Abstract

This Omnibus paper proposes a series of small improvements to the lexing of C++ and the forming of string handling, notably to clarify the behavior of Unicode encoded sources files, and to reduce implementation divergence. This paper intends to supersede N3463 [1] and P1854R0 [3] and proposes a resolution to several core issues.

While some of these changes are unrelated, the intend of this paper’s authors and SG-16 is to rewrite the description of lexing using more accurate terminology, a wording will, therefore, be provided incorporating all the desired design changes.

Updating both the design an the terminology lets us acknowledge and handle the subtilities of text which may not have been fully understood in a pre-Unicode world. The overarching goals are to reduce the divergence in implementations, the complexity of the mental model, and most importantly to make sure that the semantic of text elements is conserved through the different phases of compilation.

Proposals

All the proposed changes relate to phases 1-6 of translations. In particular, this proposal has no library impact.

Proposal 1: Mandating support for UTF-8 encoded source files in phase 1

The set of source file character sets is implementation-defined, which makes writing portable C++ code impossible. We proposed to mandate that C++ compilers must accept UTF-8 as an input format. Both to increase portability and to ensure that Unicode related features (ex P1949R3 [2] can be used widely. This would also allow us to better specify how Unicode encoded files are handled.

How the source file encoding is detected, and which other input formats are accepted would remain implementation-defined.
Most C++ compilers (GCC, EDG, MSVC, Clang) support utf8 as one of their input format - Clang only supports utf8.

Unlike N3463 [1], we do not intend to mandate or say anything about BOMs, following Unicode recommendations. BOMs should be ignored (but using BOMs as a means to detect UTF-8 files is a valid implementation strategy, which is notably used by MSVC).

We do **not** propose to make UTF-8 source file a mandated default, nor the only supported format. Just that there must be some implementation-defined mechanism (such as a compiler flag) that would tell the compiler to read the file as utf-8.

**Proposal 2: What is a whitespace or a new-line?**

We propose to specify that the following code point sequences are line-terminators (after phase 1):

- LF: Line Feed, U+000A
- VT: Vertical Tab, U+000B
- FF: Form Feed, U+000C
- CR: Carriage Return, U+000D
- CR+LF: CR (U+000D) followed by LF (U+000A)
- NEL: Next Line, U+0085
- LS: Line Separator, U+2028
- PS: Paragraph Separator, U+2029

Line terminators and the following characters constitute whitespaces

- U+0009 HORIZONTAL TAB
- U+0020 SPACE
- U+200E LEFT-TO-RIGHT MARK
- U+200F RIGHT-TO-LEFT MARK

These correspond to characters with the Pattern_Whitespace Unicode property. The line terminator subset is derived from UAX14 - UNICODE LINE BREAKING ALGORITHM.

We intend to fix CWG1655 [7] following this clarification.

**Proposal 3: Preserve Normalization forms**

We propose to specify that Unicode encodes files are not normalized in phase 1 or phase 5, as to preserve the integrity of string literals when both the source and the literal associated character set are the Unicode character set. Instead, the exact sequence of code points of these literals is preserved. In effect, this does not change the existing behavior of tested implementations in phase 1 (and phase 5 is already specified on a per code point basis).
Proposal 4: Making trailing whitespaces non-significant

There is a divergence of implementation in how compilers handle spaces between a backslash and the end of a line.

```c
int main() {
    int i = 1
    // \ + 42
    ;
    return i;
}
```

EDG (tested with icc front-end) GCC and Clang will trimm the whitespaces after the backslash - and return 1 - MSVC will not and return 43. Both strategies are valid as part of phase 1 "implementation-defined mapping".

To avoid this surprising implementation divergence we proposed that an implementation must trim all trailing whitespaces before handling \ slicing. This is reinforced by the fact thatides and tools may discard such whitespaces. The Google-style guidelines forbid trailing whitespaces.

An additional or alternative approach is to deprecate \ that are not part of a preprocessor directive. We are not proposing this at this time.

Unlike other proposals in this paper, this maybe is a silent breaking change for code that is only compiled with MSVC. A quick analysis of the VCPKG package didn't find any trailing whitespaces after backslashes

We have not been able to measure the impact of this proposed change in the MSVC ecosystem. Other compilers, and all code supported by these compilers would be unaffected.

Proposal 5: Restricting multi-characters literals to members of the Basic Latin Block

```c
int i = 'é'; can be equivalent to either int i = '\u00e9'; or int i = 'e\u0301'; depending on source encoding and normalization. There is also a lot of divergence of implementations in how literals are interpreted.
```

<table>
<thead>
<tr>
<th>Clang</th>
<th>GCC UTF-8</th>
<th>MSVC UTF-8</th>
<th>GCC Latin1</th>
<th>MSVC latin 1</th>
</tr>
</thead>
<tbody>
<tr>
<td>'é\u0301'; ill-formed</td>
<td>int(0x65CC81) + Warning</td>
<td>int(0x65cc81)</td>
<td>ill-formed</td>
<td>int (0x653f)</td>
</tr>
<tr>
<td>'\u00e9'; ill-formed</td>
<td>int(0xC3A9) + Warning</td>
<td>int(0xC3A9)</td>
<td>0xFFFFFFFFFFFFFFE9</td>
<td>int(0x09)</td>
</tr>
</tbody>
</table>

Note the presence of two code points in the first line.

We propose to limit multi-character literals to a sequence of code points from the Unicode Basic Latin Block (~ASCII) to limit the current confusion.

(We do not propose to deprecate multi-character literals).

With the proposed change:
Proposal 6: Making wide characters literals containing multiple or unrepresentable c-char ill-formed

The following constructs are well-formed but have widely different interpretations depending on implementations

```c
wchar_t a = L 'é';
wchar_t b = L 'ab';
wchar_t c = L é;
```

- the size of wchar_t being implementation defined, L’é’ is correctly interpreted on Unix platforms where that size is 32 bits, but truncated by MSVC and other compatible windows compilers where wchar_t is 16 bits. MSVC first converts to UTF-16, and truncate to the first code unit which results in an invalid lone high surrogate \xd83d.
- L’ab’ is equivalent to L’a’ on msvc and L ‘b’ on GCC and Clang. All implementation emit a warning under different warning levels
- L’é’ can be either 1 or 2 c-char depending on the source normalization: L’\u00e9’ behaves expectedly on all platforms, while L’e\u0301’ will be e in MSVC and U+0301 in GCC AND clang.

As such, we propose to make wide characters literals with multiple-chars or char which are not representable in the execution character set ill-formed.

Proposal 7: Making conversion of character and string literals to execution and wide execution encoding ill-formed for unrepresentable c-char

Implementations diverge on how they handle unrepresentable code points when conversion to execution encodings. GCC and Clang make the conversion ill-formed while MSVC usually replaces unrepresentable characters by a single question mark ?. Strings are text which carries intent and meaning; An implementation should not be able to alter that meaning.

We proposed to make such conversion ill-formed rather than implementation-defined.

After discussions, in SG-16, we do not propose to allow implementations to consider multi-code points graphemes clusters when doing that conversion. For example considering "é\u0301",
U+301 does not have a representation in latin 1, but the abstract character é does (U+00e9 maps to 0x00E9 in ISO/IEC 8859-1).

However, it does not seems possible to guarantee that an implementation knows about all such mapping, which would lead to further implementation divergence and unnecessary burden on compilers. We, therefore, propose to be explicit about the conversion being done on each code point independently as is currently the case.

Proposal 8: Enforcing the formation of universal escape sequences in phase 2 and 4

EDG(icc), GCC, MSVC and Clang form universal character names from the following codes:

```
'\u00e9';
//---
#define CONCAT(x,y) x##y
CONCAT(\, U0001F431);
```

However, these behaviors are currently UB within the standard. As such, we propose to make both these behaviors well-defined to follow existing practices.

Proposal 9: Reaffirming Unicode as the character set of the internal representation

The standard already specifies that characters outside of the basic source character set are converted to UCNs whose values are isomorphic to Unicode. We want to make it clear that characters that do not have representation in Unicode are ill-formed. This includes some characters in some Big5 encodings and exotic languages such as Klingon.

In particular, all characters in EBCDIC ¹, GB 18030 have a unique mapping in Unicode. The intent is to avoid the use of unassigned code points or the Private Use Area by the implementers, as well as to preserve semantic meaning in phase 1. The preservation of semantic meaning would also make invalid utf-8 sequences ill-formed in phase 1, and other decoding errors (as there exist no mapping for such invalid sequence). (Note that octal/hexadecimal escape sequences can be used in string literal to form arbitrary binary data)

Notably, it is important to consider that the current specification limits the implementation-defined mapping to universal character names to valid (0-10FFF) code points, and any such valid code point can appear in the source before phase 1. It is therefore not possible for an implementation to uniquely map unrepresentable characters to a valid code point.

This proposal has no bearing on the actual internal representation strategy of already conforming implementations. Notably, mandating internal as-if Unicode representation doesn't preclude bytewise preservation of narrow and wide string literals when the execution encoding is identical to the source encoding, as long as there exists a Unicode representation, as this is

¹For EBCDIC, the mapping of control characters is specified in Unicode Technical Report 16 - UTF-EBCDIC. This mapping is not semantic-preserving, to the extent control characters have semantics.
otherwise non-observable: It is an important implementation strategy for encoding which are not roundtrip-able through Unicode such as Shift JIS to preserve the byte content of string literals when both source and execution encodings are identical. Such preservation is otherwise non-observable and doesn't need to be mandated, but it needs not to be precluded. **We only seek to mandate representability in Unicode.** Such representability is notably necessary for the concatenation of literals with different encoding which may happen in phase 6 ("foo" u"bar" // u"foobar"). Again, this constraint exists in practice as many implementations use Unicode internally and do not use the "implementation-defined mapping" leeway to nefarious ends.

Specifically:

- EBCDIC encodings would be converted to Unicode according to UTF-EBCDIC / IBM CDRA - as such, EBCDIC specific control characters are mapped to C1 control characters. C1 controls characters have no meaning on their own and are designed to be interpreted in an application-specific manner.
- GB 18030 maps to Unicode, but a handful of code points maps to the Private Use Area
- Big 5: Most abstract characters map to Unicode, with the rare exception of some spelling of some people or place names. Notably, Unicode prescribes a mapping for Windows implementation (code page 951/950)
- All other encodings have a complete, semantic preserving mapping to assigned Unicode code points.

That list does not intend to be prescriptive, but to show that the C++ standard doesn't need and shouldn't try to handle characters not representable of Unicode. Furthermore, the author hopes that the wording can convert *universal-character-names* to Unicode code points as soon as they are encountered or formed (which is exclusively a matter of wording, that would neither affect behavior nor prescribe an internal representation).

Following this clarification we hope to fix CWG1332 [6]

**Proposal 10: Make L in _Pragma ill-formed**

_Pragma(L""") is equivalent to _Pragma("""").

We propose to remove the _Pragma(L""") syntax as both strings are interpreted as sequences of Unicode code points and never as a wide execution literal. C++ handling of text is confusing enough not to add meaningless characters. This would resolve CWG897 [5]. Note that there is a divergence of implementation between C++ and C where C discard all prefixes and C++ only discards L.

Out of the 90 millions lines of code of the 1300+ open source projects available on vcpkg, a single use of that feature was found within clang's lexer test suite, for a total of 2000 uses of _Pragma. Similarly, the only uses of _Pragma (u8""""), _Pragma (u""""), _Pragma (U""""), etc were found in Clang's test suite.
Proposal 11: Make character literals in preprocessor conditional behave like they do in C++ expression

```c
#if 'A' == '\x65'
#else
if ('A' == 0x65){}
```

Both conditions are not guaranteed to yield a similar result, as the value of character literals in preprocessor conditional is not required to be identical to that of character literals in expressions.

However, a survey of the 1300+ open sources projects available on vcpkg shows that the primary use case for these macros is exactly to detect the execution encoding at compile time and all compilers available on compiler explorer treat these literals as if they were in the execution encoding.

Notably, a few libraries use that pattern to detect EBCDIC or ASCII execution encoding. Of the 50 usages of the pattern, all but one were in C libraries.

While we think there should be a better way to detect encodings in C++ [4], there is no reason to deprecate that feature.

Instead, we recommend adopting the standard practice and user expectation of converting these literals to the execution encoding before evaluating them.

This also removes yet another theoretical encoding, which further simplifies the mental model.
Proposal 12: Phase 6 needs fixing

String literal concatenation happen after each literal has been converted to its associated encoding in phase 5, and as such the standard is actively encouraging mojibake in the following scenarios:

```
L"" ""  
"" u8"" ""  
u"" ""  
U"" ""  
"" u""  
"" L""  
"" U""  
"" u8""
```

While string conversion is definitively useful, the current behavior is not.

Indeed, each fragment may be in a different encoding after concatenation: utf-8 data concatenated to shift-jis data, for example.

We propose 2 possible resolutions:

- The program is ill-formed unless each string-literal part of a concatenation has the same prefix

- The encoding of the first string-literal determines the encoding of the string-literal resulting from the concatenation and the program is ill-formed if any but the first string-literal literal has an encoding prefix:

```
"" ""  // OK, const char*
u8"" ""  // OK, equivalent to u8"" u8"", const char8_t*
u"" ""  // OK, equivalent to u"" u", const char16_t*
U"" ""  // OK, equivalent to U"" U", const char32_t*
L"" ""  // OK, equivalent to L"" L", const wchar_t*
"" L""  // ill-formed
L"" u""  // ill-formed
```

In which case, each string-literal should be interpreted as having the same prefix in phase 5, so they are converted to the same encoding prior to concatenation.

There is implementation divergence in the handling of concatenation:

MSVC converts each string to its associated execution encoding, as specified by phase 6, while GCC and clang consider the following snippets equivalent (where x is any valid encoding-prefix)

```
X"" ""  
"" X""  
X"" X"
```
Annex: Schematization of text encodings handling during compilation

Current model

The graph below is a simplification of the different encodings that appear during compilation. Each rectangle represents a possibly different encoding. The red arrows represent operations that may alter the semantic of text elements.
Proposed model

In the proposed model, all conversions are either semantic preserving\(^2\) or ill-formed. Preprocessor Conditionals use the narrow literal encoding (execution literal encodings). Outside of raw literals, *universal-character-names* are converted to codepoints as they are formed, the wording is specified in term of the Unicode character set.

\(^2\)The value of multi-character literals remains implementation-defined
Acknowledgments

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References


