<Allocator, void> || convertible_to<const Alloc&, Allocator>
```

```

    void* operator new(size_t size, const This&, allocator_arg_t, const Alloc& alloc,
                      const Args&...);

    void operator delete(void* pointer, size_t size) noexcept;

private:
    add_pointer_t<yielded> value_ = nullptr; // exposition only
    exception_ptr except_; // exposition only
};
}

generator get_return_object() noexcept;

Returns: A generator object whose member coroutine_ is coroutine_handle<promise_
type>::from_promise(*this), and whose member active_ points to an empty stack.

auto final_suspend() noexcept;

Preconditions: A handle referring to the coroutine whose promise object is *this is at
the top of *active_ of some generator object x. This function is called by that coroutine
upon reaching its final suspend point ([dcl.fct.def.coroutine]).

Returns: An awaitable object of unspecified type whose member functions arrange
for the calling coroutine to be suspended, pop the coroutine handle from the top of
*x.active_, and resume execution of the coroutine referred to by x.active_->top() if
*x.active_ is not empty. If it is empty, control flow returns to the current coroutine caller
or resumer ([dcl.fct.def.coroutine]).

suspend_always yield_value(yielded val) noexcept;

Effects: Equivalent to value_ = addressof(val).

Returns: {}.

auto yield_value(const remove_reference_t<yielded>& lval)
    requires is_rvalue_reference_v<yielded> &&
    constructible_from<remove_cvref_t<yielded>, const remove_reference_t<yielded>&&;

Preconditions: A handle referring to the coroutine whose promise object is *this is at
the top of *active_ of some generator object x.

Returns: An awaitable object of an unspecified type ([expr.await]) that stores an object of
type remove_cvref_t<yielded> direct-non-list-initialized with lval, whose member func-
tions arrange for value_ to point to that stored object and then suspend the coroutine.

Throws: Any exception thrown by the initialization of the stored object.

Remarks: A yield-expression that calls this function has type void ([expr.yield]).

template<class T2, class V2, class Alloc2, class Unused>
    requires same_as<typename generator<T2, V2, Alloc2>::yielded, yielded>
    auto yield_value(ranges::elements_of<generator<T2, V2, Alloc2>&&, Unused> g) noexcept;

```

Preconditions: A handle referring to the coroutine whose promise object is **this* is at the top of **active_* of some generator object *x*. The coroutine referred to by *g.range.coroutine_* is suspended at its initial suspend point.

Returns: An awaitable object of an unspecified type ([*expr.await*]) which takes ownership of the generator *g.range*, whose member *await_ready* returns *false*, whose member *await_suspend* pushes *g.range.coroutine_* into **x.active_* and resumes execution of the coroutine referred to by *g.range.coroutine_*, and whose member *await_resume* evaluates *rethrow_exception(except_)* if *bool(except_)* is true. If *bool(except_)* is false, the *await_resume* member has no effects.

Remarks: A *yield-expression* that calls this function has type *void* ([*expr.yield*]).

```
template<ranges::input_range Rng, class Alloc>
requires convertible_to<ranges::range_reference_t<Rng>, yielded>
auto yield_value(ranges::elements_of<Rng, Alloc> r) noexcept;
```

Effects: Equivalent to:

```
auto nested = [](allocator_arg_t, Alloc, ranges::iterator_t<Rng> i,
                 ranges::sentinel_t<Rng> s)
-> generator<yielded, ranges::range_value_t<Rng>, Alloc> {
    for (; i != s; ++i) {
        co_yield static_cast<yielded>(*i);
    }
};
return yield_value(ranges::elements_of(nested(
    allocator_arg, r.allocator, ranges::begin(r.range), ranges::end(r.range))));
```

Remarks: A *yield-expression* that calls this function has type *void* ([*expr.yield*]).

```
void unhandled_exception();
```

Preconditions: A handle referring to the coroutine whose promise object is **this* is at the top of **active_* of some generator object *x*.

Effects: If the handle referring to the coroutine whose promise object is **this* is the sole element of **x.active_*, equivalent to: *throw*; Otherwise, assigns *current_exception()* to *except_*.

```
void* operator new(size_t size)
requires same_as<Allocator, void> || default_initializable<Allocator>;
```

```
template<class Alloc, class... Args>
requires same_as<Allocator, void> || convertible_to<const Alloc&, Allocator>
void* operator new(size_t size, allocator_arg_t, const Alloc& alloc, const Args&...);
```

```
template<class This, class Alloc, class... Args>
requires same_as<Allocator, void> || convertible_to<const Alloc&, Allocator>
void* operator new(size_t size, const This&, allocator_arg_t, const Alloc& alloc,
                  const Args&...);
```

Let A be

- Allocator, if it is not void,
- Alloc for the overloads with a template parameter Alloc, or
- allocator<void> otherwise.

Let B be allocator_traits<A>::template rebind_alloc<U> where U is an unspecified type whose size and alignment are both `__STDCPP_DEFAULT_NEW_ALIGNMENT__`.

Mandates: allocator_traits::pointer is a pointer type.

Effects: Initializes an allocator b of type B with A(alloc), for the overloads with a function parameter alloc, or with A() otherwise. Uses b to allocate storage for the smallest array of U sufficient to provide storage for a coroutine state of size size, and unspecified additional state necessary to ensure that operator delete can later deallocate this memory block with an allocator equal to b.

Returns: A pointer to the allocated storage.

```
void operator delete(void* pointer, size_t size) noexcept;
```

Preconditions: pointer was returned from an invocation of one of the above overloads of operator new with a size argument equal to size.

Effects: Deallocates the storage pointed to by pointer using an allocator equivalent to that used to allocate it.



Class template generator::iterator

[coroutine.generator.iterator]

```
namespace std {
    template<class Ref, class V, class Allocator>
    class generator<Ref, V, Allocator>::iterator {
    public:
        using value_type = value;
        using difference_type = ptrdiff_t;

        iterator(iterator&& other) noexcept;

        iterator& operator=(iterator&& other) noexcept;

        reference operator*() const noexcept(is_nothrow_copy_constructible_v<reference>);

        iterator& operator++();
        void operator++(int);

        friend bool operator==(iterator i, default_sentinel_t);

    private:
        coroutine_handle<promise_type> coroutine_; // exposition only
    };
}
```

`iterator(iterator&& other) noexcept;`

Effects: Initializes `coroutine_` with `exchange(other.coroutine_, {})`.

`iterator& operator=(iterator&& other) noexcept;`

Effects: Equivalent to `coroutine_ = exchange(other.coroutine_, {})`.

Returns: `*this`.

`reference operator*() const noexcept(is_nothrow_copy_constructible_v<reference>);`

Preconditions: For some generator object `x`, `coroutine_` is in `*x.active_` and `x.active_>top()` refers to a suspended coroutine with promise object `p`.

Effects: Equivalent to: `return static_cast<reference>(*p.value_);`

`iterator& operator++();`

Preconditions: For some generator object `x`, `coroutine_` is in `*x.active_`.

Effects: Equivalent to `x.active_>top().resume()`.

Returns: `*this`.

`void operator++(int);`

Effects: Equivalent to `+++this`.

`friend bool operator==(iterator i, default_sentinel_t);`

Effects: Equivalent to: `return i.coroutine_.done();`

References

- [1] Casey Carter. P1456R1: Move-only views. <https://wg21.link/p1456r1>, 11 2019.
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