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## Invoking Algorithms Asynchronously

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## 1. Introduction

This paper describes new execution policies enabling the asynchronous execution of the parallel algorithms as defined by the Parallelism TS (N4507) [1]. This paper is part of an effort to design and propose uniform parallelism APIs in C++ with the goal to make the language independent from any external solutions (such as OpenMP or OpenACC). There have been several discussions in SG1 and SG14 during the recent committee meetings in Kona and Jacksonville expressing interest in enabling asynchronous execution of parallel algorithms.

This paper also continues the specific features and needs towards supporting Heterogeneous Devices which was discussed in an evening session at Jacksonville 2016 [2]. In that evening session, Michael Wong presented the motivation to support Heterogeneous devices and how it has been done in OpenMP, and was followed by two C++-specific designs. Hartmut Kaiser presented the HPX design which caters more to a high-performance computing viewpoint. Lee Howes presented the Khronos SYCL/OpenCL design which caters more to a consumer device viewpoint. The discussion that followed, indicated enthusiastic support to move C++ towards full support for Heterogeneous computing by 2020, likely through an initial TS.

In general, all parallel algorithms as defined in N4507 are synchronous. This means that the execution of an algorithm returns only after its operation has completely finished. It is well known, that this

form of fork-join parallelism imposes an implicit barrier onto the parallel execution flow. This is also currently the case in OpenMP parallel regions. This barrier impedes parallel efficiency and efficient resource utilization of the used processing units as the execution has to 'wait' for the thread of execution which performs the necessary join operation at the end of the execution of the algorithm. The user has no means of controlling how and when this barrier is imposed and also has no means of avoiding the resource starvation associated with it.

A possible remedy for this problem is to allow for the algorithms to be executed asynchronously. While this does not remove the implicit barrier at the end of the execution of any of the algorithms, it allows to reduce the resource starvation by allowing to perform other, unrelated tasks while the join-operation (and the associated tapering of parallel work) is being executed.

This paper proposes to enable such an asynchronous execution of all algorithms as defined by N4507 by introducing special execution policies which essentially launch the execution of the algorithm on a new thread of execution while the algorithm invocation itself now returns a std::future representing the result of its execution.

Returning a Future object from the algorithm has the additional advantage of being able to integrate the parallel algorithms with other asynchronous codes which also rely on representing their results through std::future. This is especially important in light of the proposed additions to std::future as described by the Concurrency TS (4501) [3].

The proposed extensions have been implemented in HPX [4] which has an implementation of N4507. They are in use in production codes for some time.

#### 2. Summary of the Proposed Functionality

An asynchronous execution policy is an object which fulfills the concept of an execution policy as defined in N4507. Additionally it instructs a parallel algorithm to launch its execution on a new thread of execution and changes the algorithm to return a Future object representing the result of the execution of the original algorithm.

We propose that every one of the already specified execution policies (seq, par, and par\_vec) has a corresponding asynchronous execution policy which is generated by seq(task), par(task), and par\_vec(task).

With those extensions, the following use cases of the parallel algorithms library are possible:

```
using namespace std::experimental::parallel::v1;
std::vector<int> data = { ... };
// legacy standard sequential sort
std::sort(data.begin(), data.end());
// explicitly sequential sort
sort(seq, data.begin(), data.end());
```

```
// permitting parallel execution
sort(par, data.begin(), data.end());
// permitting vectorized execution as well
sort(par_vec, data.begin(), data.end());
// NEW: asynchronous, sequential execution
std::future<void> f1 = sort(seq(task), data.begin(), data.end());
// ... perform other work
             // synchronize with the asynchronous sequential sort()
f1.get();
// NEW: asynchronous execution, allow for parallelization of the algorithm
std::future<void> f2 = sort(par(task), data.begin(), data.end());
// ... perform other work
f2.get();
             // synchronize with the asynchronous parallel sort()
// NEW: asynchronous execution, allow for parallelization and vectorization
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       of the algorithm
std::future<void> f3 = sort(par_vec(task), data.begin(), data.end());
// ... perform other work
f3.get();
             // synchronize with the asynchronous parallel vectorized
sort()
```

#### 3. Specification

```
Header <experimental/execution_policy> synopsis
```

The following definitions are being proposed to be added to this header file.

```
namespace std {
namespace experimental {
namespace parallel {
inline namespace v1 {
// 3.1, Task modifier tag type for execution policies
class task_execution_policy_tag {};
// 3.2, Task modifier instance for execution policies
constexpr task_execution_policy_tag task {};
// 3.3, Asynchronous execution policy type trait
template<class T> struct is_task_execution_policy;
template<class T> constexpr
bool is_task_execution_policy_v = is_task_execution_policy<T>::value;
```

```
// 3.4, Sequential asynchronous execution policy
class sequential_task_execution_policy;
// 3.5, Parallel asynchronous execution policy
class parallel_task_execution_policy;
// 3.6, Parallel+Vector asynchronous execution policy
class parallel vector task execution policy;
// 3.7, Generator function operators added to existing execution policies
sequential task execution policy
sequential_execution_policy::operator()(
    task_execution_policy_tag) const;
parallel_task_execution_policy
parallel_execution_policy::operator()(
    task_execution_policy_tag) const;
parallel_vector_task_execution_policy
parallel_vector_execution_policy::operator()(
    task_execution_policy_tag) const;
}
}
}
}
```

## 3.1 Task modifier tag type for execution policies

class task\_execution\_policy\_tag { unspecified };

<sup>1</sup>task\_execution\_policy\_tag is a unique type used to generate an asynchronous execution policy from its non-asynchronous counterpart

## 3.2 Task modifier instance for execution policies

```
constexpr task_execution_policy_tag task{};
```

<sup>1</sup>The header <experimental/execution\_policy> declares a global object for the task modifier tag type.

## 3.3 Asynchronous execution policy type trait

template<class T> struct is\_task\_execution\_policy { see beLow };

<sup>1</sup>is\_task\_execution\_policy can be used to detect parallel execution policies for the purpose of excluding function signatures from otherwise ambiguous overload resolution participation.

<sup>2</sup> is\_task\_execution\_policy<T> shall be a UnaryTypeTrait with a BaseCharacteristic of true\_type if T is the type of a standard or implementation-defined asynchronous execution policy, otherwise false\_type.

[Note: This provision reserves the privilege of creating non-standard asynchronous execution policies to the library implementation. —end note ]

<sup>3</sup>The behavior of a program that adds specializations for is\_task\_execution\_policy is undefined.

## 3.4 Sequential asynchronous execution policy

```
class sequential_task_execution_policy { unspecified };
```

<sup>1</sup>The class sequential\_task\_execution\_policy is an asynchronous execution policy type used as a unique type to disambiguate asynchronous parallel algorithm overloading and require that a parallel algorithm's execution may not be parallelized, that the algorithm should be executed asynchronously, and that the return type of the algorithm should be a future<T>, where T is the type as returned by the non-asynchronous version of the same algorithm.

## 3.5 Parallel asynchronous execution policy

#### class parallel\_task\_execution\_policy { unspecified };

<sup>1</sup>The class parallel\_task\_execution\_policy is an asynchronous execution policy type used as a unique type to disambiguate asynchronous parallel algorithm overloading and indicate that a parallel algorithm's execution may be parallelized, that the algorithm should be executed asynchronously, and that the return type of the algorithm should be a future<T>, where T is the type as returned by the non-asynchronous version of the same algorithm.

## 3.6 Parallel+Vector asynchronous execution policy

## class parallel\_vector\_task\_execution\_policy { unspecified };

<sup>1</sup>The class parallel\_vector\_task\_execution\_policy is an asynchronous execution policy type used as a unique type to disambiguate asynchronous parallel algorithm overloading and indicate that a parallel algorithm's execution may be vectorized and parallelized, that the algorithm should be executed asynchronously, and that the return type of the algorithm should be a future<T>, where T is the type as returned by the non-asynchronous version of the same algorithm.

# 3.7 Generator function operators added to existing execution policies

Every of the non-asynchronous execution policies as defined by N4507 has an added function operator used to generate a corresponding asynchronous execution policy.

```
sequential_task_execution_policy
sequential_execution_policy::
    operator()(task_execution_policy_tag) const;
parallel_task_execution_policy
parallel_execution_policy::
    operator()(task_execution_policy_tag) const;
parallel_vector_task_execution_policy;:
    operator()(task_execution_policy;:
    operator()(task_execution_policy;:
    operator()(task_execution_policy_tag) const;
```

## **4 Exception Handling**

All behavior regarding generating exceptions is unchanged from the parallelism TS except that none of the algorithms shall directly throw any of the generated exceptions if invoked with an asynchronous execution policy but deliver the exception through the returned future object.

## **5 Example**

Given a synchronous algorithm gather:

```
template <typename BiIter, typename Pred>
pair<BiIter, BiIter>
   gather(BiIter f, BiIter l, BiIter p, Pred pred)
{
   BiIter it1 = stable_partition(f, p, not1(pred));
   BiIter it2 = stable_partition(p, l, pred);
   return make_pair(it1, it2);
}
```

the following example demonstrates how the proposed features can be used to compose more complex asynchronous algorithms. The gather algorithm is meant to collect all elements in a given range [f, 1) at the given position p for which a given boolean predicate pred is true. The implementation above achieves that by invoking stable\_partition twice, once for all elements in the range [f, p) while using the inverted predicate, and once for the elements in the range [p, 1) using the predicate as is. The algorithm gather returns a pair of iterators marking the range of the newly inserted elements. The asynchronous version of the same algorithm (here gather\_async), is called using the same arguments, it however returns a future to the pair of result iterators.

```
template <typename BiIter, typename Pred>
future<pair<BiIter, BiIter>>
   gather_async(BiIter f, BiIter l, BiIter p, Pred pred)
{
   future<BiIter> f1 = stable_partition(par(task), f, p, not1(pred));
   future<BiIter> f2 = stable_partition(par(task), p, l, pred);
   return when_all(f1, f2).then(
      [](tuple<future<BiIter>, future<BiIter>> p)
      { return make_pair(get<0>(p).get(), get<1>(p).get()); }
      );
}
```

}

The benefit of calling the asynchronous versions of the stable\_partition algorithms is twofold: a) both sub-regions can be handled concurrently, and b) the overall algorithm can be made asynchronous. The only caveat of this implementation is the slightly complicated code necessary to convert the pair of futures into a future of pairs using when\_all().then() (as proposed by the concurrency TS).

However, by using co\_await (see P0057R3, [5]), this can be further simplified:

```
template <typename BiIter, typename Pred>
future<pair<BiIter, BiIter>>
   gather_async(BiIter f, BiIter l, BiIter p, Pred pred)
{
   future<BiIter> f1 = stable_partition(par(task), f, p, not1(pred));
   future<BiIter> f2 = stable_partition(par(task), p, l, pred);
   return make_pair(co_await f1, co_await f2);
}
```

Please note, that the version using co\_await is 100% semantically equivalent to the asynchronous version using when\_all().then().

## References

[1] N4507, Technical Specification for C++ Extensions for Parallelism, (ed) J. Hoberock, <u>http://www.open-std.org/jtc1/sc22/wg21/docs/papers/2015/n4507.pdf</u>

[2] P00234R0: Towards Massive Parallelism (aka Heterogeneous Devices/Accelerator/GPGPU) support in C++ with HPX, Michael Wong, Hartmut Kaiser, Thomas Heller, <u>http://www.open-std.org/jtc1/sc22/wg21/docs/papers/2016/p0057r3.pdf</u>

[3] N4501, Working Draft, Technical Specification for C++ Extensions for Concurrency, (ed) A. Laksberg, <u>http://www.open-std.org/jtc1/sc22/wg21/docs/papers/2015/n4501.html</u>

[4] HPX: A general purpose C++ runtime system for parallel and distributed applications of any scale, <u>https://github.com/STEIIAR-GROUP/hpx</u>.

[5] P0057R3: Wording for Coroutines, Gor Nishanov, Jens Maurer, Richard Smith, Daveed Vandevoorde, <u>http://www.open-std.org/jtc1/sc22/wg21/docs/papers/2016/p0057r3.pdf</u>