P1673R2: A free function linear algebra interface based on the BLAS

Authors

- Mark Hoemmen (mhoemmen@sandia.gov) (Sandia National Laboratories)
- David Hollman (dshollm@sandia.gov) (Sandia National Laboratories)
- Christian Trott (ctrott@sandia.gov) (Sandia National Laboratories)
- Daniel Sunderland (dsunder@sandia.gov) (Sandia National Laboratories)
- Nevin Liber (nliber@anl.gov) (Argonne National Laboratory)
- Siva Rajamanickam (srajama@sandia.gov) (Sandia National Laboratories)
- Li-Ta Lo (ollie@lanl.gov) (Los Alamos National Laboratory)
- Damien Lebrun-Grandie (lebrungrandt@ornl.gov) (Oak Ridge National Laboratories)
- Graham Lopez (lopezmg@ornl.gov) (Oak Ridge National Laboratories)
- Peter Caday (peter.caday@intel.com) (Intel)
- Sarah Knepper (sarah.knepper@intel.com) (Intel)
- Piotr Luszczek (luszczech@icl.utk.edu) (University of Tennessee)
- Timothy Costa (tcosta@nvidia.com) (NVIDIA)

Contributors

- Chip Freitag (chip.freitag@amd.com) (AMD)
- Bryce Leibach (bleibach@nvidia.com) (NVIDIA)
- Srinath Vadlamani (Srinath.Vadlamani@arm.com) (ARM)
- Rene Vanoostroom (Rene.Vanoostroom@amd.com) (AMD)

Date: 2020-01-12

Revision history

- Revision 0 (pre-Cologne) submitted 2019-06-17
  - Received feedback in Cologne from SG6, LEWGI, and (???).
- Revision 1 (pre-Belfast) to be submitted 2019-10-07
  - Account for Cologne 2019 feedback
    - Make interface more consistent with existing Standard algorithms
      - Change dot, dotc, vector_norm2, and vector_abs_sum to imitate reduce, so that they return their result, instead of taking an output parameter. Users may set the result type via optional init parameter.
    - Minor changes to "expression template" classes, based on implementation experience
  - Briefly address LEWGI request of exploring concepts for input arguments.
Lazy ranges style API was NOT explored.

• Revision 2 (pre-Cologne) to be submitted 2020-01-13
  • Add "Future work" section.
  • Remove "Options and votes" section (which were addressed in SG6, SG14, and LEWGI).
  • Remove basic_mdarray overloads.
  • Remove batched linear algebra operations.
  • Remove over- and underflow requirement for vector_norm2.
  • Mandate any extent compatibility checks that can be done at compile time.
  • Add missing functions {symmetric,hermitian}_matrix_rank_k_update and triangular_matrix_{left,right}_product.
  • Remove packed_view function.
  • Fix wording for {conjugate,transpose,conjugate_transpose}_view, so that implementations may optimize the return type. Make sure that transpose_view of a layout_blas_packed matrix returns a layout_blas_packed matrix with opposite Triangle and StorageOrder.
  • Remove second template parameter T from accessor_conjugate.
  • Make scaled_scalar and conjugated_scalar exposition only.
  • Add in-place overloads of triangular_matrix_matrix_{left,right}_solve, triangular_matrix_{left,right}_product, and triangular_matrix_vector_{product,solve}.
  • Add alpha overloads to {symmetric,hermitian}_matrix_rank_{1,k}_update.
  • Add Cholesky factorization and solve examples.

Purpose of this paper

This paper proposes a C++ Standard Library dense linear algebra interface based on the dense Basic Linear Algebra Subroutines (BLAS). This corresponds to a subset of the BLAS Standard. Our proposal implements the following classes of algorithms on matrices and vectors:

• Elementwise vector sums
• Multiplying all elements of a vector or matrix by a scalar
• 2-norms and 1-norms of vectors
• Vector-vector, matrix-vector, and matrix-matrix products (contractions)
• Low-rank updates of a matrix
• Triangular solves with one or more "right-hand side" vectors
• Generating and applying plane (Givens) rotations
Our algorithms work with most of the matrix storage formats that the BLAS Standard supports:

- "General" dense matrices, in column-major or row-major format
- Symmetric or Hermitian (for complex numbers only) dense matrices, stored either as general dense matrices, or in a packed format
- Dense triangular matrices, stored either as general dense matrices or in a packed format

Our proposal also has the following distinctive characteristics:

- It uses free functions, not arithmetic operator overloading.
- The interface is designed in the spirit of the C++ Standard Library's algorithms.
- It uses the multidimensional array data structures `basic_mdspan` ([P0009R9](http://example.com)) to represent matrices and vectors. In the future, it could support other proposals' matrix and vector data structures.
- The interface permits optimizations for matrices and vectors with small compile-time dimensions; the standard BLAS interface does not.
- Each of our proposed operations supports all element types for which that operation makes sense, unlike the BLAS, which only supports four element types.
- Our operations permit "mixed-precision" computation with matrices and vectors that have different element types. This subsumes most functionality of the Mixed-Precision BLAS specification (Chapter 4 of the BLAS Standard).
- Like the C++ Standard Library's algorithms, our operations take an optional execution policy argument. This is a hook to support parallel execution and hierarchical parallelism (through the proposed executor extensions to execution policies, see [P1019R2](http://example.com)).
- Unlike the BLAS, our proposal can be expanded to support "batched" operations (see [P1417R0](http://example.com)) with almost no interface differences. This will support machine learning and other applications that need to do many small matrix or vector operations at once.

Interoperable with other linear algebra proposals

We believe this proposal is complementary to [P1385](http://example.com), a proposal for a C++ Standard linear algebra library that introduces matrix and vector classes and overloaded arithmetic operators. In fact, we think that our proposal would make a natural foundation for a library like what P1385 proposes. However, a free function interface -- which clearly separates algorithms from data structures -- more naturally allows for a richer set of operations such as what the BLAS provides. A natural extension of the present proposal would include accepting P1385's matrix and vector objects as input for the algorithms proposed here.

Why include dense linear algebra in the C++ Standard Library?

1. C++ applications in "important application areas" (see [P0939R0](http://example.com)) have depended on linear algebra for a long time.

2. Linear algebra is like sort: obvious algorithms are slow, and the fastest implementations call for hardware-specific tuning.
3. Dense linear algebra is core functionality for most of linear algebra, and can also serve as a building block for tensor operations.

4. The C++ Standard Library includes plenty of "mathematical functions." Linear algebra operations like matrix-matrix multiply are at least as broadly useful.

5. The set of linear algebra operations in this proposal are derived from a well-established, standard set of algorithms that has changed very little in decades. It is one of the strongest possible examples of standardizing existing practice that anyone could bring to C++.

6. This proposal follows in the footsteps of many recent successful incorporations of existing standards into C++, including the UTC and TAI standard definitions from the International Telecommunications Union, the time zone database standard from the International Assigned Numbers Authority, and the ongoing effort to integrate the ISO unicode standard.

Linear algebra has had wide use in C++ applications for nearly three decades (see P1417R0 for a historical survey). For much of that time, many third-party C++ libraries for linear algebra have been available. Many different subject areas depend on linear algebra, including machine learning, data mining, web search, statistics, computer graphics, medical imaging, geolocation and mapping, engineering, and physics-based simulations.

"Directions for ISO C++" (P0939R0) offers the following in support of adding linear algebra to the C++ Standard Library:

- P0939R0 calls out "Support for demanding applications in important application areas, such as medical, finance, automotive, and games (e.g., key libraries...)" as an area of general concern that "we should not ignore." All of these areas depend on linear algebra.

- "Is my proposal essential for some important application domain?" Many large and small private companies, science and engineering laboratories, and academics in many different fields all depend on linear algebra.

- "We need better support for modern hardware": Modern hardware spends many of its cycles in linear algebra. For decades, hardware vendors, some represented at WG21 meetings, have provided and continue to provide features specifically to accelerate linear algebra operations. Some of them even implement specific linear algebra operations directly in hardware. Examples include NVIDIA’s Tensor Cores and Cerebras’ Wafer Scale Engine. Several large computer system vendors offer optimized linear algebra libraries based on or closely resembling the BLAS; these include AMD’s BLIS, ARM’s Performance Libraries, Cray’s LibSci, Intel’s Math Kernel Library (MKL), IBM’s Engineering and Scientific Subroutine Library (ESSL), and NVIDIA’s cuBLAS.

Obvious algorithms for some linear algebra operations like dense matrix-matrix multiply are asymptotically slower than less-obvious algorithms. (Please refer to a survey one of us coauthored, "Communication lower bounds and optimal algorithms for numerical linear algebra.") Furthermore, writing the fastest dense matrix-matrix multiply depends on details of a specific computer architecture. This makes such operations comparable to sort in the C++ Standard Library: worth standardizing, so that Standard Library implementers can get them
right and hardware vendors can optimize them. In fact, almost all C++ linear algebra libraries end up calling non-C++ implementations of these algorithms, especially the implementations in optimized BLAS libraries (see below). In this respect, linear algebra is also analogous to standard library features like random_device: often implemented directly in assembly or even with special hardware, and thus an essential component of allowing no room for another language "below" C++ (see notes on this philosophy in P0939R0 and Stroustrup's seminal work "The Design and Evolution of C++").

Dense linear algebra is the core component of most algorithms and applications that use linear algebra, and the component that is most widely shared over different application areas. For example, tensor computations end up spending most of their time in optimized dense linear algebra functions. Sparse matrix computations get best performance when they spend as much time as possible in dense linear algebra.

The C++ Standard Library includes many "mathematical special functions" ([sf.cmath]), like incomplete elliptic integrals, Bessel functions, and other polynomials and functions named after various mathematicians. Any of them comes with its own theory and set of applications for which robust and accurate implementations are indispensable. We think that linear algebra operations are at least as broadly useful, and in many cases significantly more so.

**Why base a C++ linear algebra library on the BLAS?**

1. The BLAS is a standard that codifies decades of existing practice.
2. The BLAS separates out "performance primitives" for hardware experts to tune, from mathematical operations that rely on those primitives for good performance.
3. Benchmarks reward hardware and system vendors for providing optimized BLAS implementations.
4. Writing a fast BLAS implementation for common element types is nontrivial, but well understood.
5. Optimized third-party BLAS implementations with liberal software licenses exist.
6. Building a C++ interface on top of the BLAS is a straightforward exercise, but has pitfalls for unaware developers.

Linear algebra has had a cross-language standard, the Basic Linear Algebra Subroutines (BLAS), since 2002. The Standard came out of a standardization process that started in 1995 and held meetings three times a year until 1999. Participants in the process came from industry, academia, and government research laboratories. The dense linear algebra subset of the BLAS codifies forty years of evolving practice, and has existed in recognizable form since 1990 (see P1417R0).

The BLAS interface was specifically designed as the distillation of the "computer science"/performance-oriented parts of linear algebra algorithms. It cleanly separates operations most critical for performance, from operations whose implementation takes expertise in mathematics and rounding-error analysis. This gives vendors opportunities to add value, without asking for expertise outside the typical required skill set of a Standard Library implementer.
Well-established benchmarks such as the LINPACK benchmark reward computer hardware vendors for optimizing their BLAS implementations. Thus, many vendors provide an optimized BLAS library for their computer architectures. Writing fast BLAS-like operations is not trivial, and depends on computer architecture. However, it is a well-understood problem whose solutions could be parameterized for a variety of computer architectures. See, for example, Goto and van de Geijn 2008. There are optimized third-party BLAS implementations for common architectures, like ATLAS and GotoBLAS. A (slow but correct) reference implementation of the BLAS exists and it has a liberal software license for easy reuse.

We have experience in the exercise of wrapping a C or Fortran BLAS implementation for use in portable C++ libraries. We describe this exercise in detail in our paper "Evolving a Standard C++ Linear Algebra Library from the BLAS" (P1674). It is straightforward for vendors, but has pitfalls for developers. For example, Fortran's application binary interface (ABI) differs across platforms in ways that can cause run-time errors (even incorrect results, not just crashing). Historical examples of vendors' C BLAS implementations have also had ABI issues that required work-arounds. This dependence on ABI details makes availability in a standard C++ library valuable.

Notation and conventions

The BLAS uses Fortran terms

The BLAS' "native" language is Fortran. It has a C binding as well, but the BLAS Standard and documentation use Fortran terms. Where applicable, we will call out relevant Fortran terms and highlight possibly confusing differences with corresponding C++ ideas. Our paper P1674R0 ("Evolving a Standard C++ Linear Algebra Library from the BLAS") goes into more detail on these issues.

We call "subroutines" functions

Like Fortran, the BLAS distinguishes between functions that return a value, and subroutines that do not return a value. In what follows, we will refer to both as "BLAS functions" or "functions."

Element types and BLAS function name prefix

The BLAS implements functionality for four different matrix, vector, or scalar element types:

- **REAL** (float in C++ terms)
- **DOUBLE PRECISION** (double in C++ terms)
- **COMPLEX** (complex<float> in C++ terms)
- **DOUBLE COMPLEX** (complex<double> in C++ terms)

The BLAS' Fortran 77 binding uses a function name prefix to distinguish functions based on element type:

- **S** for **REAL** ("single")
- **D** for **DOUBLE PRECISION**
- **C** for **COMPLEX**
- **Z** for **DOUBLE COMPLEX**
For example, the four BLAS functions \texttt{SAXPY}, \texttt{DAXPY}, \texttt{CAXPY}, and \texttt{ZAXPY} all perform the vector update \( Y = Y + \text{ALPHA} \times X \) for vectors \( X \) and \( Y \) and scalar \texttt{ALPHA}, but for different vector and scalar element types.

The convention is to refer to all of these functions together as \texttt{xAXPY}. In general, a lower-case \( x \) is a placeholder for all data type prefixes that the BLAS provides. For most functions, the \( x \) is a prefix, but for a few functions like \texttt{1xMAX}, the data type "prefix" is not the first letter of the function name. (\texttt{1xMAX} is a Fortran function that returns \texttt{INTEGER}, and therefore follows the old Fortran implicit naming rule that integers start with \( I, J \), etc.)

Not all BLAS functions exist for all four data types. These come in three categories:

1. The BLAS provides only real-arithmetic (\( S \) and \( D \)) versions of the function, since the function only makes mathematical sense in real arithmetic.

2. The complex-arithmetic versions perform a slightly different mathematical operation than the real-arithmetic versions, so they have a different base name.

3. The complex-arithmetic versions offer a choice between nonconjugated or conjugated operations.

As an example of the second category, the BLAS functions \texttt{SASUM} and \texttt{DASUM} compute the sums of absolute values of a vector's elements. Their complex counterparts \texttt{CSASUM} and \texttt{DZASUM} compute the sums of absolute values of real and imaginary components of a vector \( v \), that is, the sum of \( \text{abs} (\text{real}(v(i))) + \text{abs} (\text{imag}(v(i))) \) for all \( i \) in the domain of \( v \). The latter operation is still useful as a vector norm, but it requires fewer arithmetic operations.

Examples of the third category include the following:

- nonconjugated dot product \texttt{xDOTU} and conjugated dot product \texttt{xDOTC}; and
- rank-1 symmetric (\texttt{xGERU}) vs. Hermitian (\texttt{xGERC}) matrix update.

The conjugate transpose and the (nonconjugated) transpose are the same operation in real arithmetic (if one considers real arithmetic embedded in complex arithmetic), but differ in complex arithmetic. Different applications have different reasons to want either. The C++ Standard includes complex numbers, so a Standard linear algebra library needs to respect the mathematical structures that go along with complex numbers.

What we exclude from the design

Functions not in the Reference BLAS

The BLAS Standard includes functionality that appears neither in the Reference BLAS library, nor in the classic BLAS "level" 1, 2, and 3 papers. (For history of the BLAS "levels" and a bibliography, see \texttt{P1417R0}. For a paper describing functions not in the Reference BLAS, see "An updated set of basic linear algebra subprograms (BLAS)," listed in "Other references" below.) For example, the BLAS Standard has

- several new dense functions, like a fused vector update and dot product;
- sparse linear algebra functions, like sparse matrix-vector multiply and an interface for constructing sparse matrices; and
• extended- and mixed-precision dense functions (though we subsume some of their functionality; see below).

Our proposal only includes core Reference BLAS functionality, for the following reasons:

1. Vendors who implement a new component of the C++ Standard Library will want to see and test against an existing reference implementation.

2. Many applications that use sparse linear algebra also use dense, but not vice versa.

3. The Sparse BLAS interface is a stateful interface that is not consistent with the dense BLAS, and would need more extensive redesign to translate into a modern C++ idiom. See discussion in P1417R0.

4. Our proposal subsumes some dense mixed-precision functionality (see below).

LAPACK or related functionality

The LAPACK Fortran library implements solvers for the following classes of mathematical problems:

• linear systems,
• linear least-squares problems, and
• eigenvalue and singular value problems.

It also provides matrix factorizations and related linear algebra operations. LAPACK deliberately relies on the BLAS for good performance; in fact, LAPACK and the BLAS were designed together. See history presented in P1417R0.

Several C++ libraries provide slices of LAPACK functionality. Here is a brief, noninclusive list, in alphabetical order, of some libraries actively being maintained:

• Armadillo,
• Boost.uBLAS,
• Eigen,
• Matrix Template Library, and
• Trilinos.

P1417R0 gives some history of C++ linear algebra libraries. The authors of this proposal have designed, written, and maintained LAPACK wrappers in C++. Some authors have LAPACK founders as PhD advisors. Nevertheless, we have excluded LAPACK-like functionality from this proposal, for the following reasons:

1. LAPACK is a Fortran library, unlike the BLAS, which is a multilanguage standard.

2. We intend to support more general element types, beyond the four that LAPACK supports. It’s much more straightforward to make a C++ BLAS work for general element types, than to make LAPACK algorithms work generically.

First, unlike the BLAS, LAPACK is a Fortran library, not a standard. LAPACK was developed concurrently with the "level 3" BLAS functions, and the two projects share contributors. Nevertheless, only the BLAS and not LAPACK got standardized. Some vendors supply LAPACK
implementations with some optimized functions, but most implementations likely depend heavily on "reference" LAPACK. There have been a few efforts by LAPACK contributors to develop C++ LAPACK bindings, from Lapack++ in pre-templates C++ circa 1993, to the recent "C++ API for BLAS and LAPACK". (The latter shares coauthors with this proposal.) However, these are still just C++ bindings to a Fortran library. This means that if vendors had to supply C++ functionality equivalent to LAPACK, they would either need to start with a Fortran compiler, or would need to invest a lot of effort in a C++ reimplementation. Mechanical translation from Fortran to C++ introduces risk, because many LAPACK functions depend critically on details of floating-point arithmetic behavior.

Second, we intend to permit use of matrix or vector element types other than just the four types that the BLAS and LAPACK support. This includes "short" floating-point types, fixed-point types, integers, and user-defined arithmetic types. Doing this is easier for BLAS-like operations than for the much more complicated numerical algorithms in LAPACK. LAPACK strives for a "generic" design (see Jack Dongarra interview summary in P1417R0), but only supports two real floating-point types and two complex floating-point types. Directly translating LAPACK source code into a "generic" version could lead to pitfalls. Many LAPACK algorithms only make sense for number systems that aim to approximate real numbers (or their complex extensions). Some LAPACK functions output error bounds that rely on properties of floating-point arithmetic.

For these reasons, we have left LAPACK-like functionality for future work. It would be natural for a future LAPACK-like C++ library to build on our proposal.

Extended-precision BLAS

Our interface subsumes some functionality of the Mixed-Precision BLAS specification (Chapter 4 of the BLAS Standard). For example, users may multiply two 16-bit floating-point matrices (assuming that a 16-bit floating-point type exists) and accumulate into a 32-bit floating-point matrix, just by providing a 32-bit floating-point matrix as output. Users may specify the precision of a dot product result. If it is greater than the input vectors' element type precisions (e.g., double vs. float), then this effectively performs accumulation in higher precision. Our proposal imposes semantic requirements on some functions, like vector_norm2, to behave in this way.

However, we do not include the "Extended-Precision BLAS" in this proposal. The BLAS Standard lets callers decide at run time whether to use extended precision floating-point arithmetic for internal evaluations. We could support this feature at a later time. Implementations of our interface also have the freedom to use more accurate evaluation methods than typical BLAS implementations. For example, it is possible to make floating-point sums completely independent of parallel evaluation order.

Arithmetic operators and associated expression templates

Our proposal omits arithmetic operators on matrices and vectors. We do so for the following reasons:

1. We propose a low-level, minimal interface.
2. `operator*` could have multiple meanings for matrices and vectors. Should it mean elementwise product (like `valarray`) or matrix product? Should libraries reinterpret "vector times vector" as a dot product (row vector times column vector)? We prefer to let a higher-level library decide this, and make everything explicit at our lower level.

3. Arithmetic operators require defining the element type of the vector or matrix returned by an expression. Functions let users specify this explicitly, and even let users use different output types for the same input types in different expressions.

4. Arithmetic operators may require allocation of temporary matrix or vector storage. This prevents use of nonowning data structures.

5. Arithmetic operators strongly suggest expression templates. These introduce problems such as dangling references and aliasing.

Our goal is to propose a low-level interface. Other libraries, such as that proposed by P1385, could use our interface to implement overloaded arithmetic for matrices and vectors. P0939R0 advocates using "an incremental approach to design to benefit from actual experience." A constrained, function-based, BLAS-like interface builds incrementally on the many years of BLAS experience.

Arithmetic operators on matrices and vectors would require the library, not necessarily the user, to specify the element type of an expression's result. This gets tricky if the terms have mixed element types. For example, what should the element type of the result of the vector sum `x + y` be, if `x` has element type `complex<float>` and `y` has element type `double`? It's tempting to use `common_type_t`, but `common_type_t<complex<float>, double>` is `complex<float>`. This loses precision. Some users may want `complex<double>`; others may want `complex<long double>` or something else, and others may want to choose different types in the same program.

P1385 lets users customize the return type of such arithmetic expressions. However, different algorithms may call for the same expression with the same inputs to have different output types. For example, iterative refinement of linear systems `Ax=b` can work either with an extended-precision intermediate residual vector `r = b - A*x`, or with a residual vector that has the same precision as the input linear system. Each choice produces a different algorithm with different convergence characteristics, per-iteration run time, and memory requirements. Thus, our library lets users specify the result element type of linear algebra operations explicitly, by calling a named function that takes an output argument explicitly, rather than an arithmetic operator.

Arithmetic operators on matrices or vectors may also need to allocate temporary storage. Users may not want that. When LAPACK's developers switched from Fortran 77 to a subset of Fortran 90, their users rejected the option of letting LAPACK functions allocate temporary storage on their own. Users wanted to control memory allocation. Also, allocating storage precludes use of nonowning input data structures like `basic_mdspan`, that do not know how to allocate.

Arithmetic expressions on matrices or vectors strongly suggest expression templates, as a way to avoid allocation of temporaries and to fuse computational kernels. They do not require expression templates. For example, `valarray` offers overloaded operators for vector
arithmetic, but the Standard lets implementers decide whether to use expression templates. However, all of the current C++ linear algebra libraries that we mentioned above have some form of expression templates for overloaded arithmetic operators, so users will expect this and rely on it for good performance. This was, indeed, one of the major complaints about initial implementations of `valarray`: its lack of mandate for expression templates meant that initial implementations were slow, and thus users did not want to rely on it. (See Josuttis 1999, p. 547, and Vandevoorde and Josuttis 2003, p. 342, for a summary of the history. Fortran has an analogous issue, in which (under certain conditions) it is implementation defined whether the run-time environment needs to copy noncontiguous slices of an array into contiguous temporary storage.)

Expression templates work well, but have issues. Our papers P1417R0 and "Evolving a Standard C++ Linear Algebra Library from the BLAS" (P1674R0) give more detail on these issues. A particularly troublesome one is that modern C++ `auto` makes it easy for users to capture expressions before their evaluation and writing into an output array. For matrices and vectors with container semantics, this makes it easy to create dangling references. Users might not realize that they need to assign expressions to named types before actual work and storage happen. `Eigen's documentation` describes this common problem.

Our `scaled_view`, `conjugate_view`, and `transpose_view` functions make use of one aspect of expression templates, namely modifying the `basic_mdspan` array access operator. However, we intend these functions for use only as in-place modifications of arguments of a function call. Also, when modifying `basic_mdspan`, these functions merely view the same data that their input `basic_mdspan` views. They introduce no more potential for dangling references than `basic_mdspan` itself. The use of views like `basic_mdspan` is self-documenting; it tells users that they need to take responsibility for scope of the viewed data.

**Banded matrix layouts**

This proposal omits banded matrix types. It would be easy to add the required layouts and specializations of algorithms later. The packed and unpacked symmetric and triangular layouts in this proposal cover the major concerns that would arise in the banded case, like nonstrided and nonunique layouts, and matrix types that forbid access to some multi-indices in the Cartesian product of extents.

**Tensors**

We exclude tensors from this proposal, for the following reasons. First, tensor libraries naturally build on optimized dense linear algebra libraries like the BLAS, so a linear algebra library is a good first step. Second, `mdspan` has natural use as a low-level representation of dense tensors, so we are already partway there. Third, even simple tensor operations that naturally generalize the BLAS have infinitely many more cases than linear algebra. It's not clear to us which to optimize. Fourth, even though linear algebra is a special case of tensor algebra, users of linear algebra have different interface expectations than users of tensor algebra. Thus, it makes sense to have two separate interfaces.

**Design justification**
We take a step-wise approach. We begin with core BLAS dense linear algebra functionality. We then deviate from that only as much as necessary to get algorithms that behave as much as reasonable like the existing C++ Standard Library algorithms. Future work or collaboration with other proposals could implement a higher-level interface.

We propose to build the initial interface on top of `basic_mdspan`, and plan to extend that later with overloads for a new `basic_mdarray` variant of `basic_mdspan` with container semantics as well as any type implementing a `get_mdspan` customization point. We explain the value of these choices below.

Please refer to our papers "Evolving a Standard C++ Linear Algebra Library from the BLAS" (P1674R0) and "Historical lessons for C++ linear algebra library standardization" (P1417R0). They will give details and references for many of the points that we summarize here.

We do not require using the BLAS library

Our proposal is based on the BLAS interface, and it would be natural for implementers to use an existing C or Fortran BLAS library. However, we do not require an underlying BLAS C interface. Vendors should have the freedom to decide whether they want to rely on an existing BLAS library. They may also want to write a "pure" C++ implementation that does not depend on an external library. They will, in any case, need a "generic" C++ implementation for matrix and vector element types other than the four that the BLAS supports.

Why use `basic_mdspan`?

- C++ does not currently have a data structure for representing multidimensional arrays.

- The BLAS' C interface takes a large number of pointer and integer arguments that represent matrices and vectors. Using multidimensional array data structures in the C++ interface reduces the number of arguments and avoids common errors.

- `basic_mdspan` supports row-major, column-major, and strided layouts out of the box, and have `Layout` as an extension point. This lets our interface support layouts beyond what the BLAS Standard permits.

- They can exploit any dimensions or strides known at compile time.

- They have built-in "slicing" capabilities via `subspan`.

- Their layout and accessor policies will let us simplify our interfaces even further, by encapsulating transpose, conjugate, and scalar arguments. See below for details.

- `basic_mdspan` is low level; it imposes no mathematical meaning on multidimensional arrays. This gives users the freedom to develop mathematical libraries with the semantics they want. (Some users object to calling something a "matrix" or "tensor" if it doesn't have the right mathematical properties. The Standard has already taken the word `vector`.)

- They offer a hook for future expansion to support heterogenous memory spaces. (This is a key feature of `Kokkos::View`, the data structure that inspired `basic_mdspan`.)
• Their encapsulation of matrix indexing makes C++ implementations of BLAS-like operations much less error prone and easier to read.

• They make it easier to support an efficient "batched" interface.

**Defining a concept for the data structures instead**

LEWGI requested in the 2019 Cologne meeting that we explore using a concept instead of `basic_mdspan` to define the arguments for the linear algebra functions. Our investigation of this option leads us to believe that such a concept would largely replicate the definition of `basic_mdspan`. This proposal refers to almost all of its features, including extents, layout, and accessor_policy. We expect implementations to use all of them for optimizations, for example to extract the scaling factor from a scaled_view result in order to call an optimized BLAS library directly. At this point it is not clear to us that there would be significant benefits to defining linear algebra functions in terms of such a concept, instead of defining a general customization point `get_mdspan`, allowing acceptance of any object type, which provides it. After further discussion at the 2019 Belfast meeting, LEWGI accepted our position.

**Function argument aliasing and zero scalar multipliers**

Summary:

1. The BLAS Standard forbids aliasing any input (read-only) argument with any output (write-only or read-and-write) argument.

2. The BLAS uses `INTENT(INOUT)` (read-and-write) arguments to express "updates" to a vector or matrix. By contrast, C++ Standard algorithms like `transform` take input and output iterator ranges as different parameters, but may let input and output ranges be the same.

3. The BLAS uses the values of scalar multiplier arguments ("alpha" or "beta") of vectors or matrices at run time, to decide whether to treat the vectors or matrices as write only. This matters both for performance and semantically, assuming IEEE floating-point arithmetic.

4. We decide separately, based on the category of BLAS function, how to translate `INTENT(INOUT)` arguments into a C++ idiom:

   a. For triangular solve and triangular multiply, in-place behavior is essential for computing matrix factorizations in place, without requiring extra storage proportional to the input matrix's dimensions. However, in-place functions cannot be parallelized for arbitrary execution policies. Thus, we have both not-in-place and in-place overloads, and only the not-in-place overloads take an optional `ExecutionPolicy&&`.

   b. Else, if the BLAS function unconditionally updates (like `xGER`), we retain read-and-write behavior for that argument.

   c. Else, if the BLAS function uses a scalar `beta` argument to decide whether to read the output argument as well as write to it (like `xGEMM`), we provide two versions: a write-only version (as if `beta` is zero), and a read-and-write version (as if `beta` is nonzero).
For a detailed analysis, see "Evolving a Standard C++ Linear Algebra Library from the BLAS" (P1674R0).

Support for different matrix layouts

Summary:

1. The dense BLAS supports several different dense matrix "types." Type is a mixture of "storage format" (e.g., packed, banded) and "mathematical property" (e.g., symmetric, Hermitian, triangular).

2. Some "types" can be expressed as custom basic_mdspan layouts. Other types actually represent algorithmic constraints: for instance, what entries of the matrix the algorithm is allowed to access.

3. Thus, a C++ BLAS wrapper cannot overload on matrix "type" simply by overloading on basic_mdspan specialization. The wrapper must use different function names, tags, or some other way to decide what the matrix type is.

For more details, including a list and description of the matrix "types" that the dense BLAS supports, see our paper "Evolving a Standard C++ Linear Algebra Library from the BLAS" (P1674R0) lists the different matrix types.

A C++ linear algebra library has a few possibilities for distinguishing the matrix "type":

1. It could imitate the BLAS, by introducing different function names, if the layouts and accessors do not sufficiently describe the arguments.

2. It could introduce a hierarchy of higher-level classes for representing linear algebra objects, use basic_mdspan (or something like it) underneath, and write algorithms to those higher-level classes.

3. It could use the layout and accessor types in basic_mdspan simply as tags to indicate the matrix "type." Algorithms could specialize on those tags.

We have chosen Approach 1. Our view is that a BLAS-like interface should be as low-level as possible. Approach 2 is more like a "Matlab in C++"; a library that implements this could build on our proposal's lower-level library. Approach 3 sounds attractive. However, most BLAS matrix "types" do not have a natural representation as layouts. Trying to hack them in would pollute basic_mdspan -- a simple class meant to be easy for the compiler to optimize -- with extra baggage for representing what amounts to sparse matrices. We think that BLAS matrix "type" is better represented with a higher-level library that builds on our proposal.

Over- and underflow wording for vector 2-norm

SG6 recommended to us at Belfast 2019 to change the special overflow / underflow wording for vector_norm2 to imitate the BLAS Standard more closely. The BLAS Standard does say something about overflow and underflow for vector 2-norms. We reviewed this wording and conclude that it is either a nonbinding quality of implementation (QoI) recommendation, or too vaguely stated to translate directly into C++ Standard wording. Thus, we removed the special overflow / underflow wording entirely.
Previous versions of this paper asked implementations to compute vector 2-norms "without undue overflow or underflow at intermediate stages of the computation." "Undue" imitates existing C++ Standard wording for hypot. This wording hints at the stricter requirements in F.9 (normative, but optional) of the C Standard for math library functions like hypot, without mandating those requirements. In particular, paragraph 9 of F.9 says:

> Whether or when library functions raise an undeserved "underflow" floating-point exception is unspecified. Otherwise, as implied by F.7.6, the `<math.h>` functions do not raise spurious floating-point exceptions (detectable by the user) [including the "overflow" exception discussed in paragraph 6], other than the "inexact" floating-point exception.

However, these requirements are for math library functions like hypot, not for general algorithms that return floating-point values. SG6 did not raise a concern that we should treat vector_norm2 like a math library function; their concern was that we imitate the BLAS Standard's wording.

The BLAS Standard says of several operations, including vector 2-norm: "Here are the exceptional routines where we ask for particularly careful implementations to avoid unnecessary over/underflows, that could make the output unnecessarily inaccurate or unreliable" (p. 35).

The BLAS Standard does not define phrases like "unnecessary over/underflows." The likely intent is to avoid naïve implementations that simply add up the squares of the vector elements. These would overflow even if the norm in exact arithmetic is significantly less than the overflow threshold. The POSIX Standard (IEEE Std 1003.1-2017) analogously says that hypot must "take precautions against overflow during intermediate steps of the computation."

The phrase "precautions against overflow" is too vague for us to translate into a requirement. The authors likely meant to exclude naïve implementations, but not require implementations to know whether a result computed in exact arithmetic would overflow or underflow. The latter is a special case of computing floating-point sums exactly, which is costly for vectors of arbitrary length. While it would be a useful feature, it is difficult enough that we do not want to require it, especially since the BLAS Standard itself does not.

For all of the functions listed on p. 35 of the BLAS Standard as needing "particularly careful implementations," except vector norm, the BLAS Standard has an "Advice to implementors" section with extra accuracy requirements. The BLAS Standard does have an "Advice to implementors" section for matrix norms (see Section 2.8.7, p. 69), which have similar over- and underflow concerns as vector norms. However, the Standard merely states that "[h]igh-quality implementations of these routines should be accurate" and should document their accuracy, and gives examples of "accurate implementations" in LAPACK.

The BLAS Standard never defines what "Advice to implementors" means. However, the BLAS Standard shares coauthors and audience with the Message Passing Interface (MPI) Standard, which defines "Advice to implementors" as "primarily commentary to implementors" and permissible to skip (see e.g., MPI 3.0, Section 2.1, p. 9). We thus interpret "Advice to implementors" in the BLAS Standard as a nonbinding quality of implementation (QoI) recommendation.
Future work

Summary:

1. Generalize function parameters to take any type that implements the `get_mdspan` customization point, including `basic_mdarray`.

2. Add batched linear algebra overloads.

Generalize function parameters

Our functions differ from the C++ Standard algorithms, in that they take a concrete type `basic_mdspan` with template parameters, rather than any type that satisfies a concept. We think that the template parameters of `basic_mdspan` fully describe the multidimensional equivalent of a multipass iterator, and that "conceptification" of multidimensional arrays would unnecessarily delay both this proposal and P0009 (the `basic_mdspan` proposal).

In a future proposal, we plan to generalize our function's template parameters, to permit any type besides `basic_mdspan` that implements the `get_mdspan` customization point, as long as the return value of `get_mdspan` satisfies the current requirements. `get_mdspan` will return a `basic_mdspan` that views its argument's data.

`basic_mdarray`, proposed in P1684, is the container analog of `basic_mdspan`. It is a new kind of container, with the same copy behavior as containers like `vector`. It has the same extension points as `basic_mdspan`, and also has the ability to use any contiguous container (see [container.requirements.general]) for storage. Contiguity matters because `basic_mdspan` views a subset of a contiguous pointer range, and we want to be able to get a `basic_mdspan` that views the `basic_mdarray`. `basic_mdarray` will come with support for two different underlying containers: `array` and `vector`. A `subspan` (see P0009) of a `basic_mdarray` will return a `basic_mdspan` with the appropriate layout and corresponding accessor. Users must guard against dangling pointers, just as they currently must do when using `span` to view a subset of a `vector`.

Previous versions of this proposal included function overloads that took `basic_mdarray` directly. The goals were user convenience, and to avoid any potential overhead of conversion to `basic_mdspan`, especially for very small matrices and vectors. In a future revision of P1684, `basic_mdarray` will implement `get_mdspan`. This will let users use `basic_mdarray` directly in our functions. This customization point approach would also simplify using our functions with other matrix and vector types, such as those proposed by P1385. Implementations may optionally add direct overloads of our functions for `basic_mdarray` or other types. This would address any concerns about overhead of converting from `basic_mdarray` to `basic_mdspan`.

Batched linear algebra

We plan to write a separate proposal that will add "batched" versions of linear algebra functions to this proposal. "Batched" linear algebra functions solve many independent problems all at once, in a single function call. For discussion, see Section 6.2 of our background paper P1417R0. Batched interfaces have the following advantages:

- They expose more parallelism and vectorization opportunities for many small linear algebra operations.
They are useful for many different fields, including machine learning.

Hardware vendors currently offer both hardware features and optimized software libraries to support batched linear algebra.

There is an ongoing interface standardization effort, in which we participate.

The basic_mdspan data structure makes it easy to represent a batch of linear algebra objects, and to optimize their data layout.

With few exceptions, the extension of this proposal to support batched operations will not require new functions or interface changes. Only the requirements on functions will change. Output arguments can have an additional rank; if so, then the leftmost extent will refer to the batch dimension. Input arguments may also have an additional rank to match; if they do not, the function will use ("broadcast") the same input argument for all the output arguments in the batch.

Data structures and utilities borrowed from other proposals

basic_mdspan

This proposal depends on P0009R9, which is a proposal for adding multidimensional arrays to the C++ Standard Library. basic_mdspan is the main class in P0009. It is a "view" (in the sense of span) of a multidimensional array. The rank (number of dimensions) is fixed at compile time. Users may specify some dimensions at run time and others at compile time; the type of the basic_mdspan expresses this. basic_mdspan also has two customization points:

- Layout expresses the array's memory layout: e.g., row-major (C++ style), column-major (Fortran style), or strided. We use a custom Layout later in this paper to implement a "transpose view" of an existing basic_mdspan.

- Accessor defines the storage handle (i.e., pointer) stored in the mdspan, as well as the reference type returned by its access operator. This is an extension point for modifying how access happens, for example by using atomic_ref to get atomic access to every element. We use custom Accessors later in this paper to implement "scaled views" and "conjugated views" of an existing basic_mdspan.

The basic_mdspan class has an alias mdspan that uses the default Layout and Accessor. In this paper, when we refer to mdspan without other qualifiers, we mean the most general basic_mdspan.

New basic_mdspan layouts in this proposal

Our proposal uses the layout mapping policy of basic_mdspan in order to represent different matrix and vector data layouts. Layout mapping policies as described by P0009R9 have three basic properties:

- Unique
- Contiguous
- Strided
P0009R9 includes three different layouts -- `layout_left`, `layout_right`, and `layout_stride` -- all of which are unique and strided. Only `layout_left` and `layout_right` are contiguous.

This proposal includes the following additional layouts:

- **`layout_blas_general`**: Generalization of `layout_left` and `layout_right`; describes layout used by General (GE) matrix "type"

- **`layout_blas_packed`**: Describes layout used by the BLAS' Symmetric Packed (SP), Hermitian Packed (HP), and Triangular Packed (TP) "types"

These layouts have "tag" template parameters that control their properties; see below.

We do not include layouts for unpacked "types," such as Symmetric (SY), Hermitian (HE), and Triangular (TR). P1674 explains our reasoning. In summary: Their actual layout -- the arrangement of matrix elements in memory -- is the same as General. The only differences are constraints on what entries of the matrix algorithms may access, and assumptions about the matrix's mathematical properties. Trying to express those constraints or assumptions as "layouts" or "accessors" violates the spirit (and sometimes the law) of `basic_mdspan`. We address these different matrix types with different function names.

The packed matrix "types" do describe actual arrangements of matrix elements in memory that are not the same as in General. This is why we provide `layout_blas_packed`. Note that `layout_blas_packed` is the first addition to the layouts in P0009R9 that is neither always unique, nor always strided.

Algorithms cannot be written generically if they permit output arguments with nonunique layouts. Nonunique output arguments require specialization of the algorithm to the layout, since there's no way to know generically at compile time what indices map to the same matrix element. Thus, we will impose the following rule: Any `basic_mdspan` output argument to our functions must always have unique layout (`is_always_unique()` is true), unless otherwise specified.

Some of our functions explicitly require outputs with specific nonunique layouts. This includes low-rank updates to symmetric or Hermitian matrices.

**Acknowledgments**

Sandia National Laboratories is a multimission laboratory managed and operated by National Technology & Engineering Solutions of Sandia, LLC, a wholly owned subsidiary of Honeywell International, Inc., for the U.S. Department of Energy’s National Nuclear Security Administration under contract DE-NA0003525.

Special thanks to Bob Steagall and Guy Davidson for boldly leading the charge to add linear algebra to the C++ Standard Library, and for many fruitful discussions. Thanks also to Andrew Lumsdaine for his pioneering efforts and history lessons.

**References**

References by coauthors


Other references


Wording

Text in blockquotes is not proposed wording, but rather instructions for generating proposed wording. The ◆ character is used to denote a placeholder section number which the editor shall determine. First, apply all wording from P0009R9 (this proposal is a "rebase" atop the changes proposed by P0009R9). At the end of Table ◆ ("Numerics library summary") in [numerics.general], add the following: [linalg], Linear algebra, <linalg>. At the end of [numerics], add all the material that follows.

Header <linalg> synopsis [linalg.syn]

namespace std {
    // [linalg.tags.order], storage order tags
    struct column_major_t;
    inline constexpr column_major_t column_major;
    struct row_major_t;
    inline constexpr row_major_t row_major;

    // [linalg.tags.triangle], triangle tags
    struct upper_triangle_t;
    inline constexpr upper_triangle_t upper_triangle;
    struct lower_triangle_t;
    inline constexpr lower_triangle_t lower_triangle;

    // [linalg.tags.diagonal], diagonal tags
    struct implicit_unit_diagonal_t;
    inline constexpr implicit_unit_diagonal_t implicit_unit_diagonal;
    struct explicit_diagonal_t;
    inline constexpr explicit_diagonal_t explicit_diagonal;

    // [linalg.layouts.general], class template layout_blas_general
    template<class StorageOrder>
    class layout_blas_general;

    // [linalg.layouts.packed], class template layout_blas_packed
    template<class Triangle, 
             class StorageOrder>
    class layout_blas_packed;

    // [linalg.scaled.accessor_scaled], class template accessor_scaled
    template<class ScalingFactor, 
             class Accessor>
    class accessor_scaled;

    // [linalg.scaled.scaled_view], scaled in-place transformation
    template<class ScalingFactor, 
             class ElementType>

class Extents,
class Layout,
class Accessor>

/* see-below */
scaled_view(
   const ScalingFactor & s,
   const basic_mdspan<ElementType, Extents, Layout, Accessor>& a);

// [linalg.conj.accessor_conjugate], class template accessor_conjugate
template<class Accessor>
class accessor_conjugate;

// [linalg.conj.conjugate_view], conjugated in-place transformation
template<class ElementType, class Extents, class Layout, class Accessor>
/* see-below */
conjugate_view(
   basic_mdspan<ElementType, Extents, Layout, Accessor> a);

// [linalg.transp.layout_transpose], class template layout_transpose
template<class Layout>
class layout_transpose;

// [linalg.transp.transpose_view], transposed in-place transformation
template<class ElementType, class Extents, class Layout, class Accessor>
/* see-below */
transpose_view(
   basic_mdspan<ElementType, Extents, Layout, Accessor> a);

// [linalg.conj_transp],
// conjugated transposed in-place transformation
template<class ElementType, class Extents, class Layout, class Accessor>
/* see-below */
conjugate_transpose_view(
   basic_mdspan<ElementType, Extents, Layout, Accessor> a);

// [linalg.algs.blas1.givens.lartg], compute Givens rotation
template<class Real>
void givens_rotation_setup(const Real a,
   const Real b,
   Real & c,
   Real & s,
   Real & r);

template<class Real>
void givens_rotation_setup(const complex<Real>& a,
   const complex<Real>& a,
   Real & c,
   complex<Real>& s,
   complex<Real>& r);
// [linalg.algs.blas1.givens.rot], apply computed Givens rotation
template<
class inout_vector_1_t,
class inout_vector_2_t,
class Real>
void givens_rotation_apply(
inout_vector_1_t x,
inout_vector_2_t y,
const Real c,
const Real s);

template<class ExecutionPolicy,
class inout_vector_1_t,
class inout_vector_2_t,
class Real>
void givens_rotation_apply(
    ExecutionPolicy&& exec,
inout_vector_1_t x,
inout_vector_2_t y,
const Real c,
const Real s);

template<class inout_vector_1_t,
class inout_vector_2_t,
class Real>
void givens_rotation_apply(
inout_vector_1_t x,
inout_vector_2_t y,
const Real c,
const complex<Real> s);

template<class ExecutionPolicy,
class inout_vector_1_t,
class inout_vector_2_t,
class Real>
void givens_rotation_apply(
    ExecutionPolicy&& exec,
inout_vector_1_t x,
inout_vector_2_t y,
const Real c,
const complex<Real> s);

// [linalg.algs.blas1.swap], swap elements
template<class inout_object_1_t,
class inout_object_2_t>
void linalg_swap(inout_object_1_t x,
inout_object_2_t y);

template<class ExecutionPolicy,
class inout_object_1_t,
class inout_object_2_t>
void linalg_swap(ExecutionPolicy&& exec,
inout_object_1_t x,
inout_object_2_t y);

// [linalg.algs.blas1.scal], multiply elements by scalar
template<class Scalar,
class inout_object_t>
void scale(const Scalar alpha,
inout_object_t obj);

template<class ExecutionPolicy,
class Scalar,
class inout_object_t>
void scale(ExecutionPolicy&& exec,
const Scalar alpha,
inout_object_t obj);

// [linalg.algs.blas1.copy], copy elements
template<class in_object_t,
class out_object_t>
void linalg_copy(in_object_t x,
out_object_t y);

template<class ExecutionPolicy,
class in_object_t,
class out_object_t>
void linalg_copy(ExecutionPolicy&& exec,
in_object_t x,
out_object_t y);

// [linalg.algs.blas1.add], add elementwise
template<class in_object_1_t,
class in_object_2_t,
class out_object_t>
void linalg_add(in_object_1_t x,
in_object_2_t y,
out_object_t z);

template<class ExecutionPolicy,
class in_object_1_t,
class in_object_2_t,
class out_object_t>
void linalg_add(ExecutionPolicy&& exec,
in_object_1_t x,
in_object_2_t y,
out_object_t z);

// [linalg.algs.blas1.dot],
// dot product of two vectors

// [linalg.algs.blas1.dot.dotu],
// nonconjugated dot product of two vectors
template<class in_vector_1_t,
class in_vector_2_t,
class T>
T dot(in_vector_1_t v1,
in_vector_2_t v2,
T init);

template<class ExecutionPolicy,
class in_vector_1_t,
class in_vector_2_t,
class T>
T dot(ExecutionPolicy&& exec,
in_vector_1_t v1,
in_vector_2_t v2,
T init);

template<class in_vector_1_t,
class in_vector_2_t>
auto dot(in_vector_1_t v1,
in_vector_2_t v2);

template<class ExecutionPolicy,
```cpp
class in_vector_1_t,
class in_vector_2_t>
auto dot(ExecutionPolicy&& exec,
in_vector_1_t v1,
in_vector_2_t v2);

// [linalg.algs.blas1.dot.dotc],
// conjugated dot product of two vectors
template<class in_vector_1_t,
class in_vector_2_t,
class T>
T dotc(in_vector_1_t v1,
in_vector_2_t v2,
T init);
template<class ExecutionPolicy,
class in_vector_1_t,
class in_vector_2_t,
class T>
T dotc(ExecutionPolicy&& exec,
in_vector_1_t v1,
in_vector_2_t v2,
T init);
template<class in_vector_1_t,
class in_vector_2_t>
auto dotc(in_vector_1_t v1,
in_vector_2_t v2);
template<class ExecutionPolicy,
class in_vector_1_t,
class in_vector_2_t>
auto dotc(ExecutionPolicy&& exec,
in_vector_1_t v1,
in_vector_2_t v2);

// [linalg.algs.blas1.nrm2],
// Euclidean norm of a vector
template<class in_vector_t,
class T>
T vector_norm2(in_vector_t v,
T init);
template<class ExecutionPolicy,
class in_vector_t,
class T>
T vector_norm2(ExecutionPolicy&& exec,
in_vector_t v,
T init);
template<class in_vector_t>
auto vector_norm2(in_vector_t v);
template<class ExecutionPolicy,
class in_vector_t>
auto vector_norm2(ExecutionPolicy&& exec,
in_vector_t v);

// [linalg.algs.blas1.asum],
// sum of absolute values of vector elements
template<class in_vector_t,
class T>
T vector_abs_sum(in_vector_t v,
T init);
```
```cpp
template<class ExecutionPolicy,
         class in_vector_t,
         class T>
T vector_abs_sum(ExecutionPolicy&& exec,
                 in_vector_t v,
                 T init);

template<class in_vector_t>
auto vector_abs_sum(in_vector_t v);

template<class ExecutionPolicy,
         class in_vector_t>
auto vector_abs_sum(ExecutionPolicy&& exec,
                 in_vector_t v);

// [linalg.algs.blas1.iamax],
// index of maximum absolute value of vector elements
template<class in_vector_t>
ptrdiff_t idx_abs_max(in_vector_t v);

template<class ExecutionPolicy,
         class in_vector_t>
ptrdiff_t idx_abs_max(ExecutionPolicy&& exec,
                 in_vector_t v);

// [linalg.algs.blas2.gemv],
// general matrix-vector product
template<class in_vector_t,
         class in_matrix_t,
         class out_vector_t>
void matrix_vector_product(in_matrix_t A,
                          in_vector_t x,
                          out_vector_t y);

template<class ExecutionPolicy,
         class in_vector_t,
         class in_matrix_t,
         class out_vector_t>
void matrix_vector_product(ExecutionPolicy&& exec,
                          in_matrix_t A,
                          in_vector_t x,
                          out_vector_t y);

template<class in_vector_1_t,
         class in_matrix_t,
         class in_vector_2_t,
         class out_vector_t>
void matrix_vector_product(in_matrix_t A,
                          in_vector_1_t x,
                          in_vector_2_t y,
                          out_vector_t z);

template<class ExecutionPolicy,
         class in_vector_1_t,
         class in_matrix_t,
         class in_vector_2_t,
         class out_vector_t>
void matrix_vector_product(ExecutionPolicy&& exec,
                          in_matrix_t A,
                          in_vector_1_t x,
                          in_vector_2_t y,
                          out_vector_t z);

// [linalg.algs.blas2.symv],
```
template<
class in_matrix_t,
class Triangle,
class in_vector_t,
class out_vector_t>
void symmetric_matrix_vector_product(in_matrix_t A,
Triangle t,
in_vector_t x,
out_vector_t y);

template<
class ExecutionPolicy,
class in_matrix_t,
class Triangle,
class in_vector_t,
class out_vector_t>
void symmetric_matrix_vector_product(ExecutionPolicy&
exec,
in_matrix_t A,
Triangle t,
in_vector_t x,
out_vector_t y);

template<
class in_matrix_t,
class Triangle,
class in_vector_1_t,
class in_vector_2_t,
class out_vector_t>
void symmetric_matrix_vector_product(
in_matrix_t A,
Triangle t,
in_vector_1_t x,
in_vector_2_t y,
out_vector_t z);

template<
class ExecutionPolicy,
class in_matrix_t,
class Triangle,
class in_vector_1_t,
class in_vector_2_t,
class out_vector_t>
void symmetric_matrix_vector_product(
ExecutionPolicy&
exec,
in_matrix_t A,
Triangle t,
in_vector_1_t x,
in_vector_2_t y,
out_vector_t z);

// [linalg.algs.blas2.hemv], // Hermitian matrix-vector product
template<
class in_matrix_t,
class Triangle,
class in_vector_t,
class out_vector_t>
void hermitian_matrix_vector_product(in_matrix_t A,
Triangle t,
in_vector_t x,
out_vector_t y);

template<
class ExecutionPolicy,
class in_matrix_t,
class Triangle,
class out_vector_t>
void hermitian_matrix_vector_product(ExecutionPolicy&
in_matrix_t A,
Triangle t,
in_vector_t x,
out_vector_t y);
template<class in_matrix_t, class in_vector_t, class out_vector_t>
void hermitian_matrix_vector_product(ExecutionPolicy&& exec,
    in_matrix_t A,
    Triangle t,
    in_vector_t x,
    out_vector_t y);

template<class in_matrix_t, class Triangle, class in_vector_1_t, class in_vector_2_t, class out_vector_t>
void hermitian_matrix_vector_product(in_matrix_t A,
    Triangle t,
    in_vector_1_t x,
    in_vector_2_t y,
    out_vector_t z);

template<class ExecutionPolicy, class in_matrix_t, class Triangle, class in_vector_1_t, class in_vector_2_t, class out_vector_t>
void hermitian_matrix_vector_product(ExecutionPolicy&& exec,
    in_matrix_t A,
    Triangle t,
    in_vector_1_t x,
    in_vector_2_t y,
    out_vector_t z);

// [linalg.algs.blas2.trmv],
// Triangular matrix-vector product

// [linalg.algs.blas2.trmv.ov],
// Overwriting triangular matrix-vector product
template<class in_matrix_t, class Triangle, class DiagonalStorage, class in_vector_t, class out_vector_t>
void triangular_matrix_vector_product(
in_matrix_t A,
    Triangle t,
    DiagonalStorage d,
    in_vector_t x,
    out_vector_t y);

template<class ExecutionPolicy, class in_matrix_t, class Triangle, class DiagonalStorage, class in_vector_t, class out_vector_t>
void triangular_matrix_vector_product(
    ExecutionPolicy&& exec,
    in_matrix_t A,
    Triangle t,
    DiagonalStorage d,
```cpp
in_vector_t x,
  out_vector_t y);

// [linalg.algs.blas2.trmv.in-place],
// In-place triangular matrix-vector product
template<class in_matrix_t,
       class Triangle,
       class DiagonalStorage,
       class inout_vector_t>
void triangular_matrix_vector_product(
    in_matrix_t A,
    Triangle t,
    DiagonalStorage d,
    inout_vector_t y);

// [linalg.algs.blas2.trmv.up],
// Updating triangular matrix-vector product
template<class in_matrix_t,
       class Triangle,
       class DiagonalStorage,
       class in_vector_1_t,
       class in_vector_2_t,
       class out_vector_t>
void triangular_matrix_vector_product(
    in_matrix_t A,
    Triangle t,
    DiagonalStorage d,
    in_vector_1_t x,
    in_vector_2_t y,
    out_vector_t z);

template<class ExecutionPolicy,
       class in_matrix_t,
       class Triangle,
       class DiagonalStorage,
       class in_vector_1_t,
       class in_vector_2_t,
       class out_vector_t>
void triangular_matrix_vector_product(ExecutionPolicy&& exec,
    in_matrix_t A,
    Triangle t,
    DiagonalStorage d,
    in_vector_1_t x,
    in_vector_2_t y,
    out_vector_t z);

// [linalg.algs.blas2.trsv],
// Solve a triangular linear system

// [linalg.algs.blas2.trsv.not-in-place],
// Solve a triangular linear system, not in place
template<class in_matrix_t,
       class Triangle,
       class DiagonalStorage,
       class in_vector_t,
       class out_vector_t>
void triangular_matrix_vector_solve(
    in_matrix_t A,
    Triangle t,
    DiagonalStorage d,
...}
```
```cpp
in_vector_t b,
out_vector_t x);

template<class ExecutionPolicy,
  class in_matrix_t,
  class Triangle,
  class DiagonalStorage,
  class in_vector_t,
  class out_vector_t>
void triangular_matrix_vector_solve(
  ExecutionPolicy&& exec,
  in_matrix_t A,
  Triangle t,
  DiagonalStorage d,
  in_vector_t b,
  out_vector_t x);

// [linalg.algs.blas2.trsv.in-place],
// Solve a triangular linear system, in place
template<class in_matrix_t,
  class Triangle,
  class DiagonalStorage,
  class inout_vector_t>
void triangular_matrix_vector_solve(
  in_matrix_t A,
  Triangle t,
  DiagonalStorage d,
  inout_vector_t b);

// [linalg.algs.blas2.rank1.geru],
// nonconjugated rank-1 matrix update
template<class in_vector_1_t,
  class in_vector_2_t,
  class inout_matrix_t>
void matrix_rank_1_update(
  in_vector_1_t x,
  in_vector_2_t y,
  inout_matrix_t A);

template<class ExecutionPolicy,
  class in_vector_1_t,
  class in_vector_2_t,
  class inout_matrix_t>
void matrix_rank_1_update(
  ExecutionPolicy&& exec,
  in_vector_1_t x,
  in_vector_2_t y,
  inout_matrix_t A);

// [linalg.algs.blas2.rank1.gerc],
// conjugated rank-1 matrix update
template<class in_vector_1_t,
  class in_vector_2_t,
  class inout_matrix_t>
void matrix_rank_1_update_c(
  in_vector_1_t x,
  in_vector_2_t y,
  inout_matrix_t A);

template<class ExecutionPolicy,
  class in_vector_1_t,
  class in_vector_2_t,
  class inout_matrix_t>
```
class in_vector_2_t,
class inout_matrix_t>
void matrix_rank_1_update_c(
    ExecutionPolicy&& exec,
    in_vector_1_t x,
    in_vector_2_t y,
    inout_matrix_t A);

// [linalg.algs.blas2.rank1.syr],
// symmetric rank-1 matrix update

template<class in_vector_t,
    class inout_matrix_t,
    class Triangle>
void symmetric_matrix_rank_1_update(
    in_vector_t x,
    inout_matrix_t A,
    Triangle t);

template<class ExecutionPolicy,
    class in_vector_t,
    class inout_matrix_t,
    class Triangle>
void symmetric_matrix_rank_1_update(
    ExecutionPolicy&& exec,
    in_vector_t x,
    inout_matrix_t A,
    Triangle t);

template<class T,
    class in_vector_t,
    class inout_matrix_t,
    class Triangle>
void symmetric_matrix_rank_1_update(
    T alpha,
    in_vector_t x,
    inout_matrix_t A,
    Triangle t);

template<class ExecutionPolicy,
    class T,
    class in_vector_t,
    class inout_matrix_t,
    class Triangle>
void symmetric_matrix_rank_1_update(
    ExecutionPolicy&& exec,
    T alpha,
    in_vector_t x,
    inout_matrix_t A,
    Triangle t);

// [linalg.algs.blas2.rank1.her],
// Hermitian rank-1 matrix update

template<class in_vector_t,
    class inout_matrix_t,
    class Triangle>
void hermitian_matrix_rank_1_update(
    in_vector_t x,
    inout_matrix_t A,
    Triangle t);

template<class ExecutionPolicy,
    class in_vector_t,
    class inout_matrix_t,
    class Triangle>
void hermitian_matrix_rank_1_update(
    ExecutionPolicy&& exec,
    in_vector_t x,
    inout_matrix_t A,
    Triangle t);
class inout_matrix_t,  
class Triangle>
void hermitian_matrix_rank_1_update(  
    ExecutionPolicy&& exec,  
    in_vector_t x,  
    inout_matrix_t A,  
    Triangle t);  

template<class T,  
    class in_vector_t,  
    class inout_matrix_t,  
    class Triangle>
void hermitian_matrix_rank_1_update(  
    T alpha,  
    in_vector_t x,  
    inout_matrix_t A,  
    Triangle t);  

template<class ExecutionPolicy,  
    class T,  
    class in_vector_t,  
    class inout_matrix_t,  
    class Triangle>
void hermitian_matrix_rank_1_update(  
    ExecutionPolicy&& exec,  
    T alpha,  
    in_vector_t x,  
    inout_matrix_t A,  
    Triangle t);

// [linalg.algs.blas2.rank2.syr2],  
// symmetric rank-2 matrix update  
template<class in_vector_1_t,  
    class in_vector_2_t,  
    class inout_matrix_t,  
    class Triangle>
void symmetric_matrix_rank_2_update(  
    in_vector_1_t x,  
    in_vector_2_t y,  
    inout_matrix_t A,  
    Triangle t);  

template<class ExecutionPolicy,  
    class in_vector_1_t,  
    class in_vector_2_t,  
    class inout_matrix_t,  
    class Triangle>
void symmetric_matrix_rank_2_update(  
    ExecutionPolicy&& exec,  
    in_vector_1_t x,  
    in_vector_2_t y,  
    inout_matrix_t A,  
    Triangle t);

// [linalg.algs.blas2.rank2.her2],  
// Hermitian rank-2 matrix update  
template<class in_vector_1_t,  
    class in_vector_2_t,  
    class inout_matrix_t,  
    class Triangle>
void hermitian_matrix_rank_2_update(  
    in_vector_1_t x,  
    in_vector_2_t y,  
    inout_matrix_t A,  
    Triangle t);
template<class ExecutionPolicy, class in_vector_1_t, class in_vector_2_t, class inout_matrix_t, class Triangle t>
void hermitian_matrix_rank_2_update(ExecutionPolicy&& exec, in_vector_1_t x, in_vector_2_t y, inout_matrix_t A, Triangle t);

// [linalg.algs.blas3.gemm],
// general matrix–matrix product
template<class in_matrix_1_t, class in_matrix_2_t, class out_matrix_t>
void matrix_product(in_matrix_1_t A, in_matrix_2_t B, out_matrix_t C);

template<class ExecutionPolicy, class in_matrix_1_t, class in_matrix_2_t, class out_matrix_t>
void matrix_product(ExecutionPolicy&& exec, in_matrix_1_t A, in_matrix_2_t B, out_matrix_t C);

template<class in_matrix_1_t, class in_matrix_2_t, class in_matrix_3_t, class out_matrix_t>
void matrix_product(in_matrix_1_t A, in_matrix_2_t B, in_matrix_3_t E, out_matrix_t C);

template<class ExecutionPolicy, class in_matrix_1_t, class in_matrix_2_t, class in_matrix_3_t, class out_matrix_t>
void matrix_product(ExecutionPolicy&& exec, in_matrix_1_t A, in_matrix_2_t B, in_matrix_3_t E, out_matrix_t C);

// [linalg.algs.blas3.symm],
// symmetric matrix–matrix product

// [linalg.algs.blas3.symm.ov.left],
// overwriting symmetric matrix–matrix left product
template<class in_matrix_1_t, class Triangle,
void symmetric_matrix_left_product(class in_matrix_1_t A, Triangle t, class in_matrix_2_t B, out_matrix_t C);

template<class ExecutionPolicy, class in_matrix_1_t, class Triangle, class in_matrix_2_t, class out_matrix_t>
void symmetric_matrix_left_product(ExecutionPolicy& exec, in_matrix_1_t A, Triangle t, in_matrix_2_t B, out_matrix_t C);

// [linalg.algs.blas3.symm.ov.right],
// overwriting symmetric matrix-matrix right product

void symmetric_matrix_right_product(class in_matrix_1_t A, Triangle t, class in_matrix_2_t B, out_matrix_t C);

template<class ExecutionPolicy, class in_matrix_1_t, class Triangle, class in_matrix_2_t, class out_matrix_t>
void symmetric_matrix_right_product(ExecutionPolicy& exec, in_matrix_1_t A, Triangle t, in_matrix_2_t B, out_matrix_t C);

// [linalg.algs.blas3.symm.up.left],
// updating symmetric matrix-matrix left product

void symmetric_matrix_left_product(class in_matrix_1_t A, Triangle t, class in_matrix_2_t B, class in_matrix_3_t E, out_matrix_t C);

template<class ExecutionPolicy, class in_matrix_1_t, class Triangle, class in_matrix_2_t, class in_matrix_3_t, class out_matrix_t>
void symmetric_matrix_left_product(ExecutionPolicy& exec, in_matrix_1_t A, Triangle t, in_matrix_2_t B, in_matrix_3_t E, out_matrix_t C);
class in_matrix_2_t,
class in_matrix_3_t,
class out_matrix_t>
void symmetric_matrix_left_product(
    ExecutionPolicy&& exec,
    in_matrix_1_t A,
    Triangle t,
    in_matrix_2_t B,
    in_matrix_3_t E,
    out_matrix_t C);

// [linalg.algs.blas3.symm.up.right],
// updating symmetric matrix-matrix right product
template<class in_matrix_1_t,
    class Triangle,
    class in_matrix_2_t,
    class in_matrix_3_t,
    class out_matrix_t>
void symmetric_matrix_right_product(
    in_matrix_1_t A,
    Triangle t,
    in_matrix_2_t B,
    in_matrix_3_t E,
    out_matrix_t C);

template<class ExecutionPolicy,
    class in_matrix_1_t,
    class Triangle,
    class in_matrix_2_t,
    class in_matrix_3_t,
    class out_matrix_t>
void symmetric_matrix_right_product(
    ExecutionPolicy&& exec,
    in_matrix_1_t A,
    Triangle t,
    in_matrix_2_t B,
    in_matrix_3_t E,
    out_matrix_t C);

// [linalg.algs.blas3.hemm],
// Hermitian matrix-matrix product

// [linalg.algs.blas3.hemm.ov.left],
// overwriting Hermitian matrix-matrix left product
template<class in_matrix_1_t,
    class Triangle,
    class in_matrix_2_t,
    class out_matrix_t>
void hermitian_matrix_left_product(
    in_matrix_1_t A,
    Triangle t,
    in_matrix_2_t B,
    out_matrix_t C);

template<class ExecutionPolicy,
    class in_matrix_1_t,
    class Triangle,
    class in_matrix_2_t,
    class out_matrix_t>
void hermitian_matrix_left_product(
    ExecutionPolicy&& exec,
    in_matrix_1_t A,
    Triangle t,
    in_matrix_2_t B,
    out_matrix_t C);
// [linalg.algs.blas3.hemm.ov.right],
// overwriting Hermitian matrix-matrix right product

template<class in_matrix_1_t,
        class Triangle,
        class in_matrix_2_t,
        class out_matrix_t>
void hermitian_matrix_right_product(
    in_matrix_1_t A,
    Triangle t,
    in_matrix_2_t B,
    out_matrix_t C);

template<class ExecutionPolicy,
        class in_matrix_1_t,
        class Triangle,
        class in_matrix_2_t,
        class out_matrix_t>
void hermitian_matrix_right_product(
    ExecutionPolicy&& exec,
    in_matrix_1_t A,
    Triangle t,
    in_matrix_2_t B,
    out_matrix_t C);

// [linalg.algs.blas3.hemm.up.left],
// updating Hermitian matrix-matrix left product

template<class in_matrix_1_t,
        class Triangle,
        class in_matrix_2_t,
        class in_matrix_3_t,
        class out_matrix_t>
void hermitian_matrix_left_product(
    in_matrix_1_t A,
    Triangle t,
    in_matrix_2_t B,
    in_matrix_3_t E,
    out_matrix_t C);

template<class ExecutionPolicy,
        class in_matrix_1_t,
        class Triangle,
        class in_matrix_2_t,
        class in_matrix_3_t,
        class out_matrix_t>
void hermitian_matrix_left_product(
    ExecutionPolicy&& exec,
    in_matrix_1_t A,
    Triangle t,
    in_matrix_2_t B,
    in_matrix_3_t E,
    out_matrix_t C);

// [linalg.algs.blas3.hemm.up.right],
// updating Hermitian matrix-matrix right product
template<
  class in_matrix_1_t,
  class Triangle,
  class in_matrix_2_t,
  class in_matrix_3_t,
  class out_matrix_t>
void hermitian_matrix_right_product(
  in_matrix_1_t A,
  Triangle t,
  in_matrix_2_t B,
  in_matrix_3_t E,
  out_matrix_t C);

// [linalg.algs.blas3.trmm],
// triangular matrix-matrix product

// [linalg.algs.blas3.trmm.ov.left],
// overwriting triangular matrix-matrix left product
template<
  class in_matrix_1_t,
  class Triangle,
  class DiagonalStorage,
  class in_matrix_2_t,
  class out_matrix_t>
void triangular_matrix_left_product(
  in_matrix_1_t A,
  Triangle t,
  DiagonalStorage d,
  in_matrix_2_t B,
  out_matrix_t C);

// [linalg.algs.blas3.trmm],
// triangular matrix-matrix product

// [linalg.algs.blas3.trmm.ov.left],
// overwriting triangular matrix-matrix left product
template<
  class in_matrix_1_t,
  class Triangle,
  class DiagonalStorage,
  class in_matrix_2_t,
  class out_matrix_t>
void triangular_matrix_left_product(
  in_matrix_1_t A,
  Triangle t,
  DiagonalStorage d,
  in_matrix_2_t B,
  out_matrix_t C);
template<
class in_matrix_1_t,
class Triangle,
class DiagonalStorage,
class in_matrix_2_t,
class out_matrix_t>
void triangular_matrix_left_product(
in_matrix_1_t A,
Triangle t,
DiagonalStorage d,
inout_matrix_t C);

// [linalg.algs.blas3.trmm.ov.right],
// overwriting triangular matrix–matrix right product

void triangular_matrix_right_product(
in_matrix_1_t A,
Triangle t,
DiagonalStorage d,
in_matrix_2_t B,
out_matrix_t C);

template<
class ExecutionPolicy,
class triangle_matrix_1_t,
class Triangle,
class DiagonalStorage,
class in_matrix_2_t,
class out_matrix_t>
void triangular_matrix_right_product(
ExecutionPolicy && exec,
in_matrix_1_t A,
Triangle t,
DiagonalStorage d,
in_matrix_2_t B,
out_matrix_t C);

template<
class inout_matrix_t>
void triangular_matrix_right_product(
in_matrix_1_t A,
Triangle t,
DiagonalStorage d,
inout_matrix_t C);

// [linalg.algs.blas3.trmm.up.left],
// updating triangular matrix–matrix left product

void triangular_matrix_left_product(
in_matrix_1_t A,
Triangle t,
DiagonalStorage d,
in_matrix_2_t B,
in_matrix_3_t E,
out_matrix_t C);

template<
class ExecutionPolicy>
void triangular_matrix_left_product(
in_matrix_1_t A,
Triangle t,
DiagonalStorage d,
in_matrix_2_t B,
in_matrix_3_t E,
out_matrix_t C);
class in_matrix_1_t,
class Triangle,
class DiagonalStorage,
class in_matrix_2_t,
class in_matrix_3_t,
class out_matrix_t>
void triangular_matrix_left_product(
    ExecutionPolicy&& exec,
    in_matrix_1_t A,
    Triangle t,
    DiagonalStorage d,
    in_matrix_2_t B,
    in_matrix_3_t E,
    out_matrix_t C);

// [linalg.algs.blas3.trmm.up.right],
// updating triangular matrix-matrix right product

template<class in_matrix_1_t,
class Triangle,
class DiagonalStorage,
class in_matrix_2_t,
class in_matrix_3_t,
class out_matrix_t>
void triangular_matrix_right_product(
    in_matrix_1_t A,
    Triangle t,
    DiagonalStorage d,
    in_matrix_2_t B,
    in_matrix_3_t E,
    out_matrix_t C);

template<class ExecutionPolicy,
class in_matrix_1_t,
class Triangle,
class DiagonalStorage,
class in_matrix_2_t,
class in_matrix_3_t,
class out_matrix_t>
void triangular_matrix_right_product(
    ExecutionPolicy&& exec,
    in_matrix_1_t A,
    Triangle t,
    DiagonalStorage d,
    in_matrix_2_t B,
    in_matrix_3_t E,
    out_matrix_t C);

// [linalg.algs.blas3.rank-k.syrk],
// rank-k symmetric matrix update

template<class in_matrix_1_t,
class inout_matrix_t,
class Triangle>
void symmetric_matrix_rank_k_update(
    in_matrix_1_t A,
    inout_matrix_t C,
    Triangle t);

template<class ExecutionPolicy,
class in_matrix_1_t,
class inout_matrix_t,
class Triangle>
void symmetric_matrix_rank_k_update(
    ExecutionPolicy&& exec,
    in_matrix_1_t A,
    inout_matrix_t C,
    Triangle t);

// [linalg.alg.blas3.rank-k.herk],
// rank-k Hermitian matrix update
void hermitian_matrix_rank_k_update(
    in_matrix_1_t A,
    inout_matrix_t C,
    Triangle t);

// [linalg.alg.blas3.rank-k.herk],
// rank-k Hermitian matrix update

class Triangle>
void hermitian_matrix_rank_k_update(
    ExecutionPolicy&& exec,
    T alpha,
    in_matrix_1_t A,
    inout_matrix_t C,
    Triangle t);

// [linalg.alg.blas3.rank2k.syr2k],
// rank-2k symmetric matrix update
template<class in_matrix_1_t,
    class in_matrix_2_t,
    class inout_matrix_t,
    class Triangle>
void symmetric_matrix_rank_2k_update(
    in_matrix_1_t A,
    in_matrix_2_t B,
    inout_matrix_t C,
    Triangle t);

template<class ExecutionPolicy,
    class in_matrix_1_t,
    class in_matrix_2_t,
    class inout_matrix_t,
    class Triangle>
void symmetric_matrix_rank_2k_update(
    ExecutionPolicy&& exec,
    in_matrix_1_t A,
    in_matrix_2_t B,
    inout_matrix_t C,
    Triangle t);

// [linalg.alg.blas3.rank2k.her2k],
// rank-2k Hermitian matrix update
template<class in_matrix_1_t,
    class in_matrix_2_t,
    class inout_matrix_t,
    class Triangle>
void hermitian_matrix_rank_2k_update(
    in_matrix_1_t A,
    in_matrix_2_t B,
    inout_matrix_t C,
    Triangle t);

template<class ExecutionPolicy,
    class in_matrix_1_t,
    class in_matrix_2_t,
    class inout_matrix_t,
    class Triangle>
void hermitian_matrix_rank_2k_update(
    ExecutionPolicy&& exec,
    in_matrix_1_t A,
    in_matrix_2_t B,
    inout_matrix_t C,
    Triangle t);

// [linalg.alg.blas3.trsm],
// solve multiple triangular linear systems
// [linalg.alg.blas3.trsm.left],
// solve multiple triangular linear systems
// with triangular matrix on the left

template<class in_matrix_1_t,
class Triangle,
class DiagonalStorage,
class in_matrix_2_t,
class out_matrix_t>
void triangular_matrix_matrix_left_solve(
in_matrix_1_t A,
Triangle t,
DiagonalStorage d,
in_matrix_2_t B,
out_matrix_t X);

template<class ExecutionPolicy,
class in_matrix_1_t,
class Triangle,
class DiagonalStorage,
class in_matrix_2_t,
class out_matrix_t>
void triangular_matrix_matrix_left_solve(
ExecutionPolicy&& exec,
in_matrix_1_t A,
Triangle t,
DiagonalStorage d,
in_matrix_2_t B,
out_matrix_t X);

template<class in_matrix_1_t,
class Triangle,
class DiagonalStorage,
class inout_matrix_t>
void triangular_matrix_matrix_left_solve(
in_matrix_1_t A,
Triangle t,
DiagonalStorage d,
inout_matrix_t B);

// [linalg.alg.blas3.trsm.right],
// solve multiple triangular linear systems
// with triangular matrix on the right

template<class in_matrix_1_t,
class Triangle,
class DiagonalStorage,
class in_matrix_2_t,
class out_matrix_t>
void triangular_matrix_matrix_right_solve(
in_matrix_1_t A,
Triangle t,
DiagonalStorage d,
in_matrix_2_t B,
out_matrix_t X);

template<class ExecutionPolicy,
class in_matrix_1_t,
class Triangle,
class DiagonalStorage,
class in_matrix_2_t,
class out_matrix_t>
void triangular_matrix_matrix_right_solve(
ExecutionPolicy&& exec,
in_matrix_1_t A,
Triangle t,
DiagonalStorage d,
in_matrix_2_t B,
out_matrix_t X);
in_matrix_t A,
Triangle t,
DiagonalStorage d,
in_matrix_t B,
out_matrix_t X);

template<class in_matrix_1_t,
class Triangle,
class DiagonalStorage,
class inout_matrix_t>
void triangular_matrix_matrix_right_solve(
in_matrix_1_t A,
Triangle t,
DiagonalStorage d,
inout_matrix_t B);

Tag classes [linalg.tags]

Storage order tags [linalg.tags.order]

struct column_major_t { };
inline constexpr column_major_t column_major = { };

struct row_major_t { };
inline constexpr row_major_t row_major = { };

column_major_t indicates a column-major order, and row_major_t indicates a row-major order. The interpretation of each depends on the specific layout that uses the tag. See layout_blas_general and layout_blas_packed below.

Triangle tags [linalg.tags.triangle]

Some linear algebra algorithms distinguish between the "upper triangle," "lower triangle," and "diagonal" of a matrix.

- The upper triangle of a matrix $A$ is the set of all elements of $A$ accessed by $A(i,j)$ with $i >= j$.
- The lower triangle of $A$ is the set of all elements of $A$ accessed by $A(i,j)$ with $i <= j$.
- The diagonal is the set of all elements of $A$ accessed by $A(i,i)$. It is included in both the upper triangle and the lower triangle.

struct upper_triangle_t { };
inline constexpr upper_triangle_t upper_triangle = { };

struct lower_triangle_t { };
inline constexpr lower_triangle_t lower_triangle = { };

These tag classes specify whether algorithms and other users of a matrix (represented as a basic_mdspan) should access the upper triangle (upper_triangular_t) or lower triangle
(lower_triangular_t) of the matrix. This is also subject to the restrictions of implicit_unit_diagonal_t if that tag is also applied; see below.

**Diagonal tags [linalg.tags.diagonal]**

```cpp
struct implicit_unit_diagonal_t { }; inline constexpr implicit_unit_diagonal_t implicit_unit_diagonal = { };

struct explicit_diagonal_t { }; inline constexpr explicit_diagonal_t explicit_diagonal = { };
```

These tag classes specify what algorithms and other users of a matrix should assume about the diagonal entries of the matrix, and whether algorithms and users of the matrix should access those diagonal entries explicitly.

The `implicit_unit_diagonal_t` tag indicates two things:

- the function will never access the \(i, i\) element of the matrix, and
- the matrix has a diagonal of ones (a "unit diagonal").

The tag `explicit_diagonal_t` indicates that algorithms and other users of the viewer may access the matrix's diagonal entries directly.

**Layouts for general and packed matrix types [linalg.layouts]**

**layout_blas_general [linalg.layouts.general]**

`layout_blas_general` is a `basic_mdspan` layout mapping policy. Its `StorageOrder` template parameter determines whether the matrix's data layout is column major or row major.

`layout_blas_general<column_major_t>` represents a column-major matrix layout, where the stride between consecutive rows is always one, and the stride between consecutive columns may be greater than or equal to the number of rows. *[Note: This is a generalization of layout_left. --end note]*

`layout_blas_general<row_major_t>` represents a row-major matrix layout, where the stride between consecutive rows may be greater than or equal to the number of columns, and the stride between consecutive columns is always one. *[Note: This is a generalization of layout_right. --end note]*

*[Note:]*

`layout_blas_general` represents exactly the data layout assumed by the General (GE) matrix type in the BLAS' C binding. It has two advantages:

1. Unlike `layout_left` and `layout_right`, any "submatrix" (subspan of consecutive rows and consecutive columns) of a matrix with `layout_blas_general<StorageOrder>` layout also has `layout_blas_general<StorageOrder>` layout.
2. Unlike layout_stride, it always has compile-time unit stride in one of the matrix's two extents.

BLAS functions call the possibly nonunit stride of the matrix the "leading dimension" of that matrix. For example, a BLAS function argument corresponding to the leading dimension of the matrix \( A \) is called LDA, for "leading dimension of the matrix \( A \)."

---end note---

```cpp
template<class StorageOrder>
class layout_blas_general {
public:
  template<class Extents>
  struct mapping {
    private:
      Extents extents_; // exposition only
      const typename Extents::index_type stride_; // exposition only

    public:
      constexpr mapping(const Extents& e, const typename Extents::index_type s);

      template<class OtherExtents>
      constexpr mapping(const mapping<OtherExtents>& e) noexcept;

      typename Extents::index_type
      operator()(typename Extents::index_type i, typename Extents::index_type j) const;

      constexpr typename Extents::index_type
      required_span_size() const noexcept;

      typename Extents::index_type
      stride(typename Extents::index_type r) const noexcept;

      template<class OtherExtents>
      bool operator==(const mapping<OtherExtents>& m) const noexcept;

      template<class OtherExtents>
      bool operator!=(const mapping<OtherExtents>& m) const noexcept;

      Extents extents() const noexcept;

      static constexpr bool is_always_unique();
      static constexpr bool is_always_contiguous();
      static constexpr bool is_always_strided();

      constexpr bool is_unique() const noexcept;
      constexpr bool is_contiguous() const noexcept;
      constexpr bool is_strided() const noexcept;
  };
};
```

- **Constraints:**
- StorageOrder is either column_major_t or row_major_t.
- Extents is a specialization of extents.
- Extents::rank() equals 2.

```cpp
constexpr mapping(const Extents& e, const typename Extents::index_type s);
```

- Requires:
  - If StorageOrder is column_major_t, then s is greater than or equal to e.extent(0). Otherwise, if StorageOrder is row_major_t, then s is greater than or equal to e.extent(1).

- Effects: Initializes extents_ with e, and initializes stride_ with s.

[Note:
The BLAS Standard requires that the stride be one if the corresponding matrix dimension is zero. We do not impose this requirement here, because it is specific to the BLAS. If an implementation dispatches to a BLAS function, then the implementation must impose the requirement at run time.

--end note]

```cpp
template<class OtherExtents>
constexpr mapping(const mapping<OtherExtents>& e) noexcept;
```

- Constraints:
  - OtherExtents is a specialization of extents.
  - OtherExtents::rank() equals 2.

- Effects: Initializes extents_ with m.extents_, and initializes stride_ with m.stride_.

```cpp
typename Extents::index_type
operator() (typename Extents::index_type i, 
typename Extents::index_type j) const;
```

- Requires:
  - 0 ≤ i < extent(0), and
  - 0 ≤ j < extent(1).

- Returns:
  - If StorageOrder is column_major_t, then i + stride(1)*j;
  - else, if StorageOrder is row_major_t, then stride(0)*i + j.
```cpp
template<class OtherExtents>
bool operator==(const mapping<OtherExtents>& m) const;

• **Constraints:** OtherExtents::rank() equals rank().
• **Returns:** true if and only if for 0 ≤ r < rank(), m.extent(r) equals extent(r) and m.stride(r) equals stride(r).

```cpp
template<class OtherExtents>
bool operator!=(const mapping<OtherExtents>& m) const;
```

• **Constraints:** OtherExtents::rank() equals rank().
• **Returns:** Returns: true if and only if there exists r with 0 ≤ r < rank() such that m.extent(r) does not equal extent(r) or m.stride(r) does not equal stride(r).

```cpp
typename Extents::index_type
stride(typename Extents::index_type r) const noexcept;
```

• **Returns:**
  - If StorageOrder is column_major_t, stride_if r equals 1, else 1;
  - else, if StorageOrder is row_major_t, stride_if r equals 0, else 1.

```cpp
constexpr typename Extents::index_type
required_span_size() const noexcept;
```

• **Returns:** stride(0)*stride(1).

```cpp
Extents extents() const noexcept;
```

• **Effects:** Equivalent to return extents_;.

```cpp
static constexpr bool is_always_unique();
```

• **Returns:** true.

```cpp
static constexpr bool is_always_contiguous();
```

• **Returns:** false.

```cpp
static constexpr bool is_always_strided();
```

• **Returns:** true.
\texttt{constexpr bool is_unique()} const noexcept;

- \textit{Returns:} \texttt{true}.

\texttt{constexpr bool is_contiguous()} const noexcept;

- \textit{Returns:}
  - If \texttt{StorageOrder} is \texttt{column_major_t}, then \texttt{true} if \texttt{stride(1)} equals \texttt{extent(0)}, else \texttt{false};
  - else, if \texttt{StorageOrder} is \texttt{row_major_t}, then \texttt{true} if \texttt{stride(0)} equals \texttt{extent(1)}, else \texttt{false}.

\texttt{constexpr bool is_strided()} const noexcept;

- \textit{Returns:} \texttt{true}.

\textbf{layout\_blas\_packed} \texttt{[linalg.layouts.packed]}

\texttt{layout\_blas\_packed} is a \texttt{basic\_mdspan} layout mapping policy that represents a square matrix that stores only the entries in one triangle, in a packed contiguous format. Its \texttt{Triangle} template parameter determines whether an \texttt{basic\_mdspan} with this layout stores the upper or lower triangle of the matrix. Its \texttt{StorageOrder} template parameter determines whether the layout packs the matrix's elements in column-major or row-major order.

A \texttt{StorageOrder} of \texttt{column\_major\_t} indicates column-major ordering. This packs matrix elements starting with the leftmost (least column index) column, and proceeding column by column, from the top entry (least row index).

A \texttt{StorageOrder} of \texttt{row\_major\_t} indicates row-major ordering. This packs matrix elements starting with the topmost (least row index) row, and proceeding row by row, from the leftmost (least column index) entry.

\textit{[Note:]}

\texttt{layout\_blas\_packed} describes the data layout used by the BLAS' Symmetric Packed (SP), Hermitian Packed (HP), and Triangular Packed (TP) matrix types.

If \texttt{transpose\_view}'s input has layout \texttt{layout\_blas\_packed}, the return type also has layout \texttt{layout\_blas\_packed}, but with opposite \texttt{Triangle} and \texttt{StorageOrder}. For example, the transpose of a packed column-major upper triangle, is a packed row-major lower triangle.

--\textit{end note}--

\texttt{template<class Triangle, class StorageOrder>}
\begin{verbatim}
class layout_blas_packed {
public:
\end{verbatim}
```cpp
template<class Extents>
struct mapping {
private:
    Extents extents_; // exposition only

public:
    constexpr mapping(const Extents & e);

template<class OtherExtents>
constexpr mapping(const mapping<OtherExtents> & e) noexcept;

typename Extents::index_type
operator()(typename Extents::index_type i, typename Extents::index_type j) const;

template<class OtherExtents>
bool operator==(const mapping<OtherExtents> & m) const noexcept;

template<class OtherExtents>
bool operator!=(const mapping<OtherExtents> & m) const noexcept;

typename Extents::index_type
stride(typename Extents::index_type r) const noexcept;

typename Extents::index_type
required_span_size() const noexcept;

typename Extents::index_type
extents() const noexcept;

static constexpr bool is_always_unique();
static constexpr bool is_always_contiguous();
static constexpr bool is_always_strided();

costexpr bool is_unique() const noexcept;
costexpr bool is_contiguous() const noexcept;
costexpr bool is_strided() const noexcept;
};
```

- **Constraints:**
  - Triangle is either `upper_triangle_t` or `lower_triangle_t`.
  - StorageOrder is either `column_major_t` or `row_major_t`.
  - Extents is a specialization of extents.
  - Extents::rank() equals 2.

```cpp
constexpr mapping(const Extents & e);
```

- **Requires:** `e.extent(0)` equals `e.extent(1)`.
- **Effects:** Initializes extents_ with e.
```cpp
template<class OtherExtents>
const expr mapping(const mapping<OtherExtents>& e);
```

- **Constraints:**
  - OtherExtents is a specialization of extents.
  - OtherExtents::rank() equals 2.
- **Effects:** Initializes extents_ with e.

```cpp
typename Extents::index_type
operator() (typename Extents::index_type i,
            typename Extents::index_type j) const;
```

- **Requires:**
  - \(0 \leq i < \text{extent}(0)\), and
  - \(0 \leq j < \text{extent}(1)\).
- **Returns:** Let \(N\) equal \(\text{extent}(0)\). Then:
  - If StorageOrder is column_major_t and
    - if Triangle is upper_triangle_t, then \(i + j(j+1)/2\) if \(i \geq j\), else \(j + i(i+1)/2\);
    - else, if Triangle is lower_triangle_t, then \(i + Nj - j(j+1)/2\) if \(i \leq j\), else \(j + Ni - i(i+1)/2\);
  - else, if StorageOrder is row_major_t and
    - if Triangle is upper_triangle_t, then \(j + Ni - i(i+1)/2\) if \(j \leq i\), else \(i + Nj - j(j+1)/2\);
    - else, if Triangle is lower_triangle_t, then \(j + i(i+1)/2\) if \(j \geq i\), else \(i + j(j+1)/2\).

```cpp
template<class OtherExtents>
bool operator==(const mapping<OtherExtents>& m) const;
```

- **Constraints:** OtherExtents::rank() equals rank().
- **Returns:** \(\text{true}\) if and only if for \(0 \leq r < \text{rank()}\), \(m.\text{extent}(r)\) equals extent(r).

```cpp
template<class OtherExtents>
bool operator!=(const mapping<OtherExtents>& m) const;
```

- **Constraints:** OtherExtents::rank() equals rank().
Returns: true if and only if there exists \( r \) with \( 0 \leq r < \text{rank()} \) such that \( m.\text{extent}(r) \) does not equal \( \text{extent}(r) \).

```cpp
constexpr typename Extents::index_type stride(typename Extents::index_type r) const noexcept;
```

Returns: 1 if \( \text{extent}(0) \) is less than 2, else 0.

```cpp
constexpr typename Extents::index_type required_span_size() const noexcept;
```

Returns: \( \text{extent}(0) \times (\text{extent}(0) - 1)/2 \).

```cpp
constexpr Extents extents() const noexcept;
```

Effects: Equivalent to \( \text{return extents;} \).

```cpp
static constexpr bool is_always_unique();
```

Returns: false.

```cpp
static constexpr bool is_always_contiguous();
```

Returns: true.

```cpp
static constexpr bool is_always_strided();
```

Returns: false.

```cpp
constexpr bool is_unique() const noexcept;
```

Returns: true if \( \text{extent}(0) \) is less than 2, else false.

```cpp
constexpr bool is_contiguous() const noexcept;
```

Returns: true.

```cpp
constexpr bool is_strided() const noexcept;
```

Returns: true if \( \text{extent}(0) \) is less than 2, else false.

Scaled in-place transformation [linalg.scaled]
The `scaled_view` function takes a value `alpha` and a `basic_mdspan x`, and returns a new read-only `basic_mdspan` with the same domain as `x`, that represents the elementwise product of `alpha` with each element of `x`.

[Example:

```cpp
// z = alpha * x + y
void z_equals_alpha_times_x_plus_y(
    mdspan<double, extents<dynamic_extent>> z,
    const double alpha,
    mdspan<double, extents<dynamic_extent>> x,
    mdspan<double, extents<dynamic_extent>> y
) {
    linalg_add(scaled_view(alpha, x), y, y);
}

// w = alpha * x + beta * y
void w_equals_alpha_times_x_plus_beta_times_y(
    mdspan<double, extents<dynamic_extent>> w,
    const double alpha,
    mdspan<double, extents<dynamic_extent>> x,
    const double beta,
    mdspan<double, extents<dynamic_extent>> y
) {
    linalg_add(scaled_view(alpha, x), scaled_view(beta, y), w);
}
```

--end example]

[Note:

An implementation could dispatch to a function in the BLAS library, by noticing that the first argument has an `accessor_scaled` `Accessor` type. It could use this information to extract the appropriate run-time value(s) of the relevant BLAS function arguments (e.g., `ALPHA` and/or `BETA`), by calling `accessor_scaled::scaling_factor`.

--end note]

**Class template accessor_scaled** [linalg.scaled.accessor_scaled]

The class template `accessor_scaled` is a `basic_mdspan` accessor policy whose reference type represents the product of a fixed value (the "scaling factor") and its nested `basic_mdspan` accessor's reference. It is part of the implementation of `scaled_view`.

The exposition-only class template `scaled_scalar` represents a read-only value, which is the product of a fixed value (the "scaling factor") and the value of a reference to an element of a `basic_mdspan`. [Note: The value is read only to avoid confusion with the definition of "assigning to a scaled scalar." --end note] `scaled_scalar` is part of the implementation of `scaled_accessor`.

```cpp
template<class ScalingFactor, 
    class Reference>
```
class scaled_scalar { // exposition only
private:
    const ScalingFactor scaling_factor;
    Reference value;
    using result_type = decltype(scaling_factor * value);
public:
    scaled_scalar(const ScalingFactor& s, Reference v);
    operator result_type() const;
};

- Requires:
  - ScalingFactor and Reference shall be Cpp17CopyConstructible.

- Constraints:
  - The expression scaling_factor * value is well formed.

scaled_scalar(const ScalingFactor& s, Reference v);

- Effects: Initializes scaling_factor with s, and initializes value with v.

operator result_type() const;

- Effects: Equivalent to return scaling_factor * value;

The class template accessor_scaled is a basic_mdspan accessor policy whose reference type represents the product of a scaling factor and its nested basic_mdspan accessor's reference.

template<class ScalingFactor, class Accessor>
class accessor_scaled {
public:
    using element_type  = Accessor::element_type;
    using pointer       = Accessor::pointer;
    using reference     = scaled_scalar<ScalingFactor, Accessor::reference>;
    using offset_policy = accessor_scaled<ScalingFactor, Accessor::offset_policy>;

    accessor_scaled(const ScalingFactor& s, Accessor a);

    reference access(pointer p, ptrdiff_t i) const noexcept;

    offset_policy::pointer
    offset(pointer p, ptrdiff_t i) const noexcept;

    element_type* decay(pointer p) const noexcept;
}
ScalingFactor scaling_factor() const;

private:
    const ScalingFactor scaling_factor_; // exposition only
    Accessor accessor; // exposition only
};

Requires:

- *ScalingFactor* and *Accessor* shall be *Cpp17CopyConstructible*.
- *Accessor* shall meet the *basic_mdspan* accessor policy requirements (see [mdspan.accessor.reqs] in P0009).

accessor_scaled(const ScalingFactor& s, Accessor a);

Effects: Initializes *scaling_factor_* with *s*, and initializes *accessor* with *a*.

reference access(pointer p, ptrdiff_t i) const noexcept;

Effects: Equivalent to return *accessor.access(p, i)*;

offset_policy::pointer offset(pointer p, ptrdiff_t i) const noexcept;

Effects: Equivalent to return *accessor.offset(p, i)*;

element_type* decay(pointer p) const noexcept;

Effects: Equivalent to return *accessor.decay(p)*;

ScalingFactor scaling_factor() const;

Effects: Equivalent to return *scaling_factor_*;

**scaled_view** [linalg.scaled.scaled_view]

The *scaled_view* function takes a value *alpha* and a *basic_mdspan* *x*, and returns a new read-only *basic_mdspan* with the same domain as *x*, that represents the elementwise product of *alpha* with each element of *x*.

```cpp
template<class ScalingFactor, 
    class ElementType, 
    class Extents, 
    class Layout, 
    class Accessor>
```
Let R name the type basic_mdspan<ReturnType, Extents, Layout, ReturnAccessor>, where

- ReturnType is either ElementType or const ElementType; and
- ReturnAccessor is:
  - if Accessor is accessor_scaled<NestedScalingFactor, NestedAccessor> for some NestedScalingFactor and NestedAccessor, then either accessor_scaled<ProductScalingFactor, NestedAccessor> or accessor_scaled<ScalingFactor, Accessor>, where ProductScalingFactor is decltype(s * a.accessor().scaling_factor());
  - else, accessor_scaled<ScalingFactor, Accessor>.

- Effects:
  - If Accessor is accessor_scaled<NestedScalingFactor, NestedAccessor> and ReturnAccessor is accessor_scaled<ProductScalingFactor, NestedAccessor>, then equivalent to return R(a.data(), a.mapping(), ReturnAccessor(product_s, a.accessor().nested_accessor())); where product_s equals s * a.accessor().scaling_factor();
  - else, equivalent to return R(a.data(), a.mapping(), ReturnAccessor(s, a.accessor()));.

- Remarks: The elements of the returned basic_mdspan are read only.

[Note:

The point of ReturnAccessor is to give implementations freedom to optimize applying accessor_scaled twice in a row. However, implementations are not required to optimize arbitrary combinations of nested accessor_scaled interspersed with other nested accessors.

The point of ReturnType is that, based on P0009R9, it may not be possible to deduce the const version of Accessor for use in accessor_scaled. In general, it may not be correct or efficient to use an Accessor meant for a nonconst ElementType, with const ElementType. This is because Accessor::reference may be a type other than ElementType&. Thus, we cannot require that the return type have const ElementType as its element type, since that might not be compatible with the given Accessor. However, in some cases, like accessor_basic, it is possible to deduce the const version of Accessor. Regardless, users are not allowed to modify the elements of the returned basic_mdspan.

--end note]
void test_scaled_view(basic_mdspan<double, extents<10>> a) 
{
    auto a_scaled = scaled_view(5.0, a);
    for(int i = 0; i < a.extent(0); ++i) {
        assert(a_scaled(i) == 5.0 * a(i));
    }
}

--end example]

Conjugated in-place transformation [linalg.conj]

The conjugate_view function takes a basic_mdspan x, and returns a new read-only basic_mdspan y with the same domain as x, whose elements are the complex conjugates of the corresponding elements of x. If the element type of x is not complex<R> for some R, then y is a read-only view of the elements of x.

[Note:

An implementation could dispatch to a function in the BLAS library, by noticing that the Accessor type of a basic_mdspan input has type accessor_conjugate, and that its nested Accessor type is compatible with the BLAS library. If so, it could set the corresponding TRANS* BLAS function argument accordingly and call the BLAS function.

--end note]

Class template accessor_conjugate [linalg.conj.accessor_conjugate]

The class template accessor_conjugate is a basic_mdspan accessor policy whose reference type represents the complex conjugate of its nested basic_mdspan accessor's reference.

The exposition-only class template conjugated_scalar represents a read-only value, which is the complex conjugate of the value of a reference to an element of a basic_mdspan. [Note: The value is read only to avoid confusion with the definition of "assigning to the conjugate of a scalar." --end note] conjugated_scalar is part of the implementation of accessor_conjugate.

template<class Reference, 
    class ElementType>
class conjugated_scalar { // exposition only
public:
    conjugated_scalar(Reference v);

    operator ElementType() const;

private:
    Reference val;
};
- Requires: `Reference` shall be `Cpp17CopyConstructible`.
- Constraints:
  - The expression `conj(val)` is well formed and is convertible to `ElementType`.  
    
    [Note: This implies that `ElementType` is `complex<R>` for some type `R`. --end note]
If `element_type` is `complex<R>` for some `R`, then this names `conjugated_scalar<typename Accessor::reference, element_type>`. Otherwise, it names `typename Accessor::reference`.

```cpp
using offset_policy = /* see below */;
```

If `element_type` is `complex<R>` for some `R`, then this names `accessor_conjugate<typename Accessor::offset_policy, element_type>`. Otherwise, it names `typename Accessor::offset_policy>.

```cpp
accessor_conjugate(Accessor a);
```

- **Effects:** Initializes `acc` with `a`.

```cpp
reference access(pointer p, ptrdiff_t i) const
    noexcept(noexcept(acc.access(p, i)));
```

- **Effects:** Equivalent to `return reference(acc.access(p, i));`.

```cpp
typename offset_policy::pointer
offset(pointer p, ptrdiff_t i) const
    noexcept(noexcept(acc.offset(p, i)));
```

- **Effects:** Equivalent to `return acc.offset(p, i);`.

```cpp
element_type* decay(pointer p) const
    noexcept(noexcept(acc.decay(p)));
```

- **Effects:** Equivalent to `return acc.decay(p);`.

```cpp
Accessor nested_accessor() const;
```

- **Effects:** Equivalent to `return acc;`.

**conjugate_view** [linalg.conj.conjugate_view]

```cpp
template<class ElementType, class Extents, class Layout, class Accessor>
/* see below */
conjugate_view(
    basic_mdspan<ElementType, Extents, Layout, Accessor> a);
```
Let $R$ name the type $\text{basic_mdspan<}\text{ReturnElementType, Extents, Layout, ReturnAccessor}>$, where

- $\text{ReturnElementType}$ is either $\text{ElementType}$ or $\text{const ElementType}$; and

- $\text{ReturnAccessor}$ is:
  - if $\text{Accessor}$ is $\text{accessor_conjugate<NestedAccessor>}$ for some $\text{NestedAccessor}$, then either $\text{NestedAccessor}$ or $\text{accessor_conjugate<Accessor>}$,
  - else if $\text{ElementType}$ is $\text{complex<U>}$ or $\text{const complex<U>}$ for some $U$, then $\text{accessor_conjugate<Accessor>}$,
  - else either $\text{accessor_conjugate<Accessor>}$ or $\text{Accessor}$.

- **Effects:**
  - If $\text{Accessor}$ is $\text{accessor_conjugate<NestedAccessor>}$ and $\text{ReturnAccessor}$ is $\text{NestedAccessor}$, then equivalent to return $R(a.data(), a.mapping(), a.nested_accessor())$;
  - else, if $\text{ReturnAccessor}$ is $\text{accessor_conjugate<Accessor>}$, then equivalent to return $R(a.data(), a.mapping(), accessor_conjugate<Accessor>(a.accessor()))$;
  - else, equivalent to return $R(a.data(), a.mapping(), a.accessor())$.

- **Remarks:** The elements of the returned $\text{basic_mdspan}$ are read only.

[Note:

The point of $\text{ReturnAccessor}$ is to give implementations freedom to optimize applying $\text{accessor_conjugate}$ twice in a row. However, implementations are not required to optimize arbitrary combinations of nested $\text{accessor_conjugate}$ interspersed with other nested accessors.

--end note]

[Example:

```cpp
void test_conjugate_view_complex(
    basic_mdspan<complex<double>, extents<10>> a)
{
    auto a_conj = conjugate_view(a);
    for(int i = 0; i < a.extent(0); ++i) {
        assert(a_conj(i) == conj(a(i));
    }

    auto a_conj_conj = conjugate_view(a_conj);
    for(int i = 0; i < a.extent(0); ++i) {
        assert(a_conj_conj(i) == a(i));
    }
}
```
void test_conjugate_view_real(
    basic_mdspan<double, extents<10>> a)
{
    auto a_conj = conjugate_view(a);
    for(int i = 0; i < a.extent(0); ++i) {
        assert(a_conj(i) == a(i));
    }
    auto a_conj_conj = conjugate_view(a_conj);
    for(int i = 0; i < a.extent(0); ++i) {
        assert(a_conj_conj(i) == a(i));
    }
}

--end example]

Transpose in-place transformation [linalg.transp]

layout_transpose is a basic_mdspan layout mapping policy that swaps the rightmost two
indices, extents, and strides (if applicable) of any unique basic_mdspan layout mapping
policy.

The transpose_view function takes a rank-2 basic_mdspan representing a matrix, and
returns a new read-only basic_mdspan representing the transpose of the input matrix.

[Note:

An implementation could dispatch to a function in the BLAS library, by noticing that the first
argument has a layout_transpose Layout type, and/or an accessor_conjugate (see
below) Accessor type. It could use this information to extract the appropriate run-time
value(s) of the relevant TRANS* BLAS function arguments.

--end note]

layout_transpose [linalg.transp.layout_transpose]

layout_transpose is a basic_mdspan layout mapping policy that swaps the rightmost two
indices, extents, and strides (if applicable) of any unique basic_mdspan layout mapping
policy.

For InputExtents a specialization of extents, transpose_extents_t<InputExtents>
names the extents type OutputExtents such that

- InputExtents::static_extent(InputExtents::rank()-1) equals
  OutputExtents::static_extent(OutputExtents::rank()-2),

- InputExtents::static_extent(InputExtents::rank()-2) equals
  OutputExtents::static_extent(OutputExtents::rank()-1), and
• **InputExtents::static_extent(r)** equals **OutputExtents::static_extent(r)** for \(0 \leq r < \text{InputExtents::rank()-2}\).

• **Requires**: **InputExtents** is a specialization of **extents**.

• **Constraints**: **InputExtents::rank()** is at least 2.

```cpp
template<class InputExtents> 
transpose_extents_t<InputExtents> 
transpose_extents(const InputExtents in); // exposition only
```

• **Constraints**: **InputExtents::rank()** is at least 2.

• **Returns**: An **extents** object **out** such that

  - **out.extent(in.rank()-1)** equals **in.extent(in.rank()-2)**,
  - **out.extent(in.rank()-2)** equals **in.extent(in.rank()-1)**, and
  - **out.extent(r)** equals **in.extent(r)** for \(0 \leq r < \text{in.rank()-2}\).

```cpp
template<class Layout> 
class layout_transpose { 
public: 
    template<class Extents> 
    struct mapping { 
        private: 
            using nested_mapping_type = 
                typename Layout::template mapping< 
                    transpose_extents_t<Extents>>; // exposition only
                nested_mapping_type nested_mapping_; // exposition only
        public: 
            mapping(const nested_mapping_type& map); 
            ptdiff_t operator() (ptdiff_t i, ptdiff_t j) const 
                noexcept(noexcept(nested_mapping_(j, i))); 
            nested_mapping_type nested_mapping() const; 
        template<class OtherExtents> 
            bool operator==(const mapping<OtherExtents>& m) const; 
        template<class OtherExtents> 
            bool operator!=(const mapping<OtherExtents>& m) const; 
        Extents extents() const noexcept;
        typename Extents::index_type required_span_size() const 
            noexcept(noexcept(nested_mapping_.required_span_size()));
        bool is_unique() const 
                noexcept(noexcept(nested_mapping_.is_unique()));
        bool is_contiguous() const
    }
}
```
noexcept(noexcept(nested_mapping_.is_contiguous()));

bool is_strided() const
noexcept(noexcept(nested_mapping_.is_strided()));

static constexpr bool is_always_unique();
static constexpr bool is_always_contiguous();
static constexpr bool is_always_strided();

typename Extents::index_type
stride(typename Extents::index_type r) const
noexcept(noexcept(nested_mapping_.stride(r)));
};

- **Requires:**
  - Layout shall meet the basic_mdspan layout mapping policy requirements. [Note: See [mdspan.layout.reqs] in P0009R9. --end note]

- **Constraints:**
  - For all specializations E of extents with E::rank() equal to 2, typename Layout::template mapping<E>::is_always_unique() is true.

mapping(const nested_mapping_type& map);

- **Effects:** Initializes nested_mapping_ with map.

ptrdiff_t operator()(ptrdiff_t i, ptrdiff_t j) const
noexcept(noexcept(nested_mapping_(j, i)));

- **Effects:** Equivalent to return nested_mapping_(j, i);

nested_mapping_type nested_mapping() const;

- **Effects:** Equivalent to return nested_mapping_;

template<class OtherExtents>
bool operator==(const mapping<OtherExtents>& m) const;

- **Constraints:** OtherExtents::rank() equals rank().
- **Effects:** Equivalent to nested_mapping_ == m.nested_mapping_;

template<class OtherExtents>
bool operator!=(const mapping<OtherExtents>& m) const;
- **Constraints:** OtherExtents::rank() equals rank().

- **Effects:** Equivalent to `nested_mapping_ != m.nested_mapping_`;.

```cpp
Extents extents() const noexcept;
```

- **Effects:** Equivalent to return
  ```cpp
transpose_extents(nested_mapping_.extents());
  ```

```cpp
typename Extents::index_type
required_span_size() const
  noexcept(noexcept(nested_mapping_.required_span_size()));
```

- **Effects:** Equivalent to `return nested_mapping_.required_span_size();`;

```cpp
bool is_unique() const
  noexcept(noexcept(nested_mapping_.is_unique()));
```

- **Effects:** Equivalent to `return nested_mapping_.is_unique();`;

```cpp
bool is_contiguous() const
  noexcept(noexcept(nested_mapping_.is_contiguous()));
```

- **Effects:** Equivalent to `return nested_mapping_.is_contiguous();`;

```cpp
bool is_strided() const
  noexcept(noexcept(nested_mapping_.is_strided()));
```

- **Effects:** Equivalent to `return nested_mapping_.is_strided();`;

```cpp
static constexpr bool is_always_unique();
```

- **Effects:** Equivalent to `return nested_mapping_type::is_always_unique();`;

```cpp
static constexpr bool is_always_contiguous();
```

- **Effects:** Equivalent to `return nested_mapping_type::is_always_contiguous();`;

```cpp
static constexpr bool is_always_strided();
```

- **Effects:** Equivalent to `return nested_mapping_type::is_always_strided();`;

```cpp
typename Extents::index_type
stride(typename Extents::index_type r) const
```
```c++
#include <iostream>

int main() {
    // Example code
    return 0;
}
```

- **Constraints**: `is_always_strided()` is true.

- **Effects**: Equivalent to return `nested_mapping_.stride(s);`, where `s` is 0 if `r` is 1 and `s` is 1 if `r` is 0.

**transpose_view** [linalg.transpose_view]

The `transpose_view` function takes a rank-2 `basic_mdspan` representing a matrix, and returns a new read-only `basic_mdspan` representing the transpose of the input matrix. The input matrix's data are not modified, and the returned `basic_mdspan` accesses the input matrix's data in place. If the input `basic_mdspan`'s layout is already `layout_transpose<L>` for some layout `L`, then the returned `basic_mdspan` has layout `L`. Otherwise, the returned `basic_mdspan` has layout `layout_transpose<L>`, where `L` is the input `basic_mdspan`'s layout.

```c++
template<class ElementType,  
         class Extents,  
         class Layout,  
         class Accessor>

/* see-below */
transpose_view(  
    basic_mdspan<ElementType, Extents, Layout, Accessor> a);
```

Let `ReturnExtents` name the type `transpose_extents_t<Extents>`. Let `R` name the type `basic_mdspan<ReturnElementType, ReturnExtents, ReturnLayout, Accessor>`, where

- `ReturnElementType` is either `ElementType` or `const ElementType`; and

- `ReturnLayout` is:
  - if `Layout` is `layout_blas_packed<Triangle, StorageOrder>`, then
    `layout_blas_packed<OppositeTriangle, OppositeStorageOrder>`, where
      - `OppositeTriangle` names the type
        `conditional_t<is_same_v<Triangle, upper_triangle_t>,  
        lower_triangle_t, upper_triangle_t>`, and
      - `OppositeStorageOrder` names the type
        `conditional_t<is_same_v<StorageOrder, column_major_t>,  
        row_major_t, column_major_t>```
  - else, if `Layout` is `layout_transpose<NestedLayout>` for some `NestedLayout`, then either `NestedLayout` or `layout_transpose<Layout>`,
  - else `layout_transpose<Layout>`.

- **Effects:**
- If `Layout` is `layout_blas_packed<Triangle, StorageOrder>`, then equivalent to:
  ```cpp
to return R(a.data(),
  ReturnMapping(a.mapping().extents()), a.accessor());
```

- else, if `Layout` is `layout_transpose<NestedLayout>` and `ReturnLayout` is `NestedLayout`, then equivalent to:
  ```cpp
to return R(a.data(),
  a.mapping().nested_mapping(), a.accessor());
```

- else, equivalent to:
  ```cpp
to return R(a.data(), ReturnMapping(a.mapping()),
  a.accessor);
```

  where `ReturnMapping` names the type `typename layout_transpose<Layout>::template mapping<ReturnExtents>`.  

- **Remarks:** The elements of the returned `basic_mdspan` are read only.

**Note:**

Implementations may optimize applying `layout_transpose` twice in a row. However, implementations need not optimize arbitrary combinations of nested `layout_transpose` interspersed with other nested layouts.

---end note

**Example:**

```cpp
void test_transpose_view(basic_mdspan<double, extents<3, 4>> a)
{
    const ptrdiff_t num_rows = a.extent(0);
    const ptrdiff_t num_cols = a.extent(1);

    auto a_t = transpose_view(a);
    assert(num_rows == a_t.extent(1));
    assert(num_cols == a_t.extent(0));
    assert(a.stride(0) == a_t.stride(1));
    assert(a.stride(1) == a_t.stride(0));

    for(ptrdiff_t row = 0; row < num_rows; ++row) {
        for(ptrdiff_t col = 0; col < num_rows; ++col) {
            assert(a(row, col) == a_t(col, row));
        }
    }

    auto a_t_t = transpose_view(a_t);
    assert(num_rows == a_t_t.extent(0));
    assert(num_cols == a_t_t.extent(1));
    assert(a.stride(0) == a_t_t.stride(0));
    assert(a.stride(1) == a_t_t.stride(1));

    for(ptrdiff_t row = 0; row < num_rows; ++row) {
        for(ptrdiff_t col = 0; col < num_rows; ++col) {
            assert(a(row, col) == a_t_t(row, col));
        }
    }
}
```
Conjugate transpose transform [linalg.conj_transp]

The `conjugateTransposeView` function returns a conjugate transpose view of an object. This combines the effects of `transposeView` and `conjugateView`.

```cpp
template<
  class ElementType,
  class Extents,
  class Layout,
  class Accessor>
/* see-below */
conjugateTransposeView(
  basic_mdspan<ElementType, Extents, Layout, Accessor> a);
```

- **Effects**: Equivalent to `return conjugateView(transposeView(a));`.
- **Remarks**: The elements of the returned `basic_mdspan` are read only.

[Example:

```cpp
void test_ct_view(basic_mdspan<complex<double>, extents<3, 4>> a) {
    const ptrdiff_t num_rows = a.extent(0);
    const ptrdiff_t num_cols = a.extent(1);

    auto a_ct = conjugateTransposeView(a);
    assert(num_rows == a_ct.extent(1));
    assert(num_cols == a_ct.extent(0));
    assert(a.stride(0) == a_ct.stride(1));
    assert(a.stride(1) == a_ct.stride(0));

    for(ptrdiff_t row = 0; row < num_rows; ++row) {
        for(ptrdiff_t col = 0; col < num_rows; ++col) {
            assert(a(row, col) == conj(a_ct(col, row)));
        }
    }

    auto a_ct_ct = conjugateTransposeView(a_ct);
    assert(num_rows == a_ct_ct.extent(0));
    assert(num_cols == a_ct_ct.extent(1));
    assert(a.stride(0) == a_ct_ct.stride(0));
    assert(a.stride(1) == a_ct_ct.stride(1));

    for(ptrdiff_t row = 0; row < num_rows; ++row) {
        for(ptrdiff_t col = 0; col < num_rows; ++col) {
            assert(a(row, col) == a_ct_ct(row, col));
            assert(conj(a_ct(col, row)) == a_ct_ct(row, col));
        }
    }
}
```

--end example]
Requirements [linalg.algs.reqs]

Throughout this Clause, where the template parameters are not constrained, the names of
template parameters are used to express type requirements. In the requirements below, we use
* in a typename to denote a "wildcard," that matches zero characters, _1, _2, _3, or other
things as appropriate.

- Algorithms that have a template parameter named `ExecutionPolicy` are parallel
  algorithms [algorithms.parallel.defns].
- `Scalar` meets the requirements of `SemiRegular<Scalar>`. (Some algorithms below
  impose further requirements.)
- `Real` is any of the following types: `float`, `double`, or `long double`.
- `in_vector*_t` is a rank-1 `basic_mdspan` with a potentially `const` element type and a
  unique layout. If the algorithm accesses the object, it will do so in read-only fashion.
- `inout_vector*_t` is a rank-1 `basic_mdspan` with a non-`const` element type and a
  unique layout.
- `out_vector*_t` is a rank-1 `basic_mdspan` with a non-`const` element type and a
  unique layout. If the algorithm accesses the object, it will do so in write-only fashion.
- `in_matrix*_t` is a rank-2 `basic_mdspan` with a `const` element type. If the algorithm
  accesses the object, it will do so in read-only fashion.
- `inout_matrix*_t` is a rank-2 `basic_mdspan` with a non-`const` element type.
- `out_matrix*_t` is a rank-2 `basic_mdspan` with a non-`const` element type. If the
  algorithm accesses the object, it will do so in write-only fashion.
- `in_object*_t` is a rank-1 or rank-2 `basic_mdspan` with a potentially `const` element
  type and a unique layout. If the algorithm accesses the object, it will do so in read-only
  fashion.
- `inout_object*_t` is a rank-1 or rank-2 `basic_mdspan` with a non-`const` element type
  and a unique layout.
- `out_object*_t` is a rank-1 or rank-2 `basic_mdspan` with a non-`const` element type
  and a unique layout.
- `Triangle` is either `upper_triangle_t` or `lower_triangle_t`.
- `DiagonalStorage` is either `implicit_unit_diagonal_t` or `explicit_diagonal_t`.
- `in_*_t` template parameters may deduce a `const` lvalue reference or a (non-`const`)
rvalue reference to a `basic_mdspan`.
- `inout_*_t` and `out_*_t` template parameters may deduce a `const` lvalue reference to
  a `basic_mdspan`, or a (non-`const`) rvalue reference to a `basic_mdspan`. 
BLAS 1 functions [linalg.algs.blas1]

[Note:

The BLAS developed in three "levels": 1, 2, and 3. BLAS 1 includes vector-vector operations, BLAS 2 matrix-vector operations, and BLAS 3 matrix-matrix operations. The level coincides with the number of nested loops in a naïve sequential implementation of the operation. Increasing level also comes with increasing potential for data reuse. The BLAS traditionally lists computing a Givens rotation among the BLAS 1 operations, even though it only operates on scalars.

--end note]

Givens rotations [linalg.algs.blas1.givens]

Compute Givens rotation [linalg.algs.blas1.givens.lartg]

```cpp
template<class Real>
void givens_rotation_setup(const Real a,
    const Real b,
    Real& c,
    Real& s,
    Real& r);

template<class Real>
void givens_rotation_setup(const complex<Real>& a,
    const complex<Real>& a,
    Real& c,
    complex<Real>& s,
    complex<Real>& r);
```

This function computes the plane (Givens) rotation represented by the two values \( c \) and \( s \) such that the mathematical expression

\[
\begin{bmatrix}
c & s \\
-c & \* 
\end{bmatrix} \begin{bmatrix}
a \\
0 
\end{bmatrix} = \begin{bmatrix}
r \\
0 
\end{bmatrix}
\]

holds, where \( \text{conj} \) indicates the mathematical conjugate of \( s \), \( c \) is always a real scalar, and \( c \cdot c + \text{abs}(s) \cdot \text{abs}(s) \) equals one. That is, \( c \) and \( s \) represent a 2 x 2 matrix, that when multiplied by the right by the input vector whose components are \( a \) and \( b \), produces a result vector whose first component \( r \) is the Euclidean norm of the input vector, and whose second component as zero. [Note: The C++ Standard Library \( \text{conj} \) function always returns \( \text{complex<}T\text{>} \) for some \( T \), even though overloads exist for non-complex input. The above expression uses \( \text{conj} \) as mathematical notation, not as code. --end note]

[Note: This function corresponds to the LAPACK function \( \text{xLARTG} \). The BLAS variant \( \text{xROTG} \) takes four arguments -- \( a \), \( b \), \( c \), and \( s \)-- and overwrites the input \( a \) with \( r \). We have chosen \( \text{xLARTG} \)'s interface because it separates input and output, and to encourage following \( \text{xLARTG} \)'s more careful implementation. --end note]
Note: given_rotation_setup has an overload for complex numbers, because the output argument c (cosine) is a signed magnitude. --end note

- **Constraints:** Real is float, double, or long double.

- **Effects:** Assigns to c and s the plane (Givens) rotation corresponding to the input a and b. Assigns to r the Euclidean norm of the two-component vector formed by a and b.

- **Throws:** Nothing.

Apply a computed Givens rotation to vectors [linalg.algs.blas1.givens.rot]

```cpp
template<class inout_vector_1_t,
         class inout_vector_2_t,
         class Real>
void givens_rotation_apply(
    inout_vector_1_t x,
    inout_vector_2_t y,
    const Real c,
    const Real s);

template<class ExecutionPolicy,
         class inout_vector_1_t,
         class inout_vector_2_t,
         class Real>
void givens_rotation_apply(
    ExecutionPolicy&& exec,
    inout_vector_1_t x,
    inout_vector_2_t y,
    const Real c,
    const Real s);

template<class inout_vector_1_t,
         class inout_vector_2_t,
         class Real>
void givens_rotation_apply(
    inout_vector_1_t x,
    inout_vector_2_t y,
    const Real c,
    const complex<Real> s);

template<class ExecutionPolicy,
         class inout_vector_1_t,
         class inout_vector_2_t,
         class Real>
void givens_rotation_apply(
    ExecutionPolicy&& exec,
    inout_vector_1_t x,
    inout_vector_2_t y,
    const Real c,
    const complex<Real> s);
```

Note:
These functions correspond to the BLAS function \texttt{xROT}. \(c\) and \(s\) form a plane (Givens) rotation.
Users normally would compute \(c\) and \(s\) using \texttt{givens_rotation_setup}, but they are not required to do this.

\textit{---end note---}

- \textbf{Requires:} \(x\.\text{extent}(0)\) equals \(y\.\text{extent}(0)\).

- \textbf{Constraints:}
  
  \begin{itemize}
  
  \item \(\text{Real}\) is \texttt{float}, \texttt{double}, or \texttt{long\ double}.
  
  \item For the overloads that take the last argument \(s\) as \texttt{Real}, for \(i\) in the domain of \(x\) and \(j\) in the domain of \(y\), the expressions \(x(i) = c*x(i) + s*y(j)\) and \(y(j) = c*y(j) - s*x(i)\) are well formed.
  
  \item For the overloads that take the last argument \(s\) as \texttt{const complex<Real>}, for \(i\) in the domain of \(x\) and \(j\) in the domain of \(y\), the expressions \(x(i) = c*x(i) + s*y(j)\) and \(y(j) = c*y(j) - \text{conj}(s)*x(i)\) are well formed.
  
  \end{itemize}

- \textbf{Mandates:} If neither \(x\.\text{static\_extent}(0)\) nor \(y\.\text{static\_extent}(0)\) equals \texttt{dynamic\_extent}, then \(x\.\text{static\_extent}(0)\) equals \(y\.\text{static\_extent}(0)\).

- \textbf{Effects:} Applies the plane (Givens) rotation specified by \(c\) and \(s\) to the input vectors \(x\) and \(y\), as if the rotation were a \(2 \times 2\) matrix and the input vectors were successive rows of a matrix with two rows.

\textit{Swap matrix or vector elements [linalg.algs.blas1.swap]}

\begin{verbatim}
template<class inout_object_1_t, class inout_object_2_t>
void linalg_swap(inout_object_1_t x, inout_object_2_t y);

template<class ExecutionPolicy, class inout_object_1_t, class inout_object_2_t>
void linalg_swap(ExecutionPolicy& exec, inout_object_1_t x, inout_object_2_t y);
\end{verbatim}

\textit{[Note: These functions correspond to the BLAS function \texttt{xSWAP}. ---end note---}

- \textbf{Requires:} For all \(r\) in \(0, 1, \ldots\), \(x\.\text{rank}() - 1\), \(x\.\text{extent}(r)\) equals \(y\.\text{extent}(r)\).

- \textbf{Constraints:}
  
  \begin{itemize}
  
  \item \(x\.\text{rank}()\) equals \(y\.\text{rank}()\).
  
  \item \(x\.\text{rank}()\) is no more than 2.
  
  \item For \(i\ldots\) in the domain of \(x\) and \(y\), the expression \(x(i\ldots) = y(i\ldots)\) is well formed.
  
  \end{itemize}

2020-01-12
- **Mandates:** For all \( r \) in 0, 1, ..., \( x.\text{rank}() \) - 1, if neither \( x.\text{static_extent}(r) \) nor \( y.\text{static_extent}(r) \) equals \( \text{dynamic_extent} \), then \( x.\text{static_extent}(r) \) equals \( y.\text{static_extent}(r) \).

- **Effects:** Swap all corresponding elements of the objects \( x \) and \( y \).

**Multiply the elements of an object in place by a scalar** [linalg.algs.blas1.scal]

```cpp
template<class Scalar,  
   class inout_object_t>  
void scale(const Scalar alpha,  
            inout_object_t obj);

template<class ExecutionPolicy,  
   class Scalar,  
   class inout_object_t>  
void scale(ExecutionPolicy&& exec,  
            const Scalar alpha,  
            inout_object_t obj);
```

[Note: These functions correspond to the BLAS function \( \text{xSCAL} \). --end note]

- **Constraints:**
  - \( \text{obj.rank()} \) is no more than 3.
  - For \( i... \) in the domain of \( \text{obj} \), the expression \( \text{obj}(i...) \times= \text{alpha} \) is well formed.

- **Effects:** Multiply each element of \( \text{obj} \) in place by \( \text{alpha} \).

**Copy elements of one matrix or vector into another** [linalg.algs.blas1.copy]

```cpp
template<class in_object_t,  
   class out_object_t>  
void linalg_copy(in_object_t x,  
                 out_object_t y);

template<class ExecutionPolicy,  
   class in_object_t,  
   class out_object_t>  
void linalg_copy(ExecutionPolicy&& exec,  
                in_object_t x,  
                out_object_t y);
```

[Note: These functions correspond to the BLAS function \( \text{xCOPY} \). --end note]

- **Requires:** For all \( r \) in 0, 1, ..., \( x.\text{rank}() \) - 1, \( x.\text{extent}(r) \) equals \( y.\text{extent}(r) \).

- **Constraints:**
  - \( x.\text{rank}() \) equals \( y.\text{rank}() \).
- $x.rank()$ is no more than 2.
- For all $i...$ in the domain of $x$ and $y$, the expression $y(i...) = x(i...)$ is well formed.

**Mandates:** For all $r$ in $0, 1, ..., x.rank() - 1$, if neither $x.static_extent(r)$ nor $y.static_extent(r)$ equals $\text{dynamic_extent}$, then $x.static_extent(r)$ equals $y.static_extent(r)$.

**Effects:** Overwrite each element of $y$ with the corresponding element of $x$.

Add vectors or matrices elementwise [linalg.algs.blas1.add]

```cpp
template<class in_object_1_t, 
    class in_object_2_t, 
    class out_object_t>
void linalg_add(in_object_1_t x, 
    in_object_2_t y, 
    out_object_t z);

template<class ExecutionPolicy, 
    class in_object_1_t, 
    class in_object_2_t, 
    class out_object_t>
void linalg_add(ExecutionPolicy&& exec, 
    in_object_1_t x, 
    in_object_2_t y, 
    out_object_t z);
```

[Note: These functions correspond to the BLAS function $\text{xAXPY}$. --end note]

- **Requires:** For all $r$ in $0, 1, ..., x.rank() - 1$,
  - $x.extent(r)$ equals $z.extent(r)$.
  - $y.extent(r)$ equals $z.extent(r)$.

**Constraints:**
- $x.rank()$, $y.rank()$, and $z.rank()$ are all equal.
- $x.rank()$ is no more than 2.
- For $i...$ in the domain of $x$, $y$, and $z$, the expression $z(i...) = x(i...) + y(i...)$ is well formed.

**Mandates:** For all $r$ in $0, 1, ..., x.rank() - 1$,
- if neither $x.static_extent(r)$ nor $z.static_extent(r)$ equals $\text{dynamic_extent}$, then $x.static_extent(r)$ equals $z.static_extent(r)$; and
- if neither $y.static_extent(r)$ nor $z.static_extent(r)$ equals $\text{dynamic_extent}$, then $y.static_extent(r)$ equals $z.static_extent(r)$.

2020-01-12
if neither \( x.\text{static_extent}(r) \) nor \( y.\text{static_extent}(r) \) equals \( \text{dynamic_extent} \), then \( x.\text{static_extent}(r) \) equals \( y.\text{static_extent}(r) \);

- **Effects**: Compute the elementwise sum \( z = x + y \).

**Dot product of two vectors** [linalg.algs.blas1.dot]

**Nonconjugated dot product of two vectors** [linalg.algs.blas1.dot.dotu]

*[Note: The functions in this section correspond to the BLAS functions \( \text{xDOT} \) (for real element types) and \( \text{xDOTU} \) (for complex element types). --end note]*

**Nonconjugated dot product with specified result type**

```cpp
template<
class in_vector_1_t,
class in_vector_2_t,
class T>
T dot(in_vector_1_t v1,
in_vector_2_t v2,
T init);

template<class ExecutionPolicy,
class in_vector_1_t,
class in_vector_2_t,
class T>
T dot(ExecutionPolicy&& exec,
in_vector_1_t v1,
in_vector_2_t v2,
T init);
```

- **Requires**:
  - \( T \) shall be \textit{Cpp17MoveConstructible}.
  - \( \text{init} + v1(0) * v2(0) \) shall be convertible to \( T \).
  - \( v1.\text{extent}(0) \) equals \( v2.\text{extent}(0) \).

- **Constraints**: For all \( i \) in the domain of \( v1 \) and \( v2 \) and for \( \text{val} \) of type \( T& \), the expression \( \text{val } += v1(i) * v2(i) \) is well formed.

- **Mandates**: If neither \( v1.\text{static_extent}(0) \) nor \( v2.\text{static_extent}(0) \) equals \( \text{dynamic_extent} \), then \( v1.\text{static_extent}(0) \) equals \( v2.\text{static_extent}(0) \).

- **Effects**: Let \( N \) be \( v1.\text{extent}(0) \). If \( N \) is zero, returns \( \text{init} \), else returns \( \text{GENERALIZED}\_\text{SUM}(\text{plus} <> () , \text{init} , v1(0) * v2(0) , . . . , v1(N-1) \* v2(N-1)) \).

- **Remarks**: If \( \text{in}\_\text{vector}\_t::\text{element}\_\text{type} \) and \( T \) are both floating-point types or complex versions thereof, and if \( T \) has higher precision than \( \text{in}\_\text{vector}\_t::\text{element}\_\text{type} \), then implementations will use \( T\)'s precision or greater for intermediate terms in the sum.

*[Note: Like \( \text{reduce} \), \( \text{dot} \) applies binary \textit{operator}\+ in an unspecified order. This may yield a nondeterministic result for non-associative or non-commutative \textit{operator}\+ such as floating-]*
point addition. However, implementations may perform extra work to make the result
deterministic. They may do so for all dot overloads, or just for specific ExecutionPolicy
types. --end note]

[Note: Users can get xDOTC behavior by giving the second argument as a conjugate_view.
Alternately, they can use the shortcut dotc below. --end note]

Nonconjugated dot product with default result type

```cpp
template<class in_vector_1_t,
    class in_vector_2_t>
auto dot(in_vector_1_t v1,
    in_vector_2_t v2);

template<class ExecutionPolicy,
    class in_vector_1_t,
    class in_vector_2_t>
auto dot(ExecutionPolicy & exec,
    in_vector_1_t v1,
    in_vector_2_t v2);
```

- **Effects:** Let $T$ be `decltype(v1(0)*v2(0))`. Then, the two-parameter overload is
equivalent to `dot(v1, v2, T{})`; and the three-parameter overload is equivalent to
`dot(exec, v1, v2, T{})`.

Conjugated dot product of two vectors [linalg.algs.blas1.dot.dotc]

[Note:

The functions in this section correspond to the BLAS functions `xDOT` (for real element types)
and `xDOTC` (for complex element types).

dotc exists to give users reasonable default inner product behavior for both real and complex
element types.

--end note]

Conjugated dot product with specified result type

```cpp
template<class in_vector_1_t,
    class in_vector_2_t,
    class T>
T dotc(in_vector_1_t v1,
    in_vector_2_t v2,
    T init);

template<class ExecutionPolicy,
    class in_vector_1_t,
    class in_vector_2_t,
    class T>
T dotc(ExecutionPolicy & exec,
    in_vector_1_t v1,
    in_vector_2_t v2,
    T init);
```
• **Effects:** The three-argument overload is equivalent to `dot(v1, conjugate_view(v2), init);`. The four-argument overload is equivalent to `dot(exec, v1, conjugate_view(v2), init);`.

Conjugated dot product with default result type

```cpp
template<class in_vector_1_t, 
  class in_vector_2_t>
auto dotc(in_vector_1_t v1, 
  in_vector_2_t v2);
```

```cpp
template<class ExecutionPolicy, 
  class in_vector_1_t, 
  class in_vector_2_t>
auto dotc(ExecutionPolicy&& exec, 
  in_vector_1_t v1, 
  in_vector_2_t v2);
```

• **Effects:** If `in_vector_2_t::element_type` is `complex<R>` for some R, let `T` be `decltype(v1(0)*conj(v2(0)))`; else, let `T` be `decltype(v1(0)*v2(0))`. Then, the two-parameter overload is equivalent to `dotc(v1, v2, T{});`, and the three-parameter overload is equivalent to `dotc(exec, v1, v2, T{});`.

Euclidean norm of a vector \([\text{inneralg.algs.blas1.nrm2}]\)

Euclidean norm with specified result type

```cpp
template<class in_vector_t, 
  class T>
T vector_norm2(in_vector_t v, 
  T init);
```

```cpp
template<class ExecutionPolicy, 
  class in_vector_t, 
  class T>
T vector_norm2(ExecutionPolicy&& exec, 
  in_vector_t v, 
  T init);
```

[Note: These functions correspond to the BLAS function `xNRM2`. --end note]

• **Requires:**
  
  - `T` shall be `Cpp17MoveConstructible`.
  - `init + abs(v(0))*abs(v(0))` shall be convertible to `T`.

• **Constraints:** For all `i` in the domain of `v` and for `val` of type `T&`, the expressions `val += abs(v(i))*abs(v(i))` and `sqrt(val)` are well formed. [Note: This does not imply a recommended implementation for floating-point types. See Remarks below. --end note]

• **Effects:** Returns the Euclidean norm (also called 2-norm) of the vector `v`. 
• **Remarks:** If `in_vector_t::element_type` and `T` are both floating-point types or complex versions thereof, and if `T` has higher precision than `in_vector_type::element_type`, then implementations will use `T`'s precision or greater for intermediate terms in the sum.

*Note: We recommend that implementers document their guarantees regarding overflow and underflow of `vector_norm2` for floating-point return types. --end note*

Euclidean norm with default result type

```cpp
template<class in_vector_t>
auto vector_norm2(in_vector_t v);

template<class ExecutionPolicy, 
  class in_vector_t>
auto vector_norm2(ExecutionPolicy&& exec, 
  in_vector_t v);
```

• **Effects:** Let `T` be `decltype(abs(v(0)) * abs(v(0)))`. Then, the one-parameter overload is equivalent to `vector_norm2(v, T{})`; and the two-parameter overload is equivalent to `vector_norm2(exec, v, T{})`;

Sum of absolute values of vector elements [linalg.algs.blas1.asum]

Sum of absolute values with specified result type

```cpp
template<class in_vector_t, 
  class T>
T vector_abs_sum(in_vector_t v, 
  T init);

template<class ExecutionPolicy, 
  class in_vector_t, 
  class T>
T vector_abs_sum(ExecutionPolicy&& exec, 
  in_vector_t v, 
  T init);
```

*Note: This function corresponds to the BLAS functions `SASUM`, `DASUM`, `CSASUM`, and `DZASUM`. The different behavior for complex element types is based on the observation that this lower-cost approximation of the one-norm serves just as well as the actual one-norm for many linear algebra algorithms in practice. --end note*

• **Requires:**
  - `T` shall be `Cpp17MoveConstructible`.
  - `init + v1(0)*v2(0)` shall be convertible to `T`.

• **Constraints:** For all `i` in the domain of `v` and for `val` of type `T&`, the expression `val += abs(v(i))` is well formed.

• **Effects:** Let `N` be `v.extent(0)`.
If N is zero, returns \textit{init}.

Else, if \texttt{in\_vector\_t::element\_type} is \texttt{complex<R>} for some \texttt{R}, then returns \texttt{GENERALIZED\_SUM(plus\_<\rangle(), init, abs(real(v(0))) + abs(imag(v(0))), ..., abs(real(v(N-1))) + abs(imag(v(N-1)))}).

Else, returns \texttt{GENERALIZED\_SUM(plus\_<\rangle(), init, abs(v(0)), ..., abs(v(N-1)))}.

\textit{Remarks}: If \texttt{in\_vector\_t::element\_type} and \texttt{T} are both floating-point types or complex versions thereof, and if \texttt{T} has higher precision than \texttt{in\_vector\_t::element\_type}, then implementations will use \texttt{T}'s precision or greater for intermediate terms in the sum.

\textit{Sum of absolute values with default result type}

\begin{verbatim}
template<class in_vector_t>
auto vector_abs_sum(in_vector_t v);

template<class ExecutionPolicy,
class in_vector_t>
auto vector_abs_sum(ExecutionPolicy&& exec,
in_vector_t v);
\end{verbatim}

\textit{Effects}: Let \texttt{T} be \texttt{decltype(abs(v(0)))}. Then, the one-parameter overload is equivalent to \texttt{vector_abs_sum(v, T{});}, and the two-parameter overload is equivalent to \texttt{vector_abs_sum(exec, v, T{});}.

\textit{Index of maximum absolute value of vector elements} [\textit{linalg.algs.blas1.iamax}]

\begin{verbatim}
template<class in_vector_t>
ptrdiff_t idx_abs_max(in_vector_t v);

template<class ExecutionPolicy,
class in_vector_t>
pptrdiff_t idx_abs_max(ExecutionPolicy&& exec,
in_vector_t v);
\end{verbatim}

\textit{Note}: These functions correspond to the BLAS function \texttt{IxAMAX}. --end note

\textit{Constraints}: For \texttt{i} and \texttt{j} in the domain of \texttt{v}, the expression \texttt{abs(v(i)) < abs(v(j))} is well formed.

\textit{Effects}: Returns the index (in the domain of \texttt{v}) of the first element of \texttt{v} having largest absolute value. If \texttt{v} has zero elements, then returns \texttt{-1}.

\textit{BLAS 2 functions} [\textit{linalg.algs.blas2}]

\textit{General matrix-vector product} [\textit{linalg.algs.blas2.gemv}]

\textit{Note}: These functions correspond to the BLAS function \texttt{xGEMV}. --end note
The following requirements apply to all functions in this section.

- **Requires:**
  - A.extent(1) equals x.extent(0).
  - A.extent(0) equals y.extent(0).
  - y.extent(0) equals z.extent(0) (if applicable).

- **Constraints:** For all functions in this section:
  - in_matrix_t has unique layout; and
  - A.rank() equals 2, x.rank() equals 1, y.rank() equals 1, and z.rank() equals 1 (if applicable).

- **Mandates:**
  - If neither A.static_extent(1) nor x.static_extent(0) equals dynamic_extent, then A.static_extent(1) equals x.static_extent(0).
  - If neither A.static_extent(0) nor y.static_extent(0) equals dynamic_extent, then A.static_extent(0) equals y.static_extent(0).
  - If neither y.static_extent(0) nor z.static_extent(0) equals dynamic_extent, then y.static_extent(0) equals z.static_extent(0) (if applicable).

**Overwriting matrix-vector product**

```cpp
template<class in_vector_t,
         class in_matrix_t,
         class out_vector_t>
void matrix_vector_product(in_matrix_t A,
                           in_vector_t x,
                           out_vector_t y);

template<class ExecutionPolicy,
         class in_vector_t,
         class in_matrix_t,
         class out_vector_t>
void matrix_vector_product(ExecutionPolicy&& exec,
                           in_matrix_t A,
                           in_vector_t x,
                           out_vector_t y);
```

- **Constraints:** For i, j in the domain of A, the expression y(i) += A(i,j)*x(j) is well formed.

- **Effects:** Assigns to the elements of y the product of the matrix A with the vector x.

[Example:
constexpr ptrdiff_t num_rows = 5;
constexpr ptrdiff_t num_cols = 6;

// y = 3.0 * A * x
void scaled_matvec_1(mdspan<double, extents<num_rows, num_cols>> A,
                    mdspan<double, extents<num_cols>> x,
                    mdspan<double, extents<num_rows>> y)
{
    matrix_vector_product(scaled_view(3.0, A), x, y);
}

// y = 3.0 * A * x + 2.0 * y
void scaled_matvec_2(mdspan<double, extents<num_rows, num_cols>> A,
                    mdspan<double, extents<num_cols>> x,
                    mdspan<double, extents<num_rows>> y)
{
    matrix_vector_product(scaled_view(3.0, A), x,
                          scaled_view(2.0, y), y);
}

// z = 7.0 times the transpose of A, times y
void scaled_matvec_2(mdspan<double, extents<num_rows, num_cols>> A,
                    mdspan<double, extents<num_rows>> y,
                    mdspan<double, extents<num_cols>> z)
{
    matrix_vector_product(scaled_view(7.0, transpose_view(A)), y, z);
}

--end example]

Updating matrix-vector product

template<class in_vector_1_t,
         class in_matrix_t,
         class in_vector_2_t,
         class out_vector_t>
void matrix_vector_product(in_matrix_t A,
                           in_vector_1_t x,
                           in_vector_2_t y,
                           out_vector_t z);

template<class ExecutionPolicy,
         class in_vector_1_t,
         class in_matrix_t,
         class in_vector_2_t,
         class out_vector_t>
void matrix_vector_product(ExecutionPolicy&& exec,
                           in_matrix_t A,
                           in_vector_1_t x,
                           in_vector_2_t y,
                           out_vector_t z);

- **Constraints:** For $i,j$ in the domain of $A$, the expression $z(i) = y(i) + A(i,j) \times x(j)$ is well formed.
• *Effects:* Assigns to the elements of $z$ the elementwise sum of $y$, and the product of the matrix $A$ with the vector $x$.

Symmetric matrix-vector product [linalg.algs.blas2.symv]

[Note: These functions correspond to the BLAS functions $x$SYMV and $x$SPMV. --end note]

The following requirements apply to all functions in this section.

• *Requires:*
  
  o $A$.extent(0) equals $A$.extent(1).
  o $A$.extent(1) equals $x$.extent(0).
  o $A$.extent(0) equals $y$.extent(0).
  o $y$.extent(0) equals $z$.extent(0) (if applicable).

• *Constraints:*
  
  o $in\_matrix\_t$ either has unique layout, or $layout\_blas\_packed$ layout.
  o If $in\_matrix\_t$ has $layout\_blas\_packed$ layout, then the layout's Triangle template argument has the same type as the function's Triangle template argument.
  o $A$.rank() equals 2, $x$.rank() equals 1, $y$.rank() equals 1, and $z$.rank() equals 1 (if applicable).

• *Mandates:*
  
  o If neither $A$.static_extent(0) nor $A$.static_extent(1) equals $dynamic\_extent$, then $A$.static_extent(0) equals $A$.static_extent(1).
  o If neither $A$.static_extent(1) nor $x$.static_extent(0) equals $dynamic\_extent$, then $A$.static_extent(1) equals $x$.static_extent(0).
  o If neither $A$.static_extent(0) nor $y$.static_extent(0) equals $dynamic\_extent$, then $A$.static_extent(0) equals $y$.static_extent(0).
  o If neither $y$.static_extent(0) nor $z$.static_extent(0) equals $dynamic\_extent$, then $y$.static_extent(0) equals $z$.static_extent(0) (if applicable).

• *Remarks:* The functions will only access the triangle of $A$ specified by the Triangle argument $t$, and will assume for indices $i,j$ outside that triangle, that $A(j,i)$ equals $A(i,j)$.
class in_vector_t,
class out_vector_t>
void symmetric_matrix_vector_product(in_matrix_t A,
                                     Triangle t,
                                     in_vector_t x,
                                     out_vector_t y);

template<class ExecutionPolicy,
          class in_matrix_t,
          class Triangle,
          class in_vector_t,
          class out_vector_t>
void symmetric_matrix_vector_product(ExecutionPolicy&& exec,
                                     in_matrix_t A,
                                     Triangle t,
                                     in_vector_t x,
                                     out_vector_t y);

- **Constraints:** For \(i, j\) in the domain of \(A\), the expression \(y(i) \pm A(i,j)x(j)\) is well formed.

- **Effects:** Assigns to the elements of \(y\) the product of the matrix \(A\) with the vector \(x\).

Updating symmetric matrix-vector product

template<class in_matrix_t,
          class Triangle,
          class in_vector_1_t,
          class in_vector_2_t,
          class out_vector_t>
void symmetric_matrix_vector_product(
                                     in_matrix_t A,
                                     Triangle t,
                                     in_vector_1_t x,
                                     in_vector_2_t y,
                                     out_vector_t z);

template<class ExecutionPolicy,
          class in_matrix_t,
          class Triangle,
          class in_vector_1_t,
          class in_vector_2_t,
          class out_vector_t>
void symmetric_matrix_vector_product(ExecutionPolicy&& exec,
                                     in_matrix_t A,
                                     Triangle t,
                                     in_vector_1_t x,
                                     in_vector_2_t y,
                                     out_vector_t z);

- **Constraints:** For \(i, j\) in the domain of \(A\), the expression \(z(i) = y(i) + A(i,j)x(j)\) is well formed.
• **Effects:** Assigns to the elements of \( z \) the elementwise sum of \( y \), with the product of the matrix \( A \) with the vector \( x \).

**Hermitian matrix-vector product [linalg.algs.blas2.hemv]**

**[Note: These functions correspond to the BLAS functions xHEMV and xHPMV. --end note]**

The following requirements apply to all functions in this section.

• **Requires:**
  
  - \( A\).extent(0) equals \( A\).extent(1).
  - \( A\).extent(1) equals \( x\).extent(0).
  - \( A\).extent(0) equals \( y\).extent(0).
  - \( y\).extent(0) equals \( z\).extent(0) (if applicable).

• **Constraints:**
  
  - \( \text{in} \_\text{matrix} \_t \) either has unique layout, or \( \text{layout} \_\text{blas} \_\text{packed} \) layout.
  
  - If \( \text{in} \_\text{matrix} \_t \) has \( \text{layout} \_\text{blas} \_\text{packed} \) layout, then the layout's Triangle template argument has the same type as the function's Triangle template argument.
  
  - \( A\).rank() equals 2, \( x\).rank() equals 1, \( y\).rank() equals 1, and \( z\).rank() equals 1.

• **Mandates:**

  - If neither \( A\).static_extent(0) nor \( A\).static_extent(1) equals dynamic_extent, then \( A\).static_extent(0) equals \( A\).static_extent(1).

  - If neither \( A\).static_extent(1) nor \( x\).static_extent(0) equals dynamic_extent, then \( A\).static_extent(1) equals \( x\).static_extent(0).

  - If neither \( A\).static_extent(0) nor \( y\).static_extent(0) equals dynamic_extent, then \( A\).static_extent(0) equals \( y\).static_extent(0).

  - If neither \( y\).static_extent(0) nor \( z\).static_extent(0) equals dynamic_extent, then \( y\).static_extent(0) equals \( z\).static_extent(0) (if applicable).

• **Remarks:**

  - The functions will only access the triangle of \( A \) specified by the Triangle argument \( t \).

  - If \( \text{inout} \_\text{matrix} \_t::\text{element} \_\text{type} \) is complex\(<\text{RA}>\) for some \( \text{RA} \), then the functions will assume for indices \( i,j \) outside that triangle, that \( A(j,i) \) equals \( \text{conj}(A(i,j)) \). Otherwise, the functions will assume that \( A(j,i) \) equals \( A(i,j) \).
Overwriting Hermitian matrix-vector product

```cpp
template<class in_matrix_t,  
    class Triangle,  
    class in_vector_t,  
    class out_vector_t>
void hermitian_matrix_vector_product(in_matrix_t A,  
    Triangle t,  
    in_vector_t x,  
    out_vector_t y);
```

```cpp
template<class ExecutionPolicy,  
    class in_matrix_t,  
    class Triangle,  
    class in_vector_t,  
    class out_vector_t>
void hermitian_matrix_vector_product(ExecutionPolicy&& exec,  
    in_matrix_t A,  
    Triangle t,  
    in_vector_t x,  
    out_vector_t y);
```

- **Constraints:** For $i,j$ in the domain of $A$:
  - the expression $y(i) += A(i,j) \times (j)$ is well formed; and
  - if `in_matrix_type::element_type` is `complex<RA>` for some `RA`, then the expression $y(i) += \text{conj}(A(i,j)) \times (j)$ is well formed.

- **Effects:** Assigns to the elements of $y$ the product of the matrix $A$ with the vector $x$.

Updating Hermitian matrix-vector product

```cpp
template<class in_matrix_t,  
    class Triangle,  
    class in_vector_1_t,  
    class in_vector_2_t,  
    class out_vector_t>
void hermitian_matrix_vector_product(in_matrix_t A,  
    Triangle t,  
    in_vector_1_t x,  
    in_vector_2_t y,  
    out_vector_t z);
```

```cpp
template<class ExecutionPolicy,  
    class in_matrix_t,  
    class Triangle,  
    class in_vector_1_t,  
    class in_vector_2_t,  
    class out_vector_t>
void hermitian_matrix_vector_product(ExecutionPolicy&& exec,  
    in_matrix_t A,  
    Triangle t,  
    in_vector_1_t x,  
    in_vector_1_t y,  
    out_vector_t z);
```
• **Constraints:** For \(i, j\) in the domain of \(A\):
  
  o the expression \(z(i) = y(i) + A(i,j)\times(x(j))\) is well formed; and
  
  o if `in_matrix_t::element_type` is `complex<RA>` for some `RA`, then the expression \(z(i) = y(i) + \text{conj}(A(i,j))\times(x(j))\) is well formed.

• **Effects:** Assigns to the elements of \(z\) the elementwise sum of \(y\), and the product of the matrix \(A\) with the vector \(x\).

**Triangular matrix-vector product [linalg.algs.blas2.trmv]**

[Note: These functions correspond to the BLAS functions `xTRMV` and `xTPMV`. --end note]

The following requirements apply to all functions in this section.

• **Requires:**
  
  o \(A\.extent(0)\) equals \(A\.extent(1)\).
  
  o \(A\.extent(0)\) equals \(y\.extent(0)\).
  
  o \(A\.extent(1)\) equals \(x\.extent(0)\) (if applicable).
  
  o \(y\.extent(0)\) equals \(z\.extent(0)\) (if applicable).

• **Constraints:**
  
  o `in_matrix_t` either has unique layout, or `layout_blas_packed` layout.
  
  o If `in_matrix_t` has `layout_blas_packed` layout, then the layout's `Triangle` template argument has the same type as the function's `Triangle` template argument.
  
  o \(A\.rank()\) equals 2.
  
  o \(y\.rank()\) equals 1.
  
  o \(x\.rank()\) equals 1 (if applicable).
  
  o \(z\.rank()\) equals 1 (if applicable).

• **Mandates:**
  
  o If neither \(A\.static_extent(0)\) nor \(A\.static_extent(1)\) equals `dynamic_extent`, then \(A\.static_extent(0)\) equals \(A\.static_extent(1)\).
  
  o If neither \(A\.static_extent(0)\) nor \(y\.static_extent(0)\) equals `dynamic_extent`, then \(A\.static_extent(0)\) equals \(y\.static_extent(0)\).
  
  o If neither \(A\.static_extent(1)\) nor \(x\.static_extent(0)\) equals `dynamic_extent`, then \(A\.static_extent(1)\) equals \(x\.static_extent(0)\) (if...
If neither \( y.\text{static\_extent}(0) \) nor \( z.\text{static\_extent}(0) \) equals 
\( \text{dynamic\_extent} \), then \( y.\text{static\_extent}(0) \) equals \( z.\text{static\_extent}(0) \) (if applicable).

**Remarks:**

- The functions will only access the triangle of \( A \) specified by the \( \text{Triangle} \) argument \( t \).
- If the \( \text{DiagonalStorage} \) template argument has type 
  \( \text{implicit\_unit\_diagonal\_t} \), then the functions will not access the diagonal of 
  \( A \), and will assume that that the diagonal elements of \( A \) all equal one. [Note: This 
  does not imply that the function needs to be able to form an \( \text{element\_type} \) value 
  equal to one. --end note]

Overwriting triangular matrix-vector product [\text{linalg.algs.blas2.trmv}]

```cpp
template<class in_matrix_t, 
    class Triangle, 
    class DiagonalStorage, 
    class in_vector_t, 
    class out_vector_t>
void triangular_matrix_vector_product(
    in_matrix_t A, 
    Triangle t, 
    DiagonalStorage d, 
    in_vector_t x, 
    out_vector_t y);

template<class ExecutionPolicy, 
    class in_matrix_t, 
    class Triangle, 
    class DiagonalStorage, 
    class in_vector_t, 
    class out_vector_t>
void triangular_matrix_vector_product(
    ExecutionPolicy&& exec, 
    in_matrix_t A, 
    Triangle t, 
    DiagonalStorage d, 
    in_vector_t x, 
    out_vector_t y);
```

- **Constraints:** For \( i, j \) in the domain of \( A \), the expression \( y(i) += A(i,j)*x(j) \) is well 
  formed.
- **Effects:** Assigns to the elements of \( y \) the product of the matrix \( A \) with the vector \( x \).

In-place triangular matrix-vector product [\text{linalg.algs.blas2.trmv.in-place}]

2020-01-12
template<class in_matrix_t,  
    class Triangle,  
    class DiagonalStorage,  
    class inout_vector_t>
void triangular_matrix_vector_product(
    in_matrix_t A,  
    Triangle t,  
    DiagonalStorage d,  
    inout_vector_t y);

- **Requires:** \(A\text{.extent}(1)\) equals \(y\text{.extent}(0)\).

- **Constraints:** For \(i, j\) in the domain of \(A\), the expression \(y(i) += A(i, j) \times y(j)\) is well formed.

- **Mandates:** If neither \(A\text{.static_extent}(1)\) nor \(y\text{.static_extent}(0)\) equals \(\text{dynamic_extent}\), then \(A\text{.static_extent}(1)\) equals \(y\text{.static_extent}(0)\).

- **Effects:** Overwrites \(y\) (on output) with the product of the matrix \(A\) with the vector \(y\) (on input).

**Updating triangular matrix-vector product** [linalg.algs.blas2.trmv.up]

```cpp
template<class in_matrix_t,  
    class Triangle,  
    class DiagonalStorage,  
    class in_vector_1_t,  
    class in_vector_2_t,  
    class out_vector_t>
void triangular_matrix_vector_product(in_matrix_t A,  
                                        Triangle t,  
                                        DiagonalStorage d,  
                                        in_vector_1_t x,  
                                        in_vector_2_t y,  
                                        out_vector_t z);

template<class ExecutionPolicy,  
    class in_matrix_t,  
    class Triangle,  
    class DiagonalStorage,  
    class in_vector_1_t,  
    class in_vector_2_t,  
    class out_vector_t>
void triangular_matrix_vector_product(ExecutionPolicy&& exec,  
                                        in_matrix_t A,  
                                        Triangle t,  
                                        DiagonalStorage d,  
                                        in_vector_1_t x,  
                                        in_vector_2_t y,  
                                        out_vector_t z);

- **Constraints:** For \(i, j\) in the domain of \(A\), the expression \(z(i) = y(i) + A(i, j) \times x(j)\)
  is well formed.
- **Effects**: Assigns to the elements of \( z \) the elementwise sum of \( y \), with the product of the matrix \( A \) with the vector \( x \).

**Solve a triangular linear system** [linalg.algs.blas2.trsv]

*[Note: These functions correspond to the BLAS functions \texttt{xTRSV} and \texttt{xTPSV}. --end note]*

The following requirements apply to all functions in this section.

- **Requires**:
  - \( A.\text{extent}(0) \) equals \( A.\text{extent}(1) \).
  - \( A.\text{extent}(1) \) equals \( b.\text{extent}(0) \).

- **Constraints**:
  - \( A.\text{rank}() \) equals 2.
  - \( b.\text{rank}() \) equals 1.
  - \( \text{in}_{\text{matrix}}.t \) either has unique layout, or \texttt{layout_blas_packed} layout.
  - If \( \text{in}_{\text{matrix}}.t \) has \texttt{layout_blas_packed} layout, then the layout's \texttt{Triangle} template argument has the same type as the function's \texttt{Triangle} template argument.

- **Mandates**:
  - If neither \( A.\text{static}_{\text{extent}}(0) \) nor \( A.\text{static}_{\text{extent}}(1) \) equals \texttt{dynamic}_{\text{extent}}, then \( A.\text{static}_{\text{extent}}(0) \) equals \( A.\text{static}_{\text{extent}}(1) \).
  - If neither \( A.\text{static}_{\text{extent}}(1) \) nor \( b.\text{static}_{\text{extent}}(0) \) equals \texttt{dynamic}_{\text{extent}}, then \( A.\text{static}_{\text{extent}}(1) \) equals \( b.\text{static}_{\text{extent}}(0) \).

- **Remarks**:
  - The functions will only access the triangle of \( A \) specified by the \texttt{Triangle} argument \( t \).
  - If the \texttt{DiagonalStorage} template argument has type \texttt{implicit_unit_diagonal_t}, then the functions will not access the diagonal of \( A \), and will assume that that the diagonal elements of \( A \) all equal one. *[Note: This does not imply that the function needs to be able to form an \texttt{element}_{\text{type}} value equal to one. --*end note]*

**Not-in-place triangular solve** [linalg.algs.blas2.trsv.not-in-place]

```template<
class in_matrix_t,
class Triangle,
class DiagonalStorage,
class in_vector_t,
class out_vector_t>
```
void triangular_matrix_vector_solve(
    in_matrix_t A,
    Triangle t,
    DiagonalStorage d,
    in_vector_t b,
    out_vector_t x);

template<class ExecutionPolicy,
    class in_matrix_t,
    class Triangle,
    class DiagonalStorage,
    class in_vector_t,
    class out_vector_t>
void triangular_matrix_vector_solve(
    ExecutionPolicy&& exec,
    in_matrix_t A,
    Triangle t,
    DiagonalStorage d,
    in_vector_t b,
    out_vector_t x);

- Requires:
  - \( A.\text{extent(0)} \) equals \( x.\text{extent(0)} \).

- Constraints:
  - \( x.\text{rank()} \) equals 1.
  - If \( r \) is in the domain of \( x \) and \( b \), then the expression \( x(r) = b(r) \) is well formed.
  - If \( r \) is in the domain of \( x \) and \( c \) is in the domain of \( x \), then the expression \( x(r) = A(r,c) \times x(c) \) is well formed.
  - If \( r \) is in the domain of \( x \) and \( \text{DiagonalStorage} \) is \( \text{explicit_diagonal_t} \), then the expression \( x(r) /= A(r,r) \) is well formed.

- Mandates:
  - If neither \( A.\text{static_extent(0)} \) nor \( x.\text{static_extent(0)} \) equals \( \text{dynamic_extent} \), then \( A.\text{static_extent(0)} \) equals \( x.\text{static_extent(0)} \).

- Effects: Assigns to the elements of \( x \) the result of solving the triangular linear system \( Ax=b \).

In-place triangular solve [linalg.algs.blas2.trsv.in-place]
DiagonalStorage d,
inout_vector_t b);

[Note:
This in-place version of the function intentionally lacks an overload taking an
ExecutionPolicy&, because it is not possible to parallelize in-place triangular solve for an
arbitrary ExecutionPolicy.

--end note]

• Requires:
  o A.extent(0) equals b.extent(0).

• Constraints:
  o If r and c are in the domain of b, then the expression b(r) -= A(r,c)*b(c) is
    well formed.
  o If r is in the domain of b and DiagonalStorage is explicit_diagonal_t, then
    the expression b(r) /= A(r,r) is well formed.

• Mandates:
  o If neither A.static_extent(0) nor b.static_extent(0) equals
    dynamic_extent, then A.static_extent(0) equals b.static_extent(0).

• Effects: Overwrites b with the result of solving the triangular linear system A x = b for x.

Rank-1 (outer product) update of a matrix [linalg.algs.blas2.rank1]

Nonsymmetric nonconjugated rank-1 update [linalg.algs.blas2.rank1.geru]
Relevant code snippet:

```c
template<class in_vector_1_t,
    class in_vector_2_t,
    class inout_matrix_t>
void matrix_rank_1_update_c(
    in_vector_1_t x,
    in_vector_2_t y,
    inout_matrix_t A);

template<class ExecutionPolicy,
    class in_vector_1_t,
    class in_vector_2_t,
    class inout_matrix_t>
void matrix_rank_1_update_c(
    ExecutionPolicy&& exec,
    in_vector_1_t x,
    in_vector_2_t y,
    inout_matrix_t A);
```

**Note:** This function corresponds to the BLAS functions `xGER` (for real element types) and `xGERU` (for complex element types). --end note

- **Requires:**
  - \(A.\text{extent}(0)\) equals \(x.\text{extent}(0)\).
  - \(A.\text{extent}(1)\) equals \(y.\text{extent}(0)\).

- **Constraints:**
  - \(A.\text{rank}()\) equals 2, \(x.\text{rank}()\) equals 1, and \(y.\text{rank}()\) equals 1.
  - For \(i, j\) in the domain of \(A\), the expression \(A(i, j) := x(i) \times y(j)\) is well formed.

- **Mandates:**
  - If neither \(A.\text{static}\_\text{extent}(0)\) nor \(x.\text{static}\_\text{extent}(0)\) equals \(\text{dynamic}\_\text{extent}\), then \(A.\text{static}\_\text{extent}(0)\) equals \(x.\text{static}\_\text{extent}(0)\).
  - If neither \(A.\text{static}\_\text{extent}(1)\) nor \(y.\text{static}\_\text{extent}(0)\) equals \(\text{dynamic}\_\text{extent}\), then \(A.\text{static}\_\text{extent}(1)\) equals \(y.\text{static}\_\text{extent}(0)\).

- **Effects:** Assigns to \(A\) on output the sum of \(A\) on input, and the outer product of \(x\) and \(y\).

**Note:** Users can get `xGERC` behavior by giving the second argument as a `conjugate_view`. Alternately, they can use the shortcut `matrix_rank_1_update_c` below. --end note

Nonsymmetric conjugated rank-1 update [linalg.algs.blas.rank1.ger]

---

**Note:** This function corresponds to the BLAS functions `xGER` (for real element types) and `xGERC` (for complex element types). --end note

- **Effects:** Equivalent to `matrix_rank_1_update(x, conjugate_view(y), A);`
template<class in_vector_t,  
class inout_matrix_t,  
class Triangle>
void symmetric_matrix_rank_1_update(  
in_vector_t x,  
inout_matrix_t A,  
Triangle t);

template<class ExecutionPolicy,  
class in_vector_t,  
class inout_matrix_t,  
class Triangle>
void symmetric_matrix_rank_1_update(  
ExecutionPolicy exec,  
in_vector_t x,  
inout_matrix_t A,  
Triangle t);

template<class T,  
class in_vector_t,  
class inout_matrix_t,  
class Triangle>
void symmetric_matrix_rank_1_update(  
T alpha,  
in_vector_t x,  
inout_matrix_t A,  
Triangle t);

template<class ExecutionPolicy,  
class T,  
class in_vector_t,  
class inout_matrix_t,  
class Triangle>
void symmetric_matrix_rank_1_update(  
ExecutionPolicy exec,  
T alpha,  
in_vector_t x,  
inout_matrix_t A,  
Triangle t);

[Note:

These functions correspond to the BLAS functions xSYR and xSPR.

They take an optional scaling factor alpha, because it would be impossible to express the update \(C = C - x x^T\) otherwise.

--end note]

- Requires:
  - \(\text{A.extent(0)}\) equals \(\text{A.extent(1)}\).
  - \(\text{A.extent(0)}\) equals \(\text{x.extent(0)}\).

- Constraints:
  - \(\text{A.rank()}\) equals 2 and \(\text{x.rank()}\) equals 1.
- A either has unique layout, or `layout_blas_packed` layout.
- If A has `layout_blas_packed` layout, then the layout’s Triangle template argument has the same type as the function’s Triangle template argument.
- For overloads without alpha:
  - For i, j in the domain of A, the expression \( A(i, j) += x(i) \cdot x(j) \) is well formed.
- For overloads with alpha:
  - For i, j in the domain of C, and i, k and k, i in the domain of A, the expression \( C(i, j) += \alpha \cdot A(i, k) \cdot A(j, k) \) is well formed.

- **Mandates:**
  - If neither \( A\.static\_extent(0) \) nor \( A\.static\_extent(1) \) equals dynamic_extent, then \( A\.static\_extent(0) \) equals \( A\.static\_extent(1) \).
  - If neither \( A\.static\_extent(0) \) nor \( x\.static\_extent(0) \) equals dynamic_extent, then \( A\.static\_extent(0) \) equals \( x\.static\_extent(0) \).

- **Effects:**
  - Overloads without alpha assign to A on output, the elementwise sum of A on input, with (the outer product of \( x \) and \( x \)).
  - Overloads with alpha assign to A on output, the elementwise sum of A on input, with alpha times (the outer product of \( x \) and \( x \)).

- **Remarks:** The functions will only access the triangle of A specified by the Triangle argument t, and will assume for indices i, j outside that triangle, that \( A(j, i) \) equals \( A(i, j) \).

Rank-1 update of a Hermitian matrix [linalg.algs.blas.rank1.her]

```cpp
template<class in_vector_t,  
class inout_matrix_t,  
class Triangle>
void hermitian_matrix_rank_1_update(  
in_vector_t x,  
inout_matrix_t A,  
Triangle t);

template<class ExecutionPolicy,  
class in_vector_t,  
class inout_matrix_t,  
class Triangle>
void hermitian_matrix_rank_1_update(  
ExecutionPolicy&& exec,  
in_vector_t x,  
inout_matrix_t A,  
Triangle t);

template<class T,
```
class in_vector_t,
class inout_matrix_t,
class Triangle>
void hermitian_matrix_rank_1_update(
  T alpha,
  in_vector_t x,
  inout_matrix_t A,
  Triangle t);

template<class ExecutionPolicy,
  class T,
  class in_vector_t,
  class inout_matrix_t,
  class Triangle>
void hermitian_matrix_rank_1_update(
  ExecutionPolicy&& exec,
  T alpha,
  in_vector_t x,
  inout_matrix_t A,
  Triangle t);

[Note:

These functions correspond to the BLAS functions xHER and xHPR.

They take an optional scaling factor \( \alpha \), because it would be impossible to express the update \( A = A - x x^H \) otherwise.

--end note]

- Requires:
  - \( A\.extent(0) \) equals \( A\.extent(1) \).
  - \( A\.extent(0) \) equals \( x\.extent(0) \).

- Constraints:
  - \( A\.rank() \) equals 2 and \( x\.rank() \) equals 1.
  - \( A \) either has unique layout, or layout_blas_packed layout.
  - If \( A \) has layout_blas_packed layout, then the layout's Triangle template argument has the same type as the function's Triangle template argument.
  - For overloads without \( \alpha \):
    - For \( i,j \) in the domain of \( A \):
      - if \( \text{in_vector_t::element_type is complex<RX>} \) for some \( RX \), then the expression \( A(i,j) += x(i) * \text{conj}(x(j)) \)' is well formed;
      - else, the expression \( A(i,j) += x(i) * x(j) \) is well formed.
  - For overloads with \( \alpha \):
    - For \( i,j \) in the domain of \( A \):
- if `in_vector_t::element_type` is `complex<RA>` for some RA, then the expression `A(i,j) += alpha*x(i)*conj(x(j))` is well formed;

- else, the expression `A(i,j) += alpha*x(i)*x(j)` is well formed.

**Mandates:**

- If neither `A.static_extent(0)` nor `A.static_extent(1)` equals `dynamic_extent`, then `A.static_extent(0)` equals `A.static_extent(1)`.

- If neither `A.static_extent(0)` nor `x.static_extent(0)` equals `dynamic_extent`, then `A.static_extent(0)` equals `x.static_extent(0)`.

**Effects:**

- Overloads without `alpha` assign to `A` on output, the elementwise sum of `A` on input, with (the outer product of `x` and the conjugate of `x`).

- Overloads with `alpha` assign to `A` on output, the elementwise sum of `A` on input, with `alpha` times (the outer product of `x` and the conjugate of `x`).

**Remarks:**

- The functions will only access the triangle of `A` specified by the `Triangle` argument `t`.

- If `inout_matrix_t::element_type` is `complex<RA>` for some RA, then the functions will assume for indices `i,j` outside that triangle, that `A(j,i)` equals `conj(A(i,j))`. Otherwise, the functions will assume that `A(j,k)` equals `A(i,j)`.

**Rank-2 update of a symmetric matrix** [linalg.algs.blas2.rank2.syr2]

```cpp
template<class in_vector_1_t,
         class in_vector_2_t,
         class inout_matrix_t,
         class Triangle>
void symmetric_matrix_rank_2_update(
    in_vector_1_t x,
    in_vector_2_t y,
    inout_matrix_t A,
    Triangle t);

template<class ExecutionPolicy,
         class in_vector_1_t,
         class in_vector_2_t,
         class inout_matrix_t,
         class Triangle>
void symmetric_matrix_rank_2_update(
    ExecutionPolicy& exec,
    in_vector_1_t x,
    in_vector_2_t y,
    inout_matrix_t A,
    Triangle t);
```
Mark Hoemmen

[Note: These functions correspond to the BLAS functions xSYR2 and xSPR2. --end note]

- **Requires:**
  - A.extent(0) equals A.extent(1).
  - A.extent(0) equals x.extent(0).
  - A.extent(0) equals y.extent(0).

- **Constraints:**
  - A.rank() equals 2, x.rank() equals 1, and y.rank() equals 1.
  - A either has unique layout, or layout_blas_packed layout.
  - If A has layout_blas_packed layout, then the layout's Triangle template argument has the same type as the function's Triangle template argument.
  - For i, j in the domain of A, the expression A(i,j) += x(i)*y(j) + y(i)*x(j) is well formed.

- **Mandates:**
  - If neither A.static_extent(0) nor A.static_extent(1) equals dynamic_extent, then A.static_extent(0) equals A.static_extent(1).
  - If neither A.static_extent(0) nor x.static_extent(0) equals dynamic_extent, then A.static_extent(0) equals x.static_extent(0).
  - If neither A.static_extent(0) nor y.static_extent(0) equals dynamic_extent, then A.static_extent(0) equals y.static_extent(0).

- **Effects:** Assigns to A on output the sum of A on input, the outer product of x and y, and the outer product of y and x.

- **Remarks:** The functions will only access the triangle of A specified by the Triangle argument t, and will assume for indices i, j outside that triangle, that A(j, i) equals A(i, j).

**Rank-2 update of a Hermitian matrix [linalg.algs.blas.rank2.her2]**

```cpp
template<class in_vector_1_t, class in_vector_2_t, class inout_matrix_t, class Triangle>
void hermitian_matrix_rank_2_update(
    in_vector_1_t x,
    in_vector_2_t y,
    inout_matrix_t A,
    Triangle t);

template<class ExecutionPolicy, class in_vector_1_t, class in_vector_2_t,
```
class inout_matrix_t,
class Triangle>
void hermitian_matrix_rank_2_update(
  ExecutionPolicy&& exec,
  in_vector_1_t x,
  in_vector_2_t y,
  inout_matrix_t A,
  Triangle t);

[Note: These functions correspond to the BLAS functions xHER2 and xHPR2. --end note]

- **Requires:**
  - A.extent(0) equals A.extent(1).
  - A.extent(0) equals x.extent(0).
  - A.extent(0) equals y.extent(0).

- **Constraints:**
  - A.rank() equals 2, x.rank() equals 1, and y.rank() equals 1.
  - A either has unique layout, or layout_blas_packed layout.
  - If A has layout_blas_packed layout, then the layout's Triangle template argument has the same type as the function's Triangle template argument.
  - For i, j in the domain of A:
    - If in_vector_2_t::element_type is complex<RY> for some RY,
      - if in_vector_1_t::element_type is complex<RX> for some RX,
        then the expression A(i,j) += x(i)*conj(y(j)) + y(i)*conj(x(j)) is well formed;
      - else, the expression A(i,j) += x(i)*conj(y(j)) + y(i)*x(j) is well formed;
    - else,
      - if in_vector_1_t::element_type is complex<RX> for some RX,
        then the expression A(i,j) += x(i)*y(j) + y(i)*conj(x(j)) is well formed;
      - else, the expression A(i,j) += x(i)*y(j)) + y(i)*x(j) is well formed.

- **Mandates:**
  - If neither A.static_extent(0) nor A.static_extent(1) equals dynamic_extent, then A.static_extent(0) equals A.static_extent(1).
  - If neither A.static_extent(0) nor x.static_extent(0) equals dynamic_extent, then A.static_extent(0) equals x.static_extent(0).
• If neither A.static_extent(0) nor y.static_extent(0) equals dynamic_extent, then A.static_extent(0) equals y.static_extent(0).

• **Effects:** Assigns to A on output the sum of A on input, the outer product of x and the conjugate of y, and the outer product of y and the conjugate of x.

• **Remarks:**
  
  • The functions will only access the triangle of A specified by the Triangle argument t.

  • If inout_matrix_t::element_type is complex<RA> for some RA, then the functions will assume for indices i, j outside that triangle, that A(j,i) equals conj(A(i,j)). Otherwise, the functions will assume that A(j,i) equals A(i,j).

**BLAS 3 functions [linalg.algs.blas3]**

**General matrix-matrix product [linalg.algs.blas3.gemm]**

*Note: These functions correspond to the BLAS function xGEMM.* --end note

The following requirements apply to all functions in this section.

• **Requires:**
  
  • C.extent(0) equals E.extent(0) (if applicable).

  • C.extent(1) equals E.extent(1) (if applicable).

  • A.extent(1) equals B.extent(0).

  • A.extent(0) equals C.extent(0).

  • B.extent(1) equals C.extent(1).

• **Constraints:**

  • *in_matrix_1_t, in_matrix_2_t, in_matrix_3_t* (if applicable), and *out_matrix_t* have unique layout.

  • A.rank() equals 2, B.rank() equals 2, C.rank() equals 2, and E.rank() (if applicable) equals 2.

• **Mandates:**

  • For all r in 0, 1, ..., C.rank() - 1, if neither C.static_extent(r) nor E.static_extent(r) equals dynamic_extent, then C.static_extent(r) equals E.static_extent(r) (if applicable).

  • If neither A.static_extent(1) nor B.static_extent(0) equals dynamic_extent, then A.static_extent(1) equals B.static_extent(0).

  • If neither A.static_extent(0) nor C.static_extent(0) equals dynamic_extent, then A.static_extent(0) equals C.static_extent(0).
If neither $B\.static\text{\_}extent(1)$ nor $C\.static\text{\_}extent(1)$ equals dynamic_extent, then $B\.static\text{\_}extent(1)$ equals $C\.static\text{\_}extent(1)$.

**Overwriting general matrix-matrix product**

```cpp
template<class in_matrix_1_t, 
class in_matrix_2_t, 
class out_matrix_t>
void matrix_product(in_matrix_1_t A, 
in_matrix_2_t B, 
out_matrix_t C);

template<class ExecutionPolicy, 
class in_matrix_1_t, 
class in_matrix_2_t, 
class out_matrix_t>
void matrix_product(ExecutionPolicy&& exec, 
in_matrix_1_t A, 
in_matrix_2_t B, 
out_matrix_t C);
```

- **Constraints:** For $i, j$ in the domain of $C$, $i, k$ in the domain of $A$, and $k, j$ in the domain of $B$, the expression $C(i, j) \leftarrow A(i, k) \times B(k, j)$ is well formed.

- **Effects:** Assigns to the elements of the matrix $C$ the product of the matrices $A$ and $B$.

**Updating general matrix-matrix product**

```cpp
template<class in_matrix_1_t, 
class in_matrix_2_t, 
class in_matrix_3_t, 
class out_matrix_t>
void matrix_product(in_matrix_1_t A, 
in_matrix_2_t B, 
in_matrix_3_t E, 
out_matrix_t C);

template<class ExecutionPolicy, 
class in_matrix_1_t, 
class in_matrix_2_t, 
class in_matrix_3_t, 
class out_matrix_t>
void matrix_product(ExecutionPolicy&& exec, 
in_matrix_1_t A, 
in_matrix_2_t B, 
in_matrix_3_t E, 
out_matrix_t C);
```

- **Constraints:** For $i, j$ in the domain of $C$, $i, k$ in the domain of $A$, and $k, j$ in the domain of $B$, the expression $C(i, j) \leftarrow E(i, j) + A(i, k) \times B(k, j)$ is well formed.
• **Effects:** Assigns to the elements of the matrix $C$ on output, the elementwise sum of $E$ and the product of the matrices $A$ and $B$.

• **Remarks:** $C$ and $E$ may refer to the same matrix. If so, then they must have the same layout.

**Symmetric matrix-matrix product** [linalg.algs.blas3.symm]

*Note:* These functions correspond to the BLAS function *xSYMm*.

Unlike the symmetric rank-1 update functions, these functions assume that the input matrix -- not the output matrix -- is symmetric.

--end note

The following requirements apply to all functions in this section.

• **Requires:**
  - $A.extent(0)$ equals $A.extent(1)$.
  - $C.extent(0)$ equals $E.extent(0)$ (if applicable).
  - $C.extent(1)$ equals $E.extent(1)$ (if applicable).

• **Constraints:**
  - $in\_matrix\_1\_t$ either has unique layout, or $layout\_blas\_packed$ layout.
  - $in\_matrix\_2\_t$, $in\_matrix\_3\_t$ (if applicable), and $out\_matrix\_t$ have unique layout.
  - If $in\_matrix\_t$ has $layout\_blas\_packed$ layout, then the layout's Triangle template argument has the same type as the function's Triangle template argument.
  - $A.rank()$ equals 2, $B.rank()$ equals 2, $C.rank()$ equals 2, and $E.rank()$ (if applicable) equals 2.

• **Mandates:**
  - If neither $A.static\_extent(0)$ nor $A.static\_extent(1)$ equals $dynamic\_extent$, then $A.static\_extent(0)$ equals $A.static\_extent(1)$.
  - For all $r$ in $0, 1, ...$, $C.rank() - 1$, if neither $C.static\_extent(r)$ nor $E.static\_extent(r)$ equals $dynamic\_extent$, then $C.static\_extent(r)$ equals $E.static\_extent(r)$ (if applicable).

• **Remarks:**
  - The functions will only access the triangle of $A$ specified by the Triangle argument $t$, and will assume for indices $i, j$ outside that triangle, that $A(j, i)$

2020-01-12
equals \( A(i,j) \).

- **Remarks:** \( C \) and \( E \) (if applicable) may refer to the same matrix. If so, then they must have the same layout.

The following requirements apply to all overloads of `symmetric_matrix_left_product`.

- **Requires:**
  - \( A.extent(1) \) equals \( B.extent(0) \),
  - \( A.extent(0) \) equals \( C.extent(0) \), and
  - \( B.extent(1) \) equals \( C.extent(1) \).

- **Mandates:**
  - If neither \( A.static_extent(1) \) nor \( B.static_extent(0) \) equals `dynamic_extent`, then \( A.static_extent(1) \) equals \( B.static_extent(0) \);
  - if neither \( A.static_extent(0) \) nor \( C.static_extent(0) \) equals `dynamic_extent`, then \( A.static_extent(0) \) equals \( C.static_extent(0) \); and
  - if neither \( B.static_extent(1) \) nor \( C.static_extent(1) \) equals `dynamic_extent`, then \( B.static_extent(1) \) equals \( C.static_extent(1) \).

The following requirements apply to all overloads of `symmetric_matrix_right_product`.

- **Requires:**
  - \( B.extent(1) \) equals \( A.extent(0) \),
  - \( B.extent(0) \) equals \( C.extent(0) \), and
  - \( A.extent(1) \) equals \( C.extent(1) \).

- **Mandates:**
  - If neither \( B.static_extent(1) \) nor \( A.static_extent(0) \) equals `dynamic_extent`, then \( B.static_extent(1) \) equals \( A.static_extent(0) \);
  - if neither \( B.static_extent(0) \) nor \( C.static_extent(0) \) equals `dynamic_extent`, then \( B.static_extent(0) \) equals \( C.static_extent(0) \); and
  - if neither \( A.static_extent(1) \) nor \( C.static_extent(1) \) equals `dynamic_extent`, then \( A.static_extent(1) \) equals \( C.static_extent(1) \).
class out_matrix_t>
void symmetric_matrix_left_product(
  in_matrix_1_t A,
  Triangle t,
  in_matrix_2_t B,
  out_matrix_t C);

template<class ExecutionPolicy,
  class in_matrix_1_t,
  class Triangle,
  class in_matrix_2_t,
  class out_matrix_t>
void symmetric_matrix_left_product(
  ExecutionPolicy&& exec,
  in_matrix_1_t A,
  Triangle t,
  in_matrix_2_t B,
  out_matrix_t C);

- **Constraints:** For \(i,j\) in the domain of \(C\), \(i,k\) in the domain of \(A\), and \(k,j\) in the domain of \(B\), the expression \(C(i,j) \gets A(i,k) \times B(k,j)\) is well formed.

- **Effects:** Assigns to the elements of the matrix \(C\) the product of the matrices \(A\) and \(B\).

Overwriting symmetric matrix-matrix right product [linalg.algs.blas3.symm.ov.right]

\[
\text{template<class in_matrix_1_t, class Triangle, class in_matrix_2_t, class out_matrix_t>}
void symmetric_matrix_right_product(
  in_matrix_1_t A,
  Triangle t,
  in_matrix_2_t B,
  out_matrix_t C);
\]

\[
\text{template<class ExecutionPolicy, class in_matrix_1_t, class Triangle, class in_matrix_2_t, class out_matrix_t>}
void symmetric_matrix_right_product(
  ExecutionPolicy&& exec,
  in_matrix_1_t A,
  Triangle t,
  in_matrix_2_t B,
  out_matrix_t C);
\]

- **Constraints:** For \(i,j\) in the domain of \(C\), \(i,k\) in the domain of \(B\), and \(k,j\) in the domain of \(A\), the expression \(C(i,j) \gets B(i,k) \times A(k,j)\) is well formed.

- **Effects:** Assigns to the elements of the matrix \(C\) the product of the matrices \(B\) and \(A\).

Updating symmetric matrix-matrix left product [linalg.algs.blas3.symm.up.left]
template<
class in_matrix_1_t,
class Triangle,
class in_matrix_2_t,
class in_matrix_3_t,
class out_matrix_t>
void symmetric_matrix_left_product(
in_matrix_1_t A,
Triangle t,
in_matrix_2_t B,
in_matrix_3_t E,
out_matrix_t C);

// Constraints: For i,j in the domain of C, i,k in the domain of A, and k,j in the domain of B, the expression C(i,j) += E(i,j) + A(i,k)*B(k,j) is well formed.

// Effects: assigns to the elements of the matrix C on output, the elementwise sum of E and the product of the matrices A and B.

Updating symmetric matrix-matrix right product [linalg.algs.blas3.symm.up.right]

template<
class in_matrix_1_t,
class Triangle,
class in_matrix_2_t,
class in_matrix_3_t,
class out_matrix_t>
void symmetric_matrix_right_product(
in_matrix_1_t A,
Triangle t,
in_matrix_2_t B,
in_matrix_3_t E,
out_matrix_t C);

// Constraints: For i,j in the domain of C, i,k in the domain of A, and k,j in the domain of B, the expression C(i,j) += E(i,j) + A(i,k)*B(k,j) is well formed.

// Effects: assigns to the elements of the matrix C on output, the elementwise sum of E and the product of the matrices A and B.
\texttt{in\_matrix\_3\_t E, out\_matrix\_t C};

- \textit{Constraints:} For \(i, j\) in the domain of \(C, i, k\) in the domain of \(B\), and \(k, j\) in the domain of \(A\), the expression \(C(i, j) \leftarrow E(i, j) + B(i, k) \times A(k, j)\) is well formed.

- \textit{Effects:} assigns to the elements of the matrix \(C\) on output, the elementwise sum of \(E\) and the product of the matrices \(B\) and \(A\).

\textbf{Hermitian matrix-matrix product [\texttt{inalg.algs.blas3.hemm]}]

\textit{[Note:}

These functions correspond to the BLAS function \texttt{xHEMM}.

Unlike the Hermitian rank-1 update functions, these functions assume that the input matrix -- not the output matrix -- is Hermitian.

\textit{--end note]}

The following requirements apply to all functions in this section.

- \textit{Requires:}
  - \texttt{A.extent(0)} equals \texttt{A.extent(1)}.
  - \texttt{C.extent(0)} equals \texttt{E.extent(0)} (if applicable).
  - \texttt{C.extent(1)} equals \texttt{E.extent(1)} (if applicable).

- \textit{Constraints:}
  - \texttt{in\_matrix\_1\_t} either has unique layout, or \texttt{layout\_blas\_packed} layout.
  - \texttt{in\_matrix\_2\_t, in\_matrix\_3\_t} (if applicable), and \texttt{out\_matrix\_t} have unique layout.
  - If \texttt{in\_matrix\_t} has \texttt{layout\_blas\_packed} layout, then the layout's \texttt{Triangle} template argument has the same type as the function's \texttt{Triangle} template argument.
  - \texttt{A.rank()} equals 2, \texttt{B.rank()} equals 2, \texttt{C.rank()} equals 2, and \texttt{E.rank()} (if applicable) equals 2.

- \textit{Mandates:}
  - If neither \texttt{A.static\_extent(0)} nor \texttt{A.static\_extent(1)} equals \texttt{dynamic\_extent}, then \texttt{A.static\_extent(0)} equals \texttt{A.static\_extent(1)}.
  - For all \(r\) in 0, 1, ..., \texttt{C.rank()} - 1, if neither \texttt{C.static\_extent(r)} nor \texttt{E.static\_extent(r)} equals \texttt{dynamic\_extent}, then \texttt{C.static\_extent(r)} equals \texttt{E.static\_extent(r)} (if applicable).

- \textit{Remarks:}
The functions will only access the triangle of \( A \) specified by the Triangle argument \( t \).

If `in_matrix_1_t::element_type` is `complex<RA>` for some RA, then the functions will assume for indices \( i,j \) outside that triangle, that \( A(j,i) \) equals \( \text{conj}(A(i,j)) \). Otherwise, the functions will assume that \( A(j,i) \) equals \( A(i,j) \).

\( C \) and \( E \) (if applicable) may refer to the same matrix. If so, then they must have the same layout.

The following requirements apply to all overloads of `hermitian_matrix_left_product`.

- **Requires:**
  
  - \( A.\text{extent}(1) \) equals \( B.\text{extent}(0) \),
  
  - \( A.\text{extent}(0) \) equals \( C.\text{extent}(0) \), and
  
  - \( B.\text{extent}(1) \) equals \( C.\text{extent}(1) \).

- **Mandates:**
  
  - If neither \( A.\text{static}_\text{extent}(1) \) nor \( B.\text{static}_\text{extent}(0) \) equals dynamic extent, then \( A.\text{static}_\text{extent}(1) \) equals \( B.\text{static}_\text{extent}(0) \);
  
  - if neither \( A.\text{static}_\text{extent}(0) \) nor \( C.\text{static}_\text{extent}(0) \) equals dynamic extent, then \( A.\text{static}_\text{extent}(0) \) equals \( C.\text{static}_\text{extent}(0) \); and
  
  - if neither \( B.\text{static}_\text{extent}(1) \) nor \( C.\text{static}_\text{extent}(1) \) equals dynamic extent, then \( B.\text{static}_\text{extent}(1) \) equals \( C.\text{static}_\text{extent}(1) \).

The following requirements apply to all overloads of `hermitian_matrix_right_product`.

- **Requires:**
  
  - \( B.\text{extent}(1) \) equals \( A.\text{extent}(0) \),
  
  - \( B.\text{extent}(0) \) equals \( C.\text{extent}(0) \), and
  
  - \( A.\text{extent}(1) \) equals \( C.\text{extent}(1) \).

- **Mandates:**
  
  - If neither \( B.\text{static}_\text{extent}(1) \) nor \( A.\text{static}_\text{extent}(0) \) equals dynamic extent, then \( B.\text{static}_\text{extent}(1) \) equals \( A.\text{static}_\text{extent}(0) \);
  
  - if neither \( B.\text{static}_\text{extent}(0) \) nor \( C.\text{static}_\text{extent}(0) \) equals dynamic extent, then \( B.\text{static}_\text{extent}(0) \) equals \( C.\text{static}_\text{extent}(0) \); and
  
  - if neither \( A.\text{static}_\text{extent}(1) \) nor \( C.\text{static}_\text{extent}(1) \) equals dynamic extent, then \( A.\text{static}_\text{extent}(1) \) equals \( C.\text{static}_\text{extent}(1) \).


- **Constraints**: For \(i,j\) in the domain of \(C\), \(i,k\) in the domain of \(A\), and \(k,j\) in the domain of \(B\), the expression \(C(i,j) \leftarrow A(i,k) \times B(k,j)\) is well formed.

- **Effects**: Assigns to the elements of the matrix \(C\) the product of the matrices \(A\) and \(B\).

---

---

- **Constraints**: For \(i,j\) in the domain of \(C\), \(i,k\) in the domain of \(B\), and \(k,j\) in the domain of \(A\), the expression \(C(i,j) \leftarrow B(i,k) \times A(k,j)\) is well formed.
• **Effects:** Assigns to the elements of the matrix $C$ the product of the matrices $B$ and $A$.

**Updating Hermitian matrix–matrix left product [linalg.algs.blas3.hemm.up.left]**

```cpp
template<class in_matrix_1_t,
    class Triangle,
    class in_matrix_2_t,
    class in_matrix_3_t,
    class out_matrix_t>
void hermitian_matrix_left_product(
    in_matrix_1_t A,
    Triangle t,
    in_matrix_2_t B,
    in_matrix_3_t E,
    out_matrix_t C);

template<class ExecutionPolicy,
    class in_matrix_1_t,
    class Triangle,
    class in_matrix_2_t,
    class in_matrix_3_t,
    class out_matrix_t>
void hermitian_matrix_left_product(
    ExecutionPolicy&& exec,
    in_matrix_1_t A,
    Triangle t,
    in_matrix_2_t B,
    in_matrix_3_t E,
    out_matrix_t C);
```

• **Constraints:** For $i, j$ in the domain of $C$, $i, k$ in the domain of $A$, and $k, j$ in the domain of $B$, the expression $C(i, j) += E(i, j) + A(i, k) * B(k, j)$ is well formed.

• **Effects:** Assigns to the elements of the matrix $C$ on output, the elementwise sum of $E$ and the product of the matrices $A$ and $B$.

**Updating Hermitian matrix–matrix right product [linalg.algs.blas3.hemm.up.right]**

```cpp
template<class in_matrix_1_t,
    class Triangle,
    class in_matrix_2_t,
    class in_matrix_3_t,
    class out_matrix_t>
void hermitian_matrix_right_product(
    in_matrix_1_t A,
    Triangle t,
    in_matrix_2_t B,
    in_matrix_3_t E,
    out_matrix_t C);

template<class ExecutionPolicy,
    class in_matrix_1_t,
    class Triangle,
    class in_matrix_2_t,
    class in_matrix_3_t,
    class in_matrix_3_t,
    class out_matrix_t>
void hermitian_matrix_right_product(
    ExecutionPolicy&& exec,
    in_matrix_1_t A,
    Triangle t,
    in_matrix_2_t B,
    in_matrix_3_t E,
    out_matrix_t C);
```


```c
class out_matrix_t>
void hermitian_matrix_right_product(
    ExecutionPolicy & exec,
    in_matrix_1_t A,
    Triangle t,
    in_matrix_2_t B,
    in_matrix_3_t E,
    out_matrix_t C);
```

- **Constraints:** For \( i, j \) in the domain of \( C, i, k \) in the domain of \( B \), and \( k, j \) in the domain of \( A \), the expression \( C(i, j) \leftarrow E(i, j) + B(i, k) \ast A(k, j) \) is well formed.

- **Effects:** Assigns to the elements of the matrix \( C \) on output, the elementwise sum of \( E \) and the product of the matrices \( B \) and \( A \).

Triangular matrix-matrix product [linalg.algs.blas3.trmm]

*[Note: These functions correspond to the BLAS function \texttt{xTRMM}. --end note]*

The following requirements apply to all functions in this section.

- **Requires:**
  - \( A.\text{extent}(0) \) equals \( A.\text{extent}(1) \).
  - \( C.\text{extent}(0) \) equals \( E.\text{extent}(0) \) (if applicable).
  - \( C.\text{extent}(1) \) equals \( E.\text{extent}(1) \) (if applicable).

- **Constraints:**
  - \( \text{in\_matrix\_1\_t} \) either has unique layout, or \texttt{layout\_blas\_packed} layout.
  - \( \text{in\_matrix\_2\_t, in\_matrix\_3\_t} \) (if applicable), \( \text{out\_matrix\_t} \), and \( \text{inout\_matrix\_t} \) (if applicable) have unique layout.
  - If \( \text{in\_matrix\_1\_t} \) has \texttt{layout\_blas\_packed} layout, then the layout's \texttt{Triangle} template argument has the same type as the function's \texttt{Triangle} template argument.
  - \( A.\text{rank()} \) equals 2, \( B.\text{rank()} \) equals 2, \( C.\text{rank()} \) equals 2, and \( E.\text{rank()} \) (if applicable) equals 2.

- **Mandates:**
  - If neither \( A.\text{static\_extent}(0) \) nor \( A.\text{static\_extent}(1) \) equals \texttt{dynamic\_extent}, then \( A.\text{static\_extent}(0) \) equals \( A.\text{static\_extent}(1) \).
  - For all \( r \) in \( 0, 1, ... \), \( C.\text{rank()} - 1 \), if neither \( C.\text{static\_extent}(r) \) nor \( E.\text{static\_extent}(r) \) equals \texttt{dynamic\_extent}, then \( C.\text{static\_extent}(r) \) equals \( E.\text{static\_extent}(r) \) (if applicable).

- **Remarks:**
The functions will only access the triangle of A specified by the Triangle argument t.

If the DiagonalStorage template argument has type implicit_unit_diagonal_t, then the functions will not access the diagonal of A, and will assume that that the diagonal elements of A all equal one. [Note: This does not imply that the function needs to be able to form an element_type value equal to one. --*end note]

C and E (if applicable) may refer to the same matrix. If so, then they must have the same layout.

The following requirements apply to all overloads of triangular_matrix_left_product.

- **Requires:**
  - A.extent(1) equals B.extent(0) (if applicable),
  - A.extent(0) equals C.extent(0), and
  - B.extent(1) equals C.extent(1) (if applicable).

- **Mandates:**
  - If neither A.static_extent(1) nor B.static_extent(0) equals dynamic_extent, then A.static_extent(1) equals B.static_extent(0) (if applicable);
  - if neither A.static_extent(0) nor C.static_extent(0) equals dynamic_extent, then A.static_extent(0) equals C.static_extent(0);
  - if neither B.static_extent(1) nor C.static_extent(1) equals dynamic_extent, then B.static_extent(1) equals C.static_extent(1) (if applicable).

The following requirements apply to all overloads of triangular_matrix_right_product.

- **Requires:**
  - B.extent(1) equals A.extent(0) (if applicable),
  - B.extent(0) equals C.extent(0) (if applicable), and
  - A.extent(1) equals C.extent(1).

- **Mandates:**
  - If neither B.static_extent(1) nor A.static_extent(0) equals dynamic_extent, then B.static_extent(1) equals A.static_extent(0) (if applicable);
  - if neither B.static_extent(0) nor C.static_extent(0) equals dynamic_extent, then B.static_extent(0) equals C.static_extent(0) (if applicable).
Overwriting triangular matrix-matrix left product [linalg.algs.blas3.tmm.ov.left]

Not-in-place overwriting triangular matrix-matrix left product

```cpp
template<
class in_matrix_1_t,
class Triangle,
class DiagonalStorage,
class in_matrix_2_t,
class out_matrix_t>
void triangular_matrix_left_product(
in_matrix_1_t A,
Triangle t,
DiagonalStorage d,
in_matrix_2_t B,
out_matrix_t C);
template<
class ExecutionPolicy,
class in_matrix_1_t,
class Triangle,
class DiagonalStorage,
class in_matrix_2_t,
class out_matrix_t>
void triangular_matrix_left_product(
ExecutionPolicy&& exec,
in_matrix_1_t A,
Triangle t,
DiagonalStorage d,
in_matrix_2_t B,
out_matrix_t C);
```

- **Constraints:** For $i, j$ in the domain of $C$, $i, k$ in the domain of $A$, and $k, j$ in the domain of $B$, the expression $C(i, j) += A(i, k) \times B(k, j)$ is well formed.

- **Effects:** Assigns to the elements of the matrix $C$ the product of the matrices $A$ and $B$.

In-place overwriting triangular matrix-matrix left product

```cpp
template<
class in_matrix_1_t,
class Triangle,
class DiagonalStorage,
class inout_matrix_t>
void triangular_matrix_left_product(
in_matrix_1_t A,
Triangle t,
DiagonalStorage d,
inout_matrix_t C);
```

- **Requires:** $A.extent(1)$ equals $C.extent(0)$.
• **Constraints:** For \( i, j \) in the domain of \( C \), and \( i, k \) in the domain of \( A \), the expression \( C(i,j) += A(i,k) * C(k,j) \) is well formed.

• **Mandates:** If neither \( A.\text{static_extent}(1) \) nor \( C.\text{static_extent}(0) \) equals 
  \( \text{dynamic_extent} \), then \( A.\text{static_extent}(1) \) equals \( C.\text{static_extent}(0) \).

• **Effects:** Overwrites \( C \) on output with the product of the matrices \( A \) and \( C \) (on input).

Overwriting triangular matrix-matrix right product ([linalg.algs.blas3.trmm.ov])

Not-in-place overwriting triangular matrix-matrix right product

```cpp
template<class in_matrix_1_t,
         class Triangle,
         class DiagonalStorage,
         class in_matrix_2_t,
         class out_matrix_t>
void triangular_matrix_right_product(
    in_matrix_1_t A,
    Triangle t,
    DiagonalStorage d,
    in_matrix_2_t B,
    out_matrix_t C);

template<class ExecutionPolicy,
          class in_matrix_1_t,
          class Triangle,
          class DiagonalStorage,
          class in_matrix_2_t,
          class out_matrix_t>
void triangular_matrix_right_product(  
    ExecutionPolicy&& exec,
    in_matrix_1_t A,
    Triangle t,
    DiagonalStorage d,
    in_matrix_2_t B,
    out_matrix_t C);
```

• **Constraints:** For \( i, j \) in the domain of \( C \), \( i, k \) in the domain of \( B \), and \( k, j \) in the domain of 
  \( A \), the expression \( C(i,j) += B(i,k) * A(k,j) \) is well formed.

• **Effects:** Assigns to the elements of the matrix \( C \) the product of the matrices \( B \) and \( A \).

In-place overwriting triangular matrix-matrix right product

```cpp
template<class in_matrix_1_t,
         class Triangle,
         class DiagonalStorage,
         class inout_matrix_t>
void triangular_matrix_right_product(
    in_matrix_1_t A,
    Triangle t,
    DiagonalStorage d,
    inout_matrix_t C);
```
- **Requires**: `C.extent(1)` equals `A.extent(0)`.

- **Constraints**: For `i, j` and `i, k` in the domain of `C`, and `k, j` in the domain of `A`, the expression `C(i, j) += C(i, k) * A(k, j)` is well formed.

- **Mandates**: If neither `C.static_extent(1)` nor `A.static_extent(0)` equals `dynamic_extent`, then `C.static_extent(1)` equals `A.static_extent(0)`.

- **Effects**: Overwrites `C` on output with the product of the matrices `C` (on input) and `A`.

---

**Updating triangular matrix-matrix left product** [linalg.algs.blas3.trmm.left]

```cpp
template<class in_matrix_1_t,
         class Triangle,
         class DiagonalStorage,
         class in_matrix_2_t,
         class in_matrix_3_t,
         class out_matrix_t>
void triangular_matrix_left_product(
   in_matrix_1_t A,
   Triangle t,
   DiagonalStorage d,
   in_matrix_2_t B,
   in_matrix_3_t E,
   out_matrix_t C);

```

- **Constraints**: For `i, j` in the domain of `C`, `i, k` in the domain of `A`, and `k, j` in the domain of `B`, the expression `C(i, j) += E(i, j) + A(i, k) * B(k, j)` is well formed.

- **Effects**: Assigns to the elements of the matrix `C` on output, the elementwise sum of `E` and the product of the matrices `A` and `B`.

---

**Updating triangular matrix-matrix right product** [linalg.algs.blas3.trmm.right]

```cpp
template<class in_matrix_1_t,
         class Triangle,
         class DiagonalStorage,
         class in_matrix_2_t,
```
void triangular_matrix_right_product(
  in_matrix_1_t A,
  Triangle t,
  DiagonalStorage d,
  in_matrix_2_t B,
  in_matrix_3_t E,
  out_matrix_t C);

// Constraints: For $i, j$ in the domain of $C$, $i, k$ in the domain of $B$, and $k, j$ in the domain of $A$, the expression $C(i, j) += E(i, j) + B(i, k) \cdot A(k, j)$ is well formed.

// Effects: Assigns to the elements of the matrix $C$ on output, the elementwise sum of $E$ and the product of the matrices $B$ and $A$.

Rank-k update of a symmetric or Hermitian matrix [linalg.blas3.rank-k]

[Note: Users can achieve the effect of the TRANS argument of these BLAS functions, by applying transpose_view or conjugate_transpose_view to the input matrix. --end note]

Rank-k symmetric matrix update [linalg.blas3.rank-k.syrk]
template<class T,
class in_matrix_1_t,
class inout_matrix_t,
class Triangle>
void symmetric_matrix_rank_k_update(
T alpha,
in_matrix_1_t A,
inout_matrix_t C,
Triangle t);

template<class ExecutionPolicy,
class T,
class in_matrix_1_t,
class inout_matrix_t,
class Triangle>
void symmetric_matrix_rank_k_update(
    ExecutionPolicy&& exec,
    T alpha,
in_matrix_1_t A,
inout_matrix_t C,
    Triangle t);

[Note:

These functions correspond to the BLAS function xSYRK.

They take an optional scaling factor alpha, because it would be impossible to express the
update $C = C - A A^T$ otherwise.

--end note]

- Requires:
  - $A.extent(0)$ equals $C.extent(0)$.
  - $C.extent(0)$ equals $C.extent(1)$.

- Constraints:
  - $A.rank()$ equals 2 and $C.rank()$ equals 2.
  - C either has unique layout, or layout_blas_packed layout.
  - If $C$ has layout_blas_packed layout, then the layout's Triangle template argument has the same type as the function's Triangle template argument.
  - For $i, j$ in the domain of $C$, and $i, k$ and $k, j$ in the domain of $A$, the expression $C(i, j) := A(i, k) * A(j, k)$ is well formed.
  - For $i, j$ in the domain of $C$, and $i, k$ and $k, j$ in the domain of $A$, the expression $C(i, j) := \alpha * A(i, k) * A(j, k)$ is well formed (if applicable).

- Mandates:
  - If neither $A.static_extent(0)$ nor $C.static_extent(0)$ equals dynamic_extent, then $A.static_extent(0)$ equals $C.static_extent(0)$.  

- If neither \(C\text{.static_extent}(0)\) nor \(C\text{.static_extent}(1)\) equals \(\text{dynamic_extent}\), then \(C\text{.static_extent}(0)\) equals \(C\text{.static_extent}(1)\).

- **Effects:**
  - Overloads without \(\alpha\) assign to \(C\) on output, the elementwise sum of \(C\) on input with (the matrix product of \(A\) and the nonconjugated transpose of \(A\)).
  - Overloads with \(\alpha\) assign to \(C\) on output, the elementwise sum of \(C\) on input with \(\alpha\) times (the matrix product of \(A\) and the nonconjugated transpose of \(A\)).

- **Remarks:** The functions will only access the triangle of \(C\) specified by the Triangle argument \(t\), and will assume for indices \(i,j\) outside that triangle, that \(C(j,i)\) equals \(C(i,j)\).

---

**Rank-k symmetric matrix update [linalg.alg.blas3.rank-k.herk]**

```cpp
template<
class in_matrix_1_t,
class inout_matrix_t,
class Triangle>
void hermitian_matrix_rank_k_update(
  in_matrix_1_t A,
  inout_matrix_t C,
  Triangle t);

template<
class ExecutionPolicy,
class in_matrix_1_t,
class inout_matrix_t,
class Triangle>
void hermitian_matrix_rank_k_update(
  ExecutionPolicy&& exec,
  in_matrix_1_t A,
  inout_matrix_t C,
  Triangle t);

template<
class T,
class in_matrix_1_t,
class inout_matrix_t,
class Triangle>
void hermitian_matrix_rank_k_update(
  T alpha,
  in_matrix_1_t A,
  inout_matrix_t C,
  Triangle t);

template<
class ExecutionPolicy,
class T,
class in_matrix_1_t,
class inout_matrix_t,
class Triangle>
void hermitian_matrix_rank_k_update(
  ExecutionPolicy&& exec,
  T alpha,
  in_matrix_1_t A,
  inout_matrix_t C,
  Triangle t);
```
These functions correspond to the BLAS function \texttt{xHERK}.

They take an optional scaling factor \texttt{alpha}, because it would be impossible to express the updates \( C = C - A A^T \) or \( C = C - A A^H \) otherwise.

--end note]

- **Requires:**
  - \texttt{A.extent(0)} equals \texttt{C.extent(0)}.
  - \texttt{C.extent(0)} equals \texttt{C.extent(1)}.

- **Constraints:**
  - \texttt{A.rank()} equals 2 and \texttt{C.rank()} equals 2.
  - \texttt{C} either has unique layout, or \texttt{layout_blas_packed} layout.
  - If \texttt{C} has \texttt{layout_blas_packed} layout, then the layout's \texttt{Triangle} template argument has the same type as the function's \texttt{Triangle} template argument.
  - For overloads without \texttt{alpha}: For \( i,j \) in the domain of \( C \), and \( i,k \) and \( k,i \) in the domain of \( A \),
    - if \texttt{in\_matrix\_1\_t::element\_type\_is complex<\texttt{R}>} for some \( \texttt{R} \), then the expression \( C(i,j) \pm A(i,k)\ast\text{conj}(A(j,k)) \) is well formed;
    - else, the expression \( C(i,j) \pm A(i,k)\ast A(j,k) \) is well formed.
  - For overloads with \texttt{alpha}: For \( i,j \) in the domain of \( C \), and \( i,k \) and \( k,i \) in the domain of \( A \),
    - if \texttt{in\_matrix\_1\_t::element\_type\_is complex<\texttt{R}>} for some \( \texttt{R} \), then the expression \( C(i,j) \pm \text{alpha}\ast A(i,k)\ast\text{conj}(A(j,k)) \) is well formed;
    - else, the expression \( C(i,j) \pm \text{alpha}\ast A(i,k)\ast A(j,k) \) is well formed.

- **Mandates:**
  - If neither \texttt{A.static\_extent(0)} nor \texttt{C.static\_extent(0)} equals \texttt{dynamic\_extent}, then \texttt{A.static\_extent(0)} equals \texttt{C.static\_extent(0)}.
  - If neither \texttt{C.static\_extent(0)} nor \texttt{C.static\_extent(1)} equals \texttt{dynamic\_extent}, then \texttt{C.static\_extent(0)} equals \texttt{C.static\_extent(1)}.

- **Effects:**
  - Overloads without \texttt{alpha} assign to \( C \) on output, the elementwise sum of \( C \) on input with (the matrix product of \( A \) and the conjugated transpose of \( A \)).
  - Overloads with \texttt{alpha} assign to \( C \) on output, the elementwise sum of \( C \) on input with \( \text{alpha} \) times (the matrix product of \( A \) and the conjugated transpose of \( A \)).
Remarks: The functions will only access the triangle of $C$ specified by the $\text{Triangle}$ argument $t$, and will assume for indices $i,j$ outside that triangle, that $C(i,j)$ equals $C(j,i)$.

Rank-2k update of a symmetric or Hermitian matrix [linalg.blas3.rank2k]

[Note: Users can achieve the effect of the $\text{TRANS}$ argument of these BLAS functions, by applying $\text{transpose_view}$ or $\text{conjugate_transpose_view}$ to the input matrices. --end note]

Rank-2k symmetric matrix update [linalg.blas3.rank2k.sy2k]

```c++
template<
class in_matrix_1_t,
class in_matrix_2_t,
class inout_matrix_t,
class Triangle>
void symmetric_matrix_rank_2k_update(
in_matrix_1_t A,
in_matrix_2_t B,
inout_matrix_t C,
Triangle t);
```

[Note: These functions correspond to the BLAS function $\text{xSYR2K}$. The BLAS "quick reference" has a typo; the "ALPHA" argument of $\text{CSYR2K}$ and $\text{ZSYR2K}$ should not be conjugated. --end note]

- Requires:
  - $\text{A.extent(0)}$ equals $\text{C.extent(0)}$.
  - $\text{B.extent(1)}$ equals $\text{C.extent(0)}$.
  - $\text{C.extent(0)}$ equals $\text{C.extent(1)}$.

- Constraints:
  - $\text{A.rank()}$ equals 2, $\text{B.rank()}$ equals 2, and $\text{C.rank()}$ equals 2.
  - $\text{C}$ either has unique layout, or $\text{layout_blas_packed}$ layout.
  - If $\text{C}$ has $\text{layout_blas_packed}$ layout, then the layout's $\text{Triangle}$ template argument has the same type as the function's $\text{Triangle}$ template argument.
For \( i, j \) in the domain of \( C, i, k \) and \( k, i \) in the domain of \( A \), and \( j, k \) and \( k, j \) in the domain of \( B \), the expression \( C(i,j) \rightarrow A(i,k)B(j,k) + B(i,k)A(j,k) \) is well formed.

- **Mandates:**
  - If neither \( A\.static\.extent(0) \) nor \( C\.static\.extent(0) \) equals \( \text{dynamic\.extent} \), then \( A\.static\.extent(0) = C\.static\.extent(0) \).
  - If neither \( B\.static\.extent(1) \) nor \( C\.static\.extent(0) \) equals \( \text{dynamic\.extent} \), then \( B\.static\.extent(1) = C\.static\.extent(0) \).
  - If neither \( C\.static\.extent(0) \) nor \( C\.static\.extent(1) \) equals \( \text{dynamic\.extent} \), then \( C\.static\.extent(0) = C\.static\.extent(1) \).

- **Effects:** Assigns to \( C \) on output, the elementwise sum of \( C \) on input with (the matrix product of \( A \) and the nonconjugated transpose of \( B \)) and (the matrix product of \( B \) and the nonconjugated transpose of \( A \)).

- **Remarks:** The functions will only access the triangle of \( C \) specified by the\ Triangle argument \( t \), and will assume for indices \( i, j \) outside that triangle, that \( C(j,i) = C(i,j) \).

---

**Rank-2k Hermitian matrix update** [linalg.alg.blas3.rank2k.her2k]

```cpp
template<class in_matrix_1_t,
class in_matrix_2_t,
class inout_matrix_t,
class Triangle>
void hermitian_matrix_rank_2k_update(
in_matrix_1_t A,
in_matrix_2_t B,
inout_matrix_t C,
Triangle t);

template<class ExecutionPolicy,
class in_matrix_1_t,
class in_matrix_2_t,
class inout_matrix_t,
class Triangle>
void hermitian_matrix_rank_2k_update(
ExecutionPolicy&\& exec,
in_matrix_1_t A,
in_matrix_2_t B,
inout_matrix_t C,
Triangle t);
```

**[Note:** These functions correspond to the BLAS function \texttt{xHER2K}. --end note]**

- **Requires:**
  - \( A\.extent(0) = C\.extent(0) \).
\* B.extent(1) equals C.extent(0).
\* C.extent(0) equals C.extent(1).

- **Constraints:**
  \* A.rank() equals 2, B.rank() equals 2, and C.rank() equals 2.
  \* C either has unique layout, or layout_blas_packed layout.
  \* If C has layout_blas_packed layout, then the layout's Triangle template argument has the same type as the function's Triangle template argument.
  \* For i, j in the domain of C, i, k and k, i in the domain of A, and j, k and k, j in the domain of B,
    - if in_matrix_1_t::element_type is complex<RA> for some RA, then
      - if in_matrix_2_t::element_type is complex<RB> for some RB, then the expression C(i, j) += A(i, k)*conj(B(j, k)) + B(i, k)*conj(A(j, k)) is well formed;
    - else, the expression C(i, j) += A(i, k)*conj(B(j, k)) + B(i, k)*A(j, k) is well formed;
    - else,
      - if in_matrix_2_t::element_type is complex<RB> for some RB, then the expression C(i, j) += A(i, k)*conj(B(j, k)) + B(i, k)*A(j, k) is well formed;
      - else, the expression C(i, j) += A(i, k)*B(j, k) + B(i, k)*A(j, k) is well formed.

- **Mandates:**
  \* If neither A.static_extent(0) nor C.static_extent(0) equals dynamic_extent, then A.static_extent(0) equals C.static_extent(0).
  \* If neither B.static_extent(1) nor C.static_extent(0) equals dynamic_extent, then B.static_extent(1) equals C.static_extent(0).
  \* If neither C.static_extent(0) nor C.static_extent(1) equals dynamic_extent, then C.static_extent(0) equals C.static_extent(1).

- **Effects:** Assigns to C on output, the elementwise sum of C on input with (the matrix product of A and the conjugate transpose of B) and (the matrix product of B and the conjugate transpose of A.)

- **Remarks:**
  \* The functions will only access the triangle of C specified by the Triangle argument t.
If \texttt{inout_matrix_t::element_type} is \texttt{complex<RC>} for some \texttt{RC}, then the functions will assume for indices \(i,j\) outside that triangle, that \(C(j,i)\) equals \(\text{conj}(C(i,j))\). Otherwise, the functions will assume that \(C(j,i)\) equals \(C(i,j)\).

**Solve multiple triangular linear systems** \([\text{linear.alg.blas3.trsm}]\)

[Note: These functions correspond to the BLAS function \texttt{xTRSM}. The Reference BLAS does not have a \texttt{xTPSM} function. --end note]

The following requirements apply to all functions in this section.

- **Requires:**
  - For all \(r\) in \(0, 1, \ldots, B.\text{rank()} - 1\), \(X.\text{extent}(r)\) equals \(B.\text{extent}(r)\) (if applicable).
  - \(A.\text{extent}(0)\) equals \(A.\text{extent}(1)\).

- **Constraints:**
  - \(A.\text{rank}()\) equals 2 and \(B.\text{rank}()\) equals 2.
  - \(X.\text{rank}()\) equals 2 (if applicable).
  - \(\text{in\_matrix\_1\_t}\) either has unique layout, or \texttt{layout\_blas\_packed} layout.
  - \(\text{in\_matrix\_2\_t}\) has unique layout (if applicable).
  - \(\text{out\_matrix\_t}\) has unique layout.
  - \(\text{inout\_matrix\_t}\) has unique layout (if applicable).
  - If \(r,j\) is in the domain of \(X\) and \(B\), then the expression \(X(r,j) = B(r,j)\) is well formed (if applicable).
  - If \texttt{DiagonalStorage} is \texttt{explicit_diagonal\_t}, and \(i,j\) is in the domain of \(X\), then the expression \(X(i,j) /= A(i,i)\) is well formed (if applicable).

- **Mandates:**
  - For all \(r\) in \(0, 1, \ldots, X.\text{rank()} - 1\), if neither \(X.\text{static\_extent}(r)\) nor \(B.\text{static\_extent}(r)\) equals \texttt{dynamic\_extent}, then \(X.\text{static\_extent}(r)\) equals \(B.\text{static\_extent}(r)\) (if applicable).
  - If neither \(A.\text{static\_extent}(0)\) nor \(A.\text{static\_extent}(1)\) equals \texttt{dynamic\_extent}, then \(A.\text{static\_extent}(0)\) equals \(A.\text{static\_extent}(1)\).

- **Remarks:**
  - The functions will only access the triangle of \(A\) specified by the \texttt{Triangle} argument \(t\).
  - If the \texttt{DiagonalStorage} template argument has type \texttt{implicit\_unit\_diagonal\_t}, then the functions will not access the diagonal of
A, and will assume that the diagonal elements of A all equal one. [Note: This
does not imply that the function needs to be able to form an element_type value
equal to one. --*end note]

Solve multiple triangular linear systems with triangular matrix on the left [linalg.alg.blas3.trsm.left]

Not-in-place multiple triangular systems left solve

```cpp
template<class in_matrix_1_t,
class Triangle,
class DiagonalStorage,
class in_matrix_2_t,
class out_matrix_t>
void triangular_matrix_matrix_left_solve(
in_matrix_1_t A,
Triangle t,
DiagonalStorage d,
in_matrix_2_t B,
out_matrix_t X);
```

- **Requires:** A.extent(1) equals B.extent(0).

- **Constraints:** If i, j and i, k are in the domain of X, then the expression X(i, j) -=
  A(i, k) * X(k, j) is well formed.

- **Mandates:** If neither A.static_extent(1) nor B.static_extent(0) equals
dynamic_extent, then A.static_extent(1) equals B.static_extent(0).

- **Effects:** Assigns to the elements of X the result of solving the triangular linear system(s)
  AX=B for X.

In-place multiple triangular systems left solve

```cpp
template<class in_matrix_1_t,
class Triangle,
class DiagonalStorage,
class inout_matrix_t>
void triangular_matrix_matrix_left_solve(
in_matrix_1_t A,
Triangle t,
```

---
DiagonalStorage d,
inout_matrix_t B);

[Note:

This in-place version of the function intentionally lacks an overload taking an ExecutionPolicy&&, because it is not possible to parallelize in-place triangular solve for an arbitrary ExecutionPolicy.

This algorithm makes it possible to compute factorizations like Cholesky and LU in place.

---end note]

- **Requires:** A.extent(1) equals B.extent(0).

- **Constraints:**
  - If DiagonalStorage is explicit_diagonal_t, and i,j is in the domain of B, then the expression B(i,j) /= A(i,i) is well formed (if applicable).
  - If i,j and i,k are in the domain of X, then the expression B(i,j) -= A(i,k) * B(k,j) is well formed.

- **Mandates:** If neither A.static_extent(1) nor B.static_extent(0) equals dynamic_extent, then A.static_extent(1) equals B.static_extent(0).

- **Effects:** Overwrites B with the result of solving the triangular linear system(s) AX=B for X.

Solve multiple triangular linear systems with triangular matrix on the right [linalg.blas3.trsm.right]

Not-in-place multiple triangular systems right solve

```cpp
template<class in_matrix_1_t,  
class Triangle,  
class DiagonalStorage,  
class in_matrix_2_t,  
class out_matrix_t>  
void triangular_matrix_matrix_right_solve(  
in_matrix_1_t A,  
Triangle t,  
DiagonalStorage d,  
in_matrix_2_t B,  
out_matrix_t X);

template<class ExecutionPolicy,  
class in_matrix_1_t,  
class Triangle,  
class DiagonalStorage,  
class in_matrix_2_t,  
class out_matrix_t>  
void triangular_matrix_matrix_right_solve(  
ExecutionPolicy&& exec,  
in_matrix_1_t A,  
Triangle t,  
DiagonalStorage d,
```
\texttt{\texttt{in\_matrix\_2\_t} \texttt{B,}}
\texttt{\texttt{out\_matrix\_t} \texttt{X};}

- \textit{Requires:} A.extent(1) equals B.extent(1).

- \textit{Constraints:} If \(i, j\) and \(i, k\) are in the domain of \(X\), then the expression \(X(i, j) \leftarrow X(i, k) \times A(k, j)\) is well formed.

- \textit{Mandates:} If neither A.static_extent(1) nor B.static_extent(1) equals dynamic_extent, then A.static_extent(1) equals B.static_extent(1).

- \textit{Effects:} Assigns to the elements of \(X\) the result of solving the triangular linear system(s) \(XA=B\) for \(X\).

\textbf{In-place multiple triangular systems right solve}

\begin{verbatim}
template<class in_matrix_1_t,
    class Triangle,
    class DiagonalStorage,
    class inout_matrix_t>
void triangular_matrix_matrix_right_solve(
    in_matrix_1_t A,
    Triangle t,
    DiagonalStorage d,
    inout_matrix_t B);
\end{verbatim}

\textit{[Note:}

This in-place version of the function intentionally lacks an overload taking an \texttt{ExecutionPolicy\&\&}, because it is not possible to parallelize in-place triangular solve for any \texttt{ExecutionPolicy}.

This algorithm makes it possible to compute factorizations like Cholesky and LU in place.

\textit{--end note]}

- \textit{Requires:} A.extent(1) equals B.extent(1).

- \textit{Constraints:}
  \begin{itemize}
  \item If DiagonalStorage is \texttt{explicit\_diagonal\_t}, and \(i, j\) is in the domain of \(B\),
    then the expression \(B(i, j) \leftarrow A(i, i)\) is well formed (if applicable).
  \item If \(i, j\) and \(i, k\) are in the domain of \(X\), then the expression \(B(i, j) \leftarrow B(i, k) \times A(k, j)\) is well formed.
  \end{itemize}

- \textit{Mandates:} If neither A.static_extent(1) nor B.static_extent(1) equals dynamic_extent, then A.static_extent(1) equals B.static_extent(1).

- \textit{Effects:} Overwrites \(B\) with the result of solving the triangular linear system(s) \(XA=B\) for \(X\).

\textbf{Examples}
Cholesky factorization

This example shows how to compute the Cholesky factorization of a real symmetric positive definite matrix $A$ stored as a basic_mdspan with a unique non-packed layout. The algorithm imitates DPOTRF2 in LAPACK 3.9.0. If Triangle is upper_triangle_t, then it computes the Cholesky factorization $A = U^T U$. Otherwise, it computes the Cholesky factorization $A = L L^T$. The function returns 0 if success, else k+1 if row/column k has a zero or NaN (not a number) diagonal entry.

```cpp
#include <linalg>
#include <cmath>

template<class inout_matrix_t, class Triangle>
int cholesky_factor(inout_matrix_t A, Triangle t) {
  using element_type = typename inout_matrix_t::element_type;
  constexpr element_type ZERO {};
  constexpr element_type ONE (1.0);
  const ptdiff_t n = A.extent(0);

  if (n == 0) {
    return 0;
  }
  else if (n == 1) {
    if (A(0,0) <= ZERO || isnan(A(0,0))) {
      return 1;
    }
    A(0,0) = sqrt(A(0,0));
  }
  else {
    // Partition $A$ into $[A_{11}, A_{12},
    // A_{21}, A_{22}]$,
    // where $A_{21}$ is the transpose of $A_{12}$. 
    const ptdiff_t n1 = n / 2;
    const ptdiff_t n2 = n - n1;
    auto A11 = subspan(A, pair{0, n1}, pair{0, n1});
    auto A22 = subspan(A, pair{n1, n}, pair{n1, n});

    // Factor $A_{11}$
    const int info1 = cholesky_factor(A11, t);
    if (info1 != 0) {
      return info1;
    }

    if constexpr (std::is_same_v<Triangle, upper_triangle_t>) {
      // Update and scale $A_{12}$
      auto A12 = subspan(A, pair{0, n1}, pair{n1, n});
      triangular_matrix_matrix_left_solve(transpose_view(A11),
              upper_triangle, explicit_diagonal, A12);
      // $A_{22} = A_{22} - A_{12}^T * A_{12}$
      symmetric_matrix_rank_k_update(-ONE, transpose_view(A12),
                                      A22, t);
    } else {
      //
    }
  }
}
```
// Compute the Cholesky factorization A = L * L^T
// Update and scale A21
auto A21 = subspace(A, pair{n1, n}, pair{0, n1});
triangular_matrix_matrix_right_solve(transpose_view(A11), lower_triangle, explicit_diagonal, A21);
// A22 = A22 - A21 * A21^T
symmetric_matrix_rank_k_update(-ONE, A21, A22, t);
}

// Factor A22
const int info2 = cholesky_factor(A22, t);
if (info2 != 0) {
    return info2 + n1;
}
}

Solve linear system using Cholesky factorization

This example shows how to solve a symmetric positive definite linear system Ax=b, using the
Cholesky factorization computed in the previous example in-place in the matrix A. The example
assumes that cholesky_factor(A, t) returned 0, indicating no zero or NaN pivots.

template<class in_matrix_t,
    class Triangle,
    class in_vector_t,
    class out_vector_t>
void cholesky_solve(
in_matrix_t A,
    Triangle t,
    in_vector_t b,
    out_vector_t x)
{
    if constexpr (std::is_same_v<Triangle, upper_triangle_t>) {
        // Solve Ax = b where A = U^T U
        // Solve U^T c = b, using x to store c.
        triangular_matrix_vector_solve(transpose_view(A), t, explicit_diagonal, b, x);
        // Solve U x = c, overwriting x with result.
        triangular_matrix_vector_solve(A, t, explicit_diagonal, x);
    } else {
        // Solve Ax = b where A = L L^T
        // Solve L c = b, using x to store c.
        triangular_matrix_vector_solve(A, t, explicit_diagonal, b, x);
        // Solve L^T x = c, overwriting x with result.
        triangular_matrix_vector_solve(transpose_view(A), t, explicit_diagonal, x);
    }
}
Compute QR factorization of a tall skinny matrix

This example shows how to compute the QR factorization of a "tall and skinny" matrix $V$, using a cache-blocked algorithm based on rank-k symmetric matrix update and Cholesky factorization. "Tall and skinny" means that the matrix has many more rows than columns.

```cpp
// Compute QR factorization $A = Q R$, with $A$ storing $Q$.
template<class inout_matrix_t,
         class out_matrix_t>
int cholesky_tsqr_one_step(
    inout_matrix_t A, // $A$ on input, $Q$ on output
    out_matrix_t R)
{
    // One might use cache size, sizeof(element_type), and $A$.extent(1)
    // to pick the number of rows per block. For now, we just pick
    // some constant.
    constexpr ptrdiff_t max_num_rows_per_block = 500;

    using R_element_type = typename out_matrix_t::element_type;
    constexpr R_element_type ZERO {};
    for(ptrdiff_t i = 0; i < R.extent(0); ++i) {
        for(ptrdiff_t j = 0; j < R.extent(1); ++j) {
            R(0,0) = ZERO;
        }
    }

    // Cache-blocked version of $R = R + A^T \ast A$.
    const ptrdiff_t num_rows = A.extent(0);
    const ptrdiff_t rest_num_rows = num_rows;
    auto A_rest = A;
    while(A_rest.extent(0) > 0) {
        const ptrdiff num_rows_per_block =
            min(A.extent(0), max_num_rows_per_block);
        auto A_cur = subspan(A_rest, pair{0, num_rows_per_block}, all);
        A_rest = subspan(A_rest,
            pair{num_rows_per_block, A_rest.extent(0)}, all);
        // $R = R + A_{cur}^T \ast A_{cur}$
        symmetric_matrix_rank_k_update(transpose_view(A_cur),
            R, upper_triangle);
    }

    const int info = cholesky_factor(R, upper_triangle);
    if(info != 0) {
        return info;
    }
    triangular_matrix_matrix_left_solve(R, upper_triangle, A);
    return info;
}

// Compute QR factorization $A = Q R$. Use $R_{tmp}$ as temporary $R$ factor
// storage for iterative refinement.
template<class in_matrix_t,
         class out_matrix_1_t,
         class out_matrix_2_t,
         class out_matrix_3_t>
int cholesky_tsqr(
```
```c
in_matrix_t A,
out_matrix_1_t Q,
out_matrix_2_t R_tmp,
out_matrix_3_t R)
{
    assert(R.extent(0) == R.extent(1));
    assert(A.extent(1) == R.extent(0));
    assert(R_tmp.extent(0) == R_tmp.extent(1));
    assert(A.extent(0) == Q.extent(0));
    assert(A.extent(1) == Q.extent(1));

    linalg_copy(A, Q);
    const int info1 = cholesky_tsqr_one_step(Q, R);
    if(info1 != 0) {
        return info1;
    }

    // Use one step of iterative refinement to improve accuracy.
    const int info2 = cholesky_tsqr_one_step(Q, R_tmp);
    if(info2 != 0) {
        return info2;
    }

    // R = R_tmp * R
    triangular_matrix_left_product(R_tmp, upper_triangle,
                               explicit_diagonal, R);
    return 0;
}
```