1 ISO/IEC JTC 1/SC 22/WG 23 N 0327

2 Additional meeting #17 markup of Proposed vulnerability descriptions YUK and SUK

5	Date Contributed by Original file name Notes	March 25, 2011 Secretary Replaces N0324	
4 5 6	I wrote up two vul	nerabilites instead of one.	
7 8	The first one deals	with the suppression of runtime checks (as I was tasked to do).	
9 10 11		als with the de-facto suppression of compile-time checks and with operations that the language might provide.	
12 13	I simply could not they are of the same	(-	
14	Suppression of	Language-Defined Run-Time Checking (YUK)	Comment [JWM1]: Mention that some languages disable the checking by default.
15 16 17	Description of app	lication vulnerability	
18 19 20 21 22	Some languages include the provision for runtime checking to prevent vulnerabilities to arise. Canonical examples are bounds or length checks on array operations or null-value checks upon dereferencing pointers or references. In most cases, the reaction to a failed check is the raising of a language-defined exception.		
23 24 25 26 27 28	of exceptions, lang	ing requires execution time and as some project guidelines exclude the use guages may define a way to optionally suppress such checking for regions he entire program. Analogously, compiler options may be used to achieve	
29	Cross reference		
30 31 32 33			
34 35	Mechanism of Fail	lure	
36 37 38 39 40 41	suppressed checks subsections.Vulner	that should have been prevented by the checks re emerge whenever the would have failed. For their description, see the respective rabilities that could have been prevented by the run-time checks are ng in memory corruption, propagation of incorrect values or unintended	
42 43	Applicable langua	ge characteristics	
44		~	

45 46 47	This vulnerability description is intended to be applicable to languages with the following characteristics:					
48 49	• Languages that define runtime checks to prevent certain vulnerabilities and					
50 51	 Languages that allow the above checks to be suppressed, 	Formatted: List Paragraph, No bullets or				
52 53	 Languages or compilers that suppress checking by default, or 	numbering Formatted: Bulleted + Level: 1 + Aligned a				
54 55	• Languages, whose compilers or interpreters provide options to omit the above checks	0.25" + Tab after: 0.5" + Indent at: 0.5"				
56 57 58	Avoiding the vulnerability					
59 60 61	Software developers can avoid the vulnerability or mitigate its ill effects in the following ways:					
62 63 64	 <u>Do not suppress checks or restrict such suppression to the most performance critical sections of the code.</u> Do not suppress checks at all or restrict the suppression of checks to regions of the code that have been proved to be performance-critical. 					
65 66 67	 If the default behaviour of the compiler or the language is to suppress checks, then enable them. 	Formatted: Indent: Left: 0.5", No bullets on numbering				
68 69 70	• Where checks are suppressed, verify that the suppressed checks could not have failed.					
70 71 72	• Clearly identify code sections where checks are suppressed.					
73 74 75 76 77	• Do not assume that checks in code verified to satisfy all checks could not fail nevertheless due to hardware faults.					
78	Provision of Inherently Unsafe Operations (SUK)					
79 80 81	Description of application vulnerability					
82 83 84	Languages define semantic rules to be obeyed by legal programs. Compilers enforce these rules and reject violating programs.					
85 86 87 88	A canonical example are the rules of type checking, intended among other reasons to prevent semantically incorrect assignments, such as characters to pointers, meter to feet, euro to dollar, real numbers to booleans, or complex numbers to two-dimensional coordinates.					
89 90 91 92	Yet, occasionally there arises a need to step outside the rules of the type model to achieve needed functionally. A typical such-situation is the casting of memory as part of the implementation of a heap allocator to the type of object for which the memory is allocated. A type-safe assignment is impossible for this functionality. Thus, a capability for unchecked					

93 94 95	"type casting" between arbitrary types to interpret the bits in a different fashion is a necessary but inherently unsafe operation, without which the type-safe allocator cannot be programmed.			
96 97	Another example is the provision of operations known to be inherently unsafe, such as the deallocation of heap memory without prevention of dangling references.			
98				
99	A third example is any interfacing with another language, since the checks ensuring type-			
100	safeness rarely extend across language boundaries.			
101				
102	These inherently unsafe operations constitute a vulnerability, since they can (and will) be used			
103 104	by programmers in situations where their use is neither necessary nor appropriate. As the knowledge of the programmer about implementation details may be incomplete or incorrect,			
104	unintended execution semantics may result.			
105	difficited execution semantics may result.			
107	The vulnerability is eminently exploitable to violate program security.			
108				
109				
110	Cross reference			
111				
112				
113				
114 115	Mechanism of Failure			
115	The use of inherently unsafe operations or the sSuppression of checkings of the use of			
117	inherently unsafe operations circumvents the ehecks-features that are normally applied to			
118	ensure safe execution. Control flow, data values, and memory accesses can be corrupted as a			
119	consequence. See the respective vulnerabilities resulting from such corruption.			
120				
121				
122	Applicable language characteristics			
123				
124	This vulnerability description is intended to be applicable to languages with the following characteristics:			
125 126	characteristics.			
120	• Languages that allow compile-time checks for the prevention of vulnerabilities to be			
127	suppressed by compiler or interpreter options or by language constructs, or			
129	suppressed by compiler of interpreter options of by nungative constitueis, of			
130	• Languages that provide inherently unsafe operations			
131				
132				
133	Avoiding the vulnerability			
134				
135	Software developers can avoid the vulnerability or mitigate its ill effects in the following			
136	ways:			
137	• Destrict the summaries of commile time the state to sub-set the summaries i			
138 139	• Restrict the suppression of compile-time checks to where the suppression is functionally essential.			
139				
140	• Use inherently unsafe operations only when they are functionally essential.			
142	ese interently unsure operations only when they are functionally essential.			

	-	
143	•	Clearly identify program code that suppresses checks or uses unsafe operations. This
144		permits the focusing of review effort to examine whether the function could be
145		performed in a safer manner.