

Deducing the type of variable from its initializer expression

(revision 2)

Programming Language C++
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1 Introduction

This document is a minor revision of the document N1721=04-0161. The document N1721=04-0161 contained the suggested wording for new uses of keyword *auto*, which were unanimously approved by the evolution group meeting in Redmond, October 2004. Based on the discussions and straw-polls in the Lillehammer meeting in April 2005, this document now adds wording for allowing the initialization (with *auto*) of more than one variables in a single statement; N1721=04-0161 allowed only one variable initialization per statement.

2 Proposed features

We suggest that the *auto* keyword would indicate that the type of a variable is to be deduced from its initializer expression. For example:

```
auto x = 3.14; // x has type double
```

The *auto* keyword can occur as a basic type specifier (allow to be used with cv-qualifiers, *, [] and &) and the semantics of *auto* should follow exactly the rules of template argument deduction. Examples (the notation $x : T$ is read as “x has type T”):

```
int foo();  
auto x1 = foo(); // x1 : int  
const auto& x2 = foo(); // x2 : const int&  
auto& x3 = foo(); // x3 : int&: error, cannot bind a reference to a temporary
```

```
float& bar();  
auto y1 = bar(); // y1 : float  
const auto& y2 = bar(); // y2 : const float&  
auto& y3 = bar(); // y3 : float&
```

```
A* fii()
```

```

auto* z1 = fii();      // z1 : A*
auto z2 = fii();      // z2 : A*
auto* z3 = bar();     // error, bar does not return a pointer type

auto z4[] = fii();    // z4 : A*

```

A major concern in discussions of *auto*-like features has been the potential difficulty in figuring out whether the declared variable will be of a reference type or not. Particularly, is unintentional aliasing or slicing of objects likely? For example

```

class B { ... virtual void f(); }
class D : public B { ... void f(); }
B* d = new D();
...
auto b = *d; // is this casting a reference to a base or slicing an object?
b.f();      // is polymorphic behavior preserved?

```

Basing *auto* on template argument deduction rules provides a natural way for a programmer to express his intention. Controlling copying and referencing is essentially the same as with variables whose types are declared explicitly. For example:

```

A foo();
A& bar();
...
A x1 = foo(); // x1 : A
auto x1 = foo(); // x1 : A

A& x2 = foo(); // error, we cannot bind a non-lvalue to a non-const reference
auto& x2 = foo(); // error

A y1 = bar(); // y1 : A
auto y1 = bar(); // y1 : A

A& y2 = bar(); // y2 : A&
auto& y2 = bar(); // y2 : A&

```

Thus, as in the rest of the language, value semantics is the default, and reference semantics is provided through consistent use of **&**.

Multi-variable declarations

More than one variable can be declared in a single statement:

```

int i;
auto a = 1, *b = &i;

```

In the case of two or more variables, both deductions must lead to the same type. Note that the declared variables can get different types, as is the case in the above example. The requirement on the type deductions to lead to the same type is best explained by translation to template argument deduction. The deductions in the above example correspond to the deductions of template parameter *T* below:

```

template <class T>
void foo(T a, T* b);
...
foo(1, &i);

```

Here, *T* must be deduced to be the same type based on both arguments; otherwise the code is ill-defined.

Direct initialization syntax

Direct initialization syntax is allowed and is equivalent to copy initialization. For example:

```
auto x = 1; // x : int
auto x(1); // x : int
```

The semantics of a direct-initialization expression of the form $T v(x)$ with T a type expression containing an occurrence of *auto*, v as a variable name, and x an expression, is defined as a translation to the corresponding copy initialization expression $T v = x$. Examples:

```
const auto& y(x) -> const auto& y = x;
```

It follows that the direct initialization syntax is allowed with *new* expressions as well:

```
new auto(1);
```

The expression *auto*(I) has type *int*, and thus *new auto*(I) has type *int**. Combining a *new* expression using *auto* with an *auto* variable declaration gives:

```
auto* x = new auto(1);
```

Here, *new auto*(I) has type *int**, which will be the type of x too.

3 Proposed wording

Section 7.1.5.1 Type specifiers [dcl.type.simple]

Add to the paragraph 1

- *auto* can either be a storage class specifier, or a simple type specifier. *auto* can be combined with any other type specifier, in which case it is treated as a storage class specifier. If the *decl-specifier-sequence* contains no type specifier other than *auto*, then the following restrictions apply to the *decl-specifier-sequence*:
 - It must be followed by one or more *init-declarators*, each of which must have a non-empty *initializer* of either of the following two forms:

```
= initializer-clause
( initializer-clause )
```

- The only other allowed *decl-specifiers* are *cv-qualifiers* and the storage class specifier *static*.

[*Example*: The following are valid declarations:

```
auto x = 5;
const auto *v = &x, u = 6;
static auto y = 0.0;
static auto int z; // invalid, auto treated as a storage class specifier
auto int r; // ok
```

— *end example*]

Section 7.1.5.2 Simple type specifiers [dcl.type.simple]

In paragraph 1, add the following to the list of simple type specifiers:

```
auto
```

To Table 7, add the line:

auto	placeholder for a type
------	------------------------

Section 8.3 Meaning of declarators [dcl.meaning]

New paragraph after paragraph 1:

If *decl-specifier-sequence* contains the *simple-type-specifier* `auto`, the declarator is required to declare an object and to specify an initial value; the type of the declared identifier is deduced from the type of its initializer ([dcl.auto]).

Replace paragraph 4 with:

First, the *decl-specifier-seq* determines a type; or, when it contains an occurrence of `auto`, a *type scheme*. A type scheme yields a type when the occurrence of `auto` in the type scheme is replaced by a type. In a declaration

```
T D
```

the *decl-specifier-seq* `T` determines the type, or type scheme, “`T`”. [Example: in the declarations

```
int unsigned i;
const auto& p = f();
```

the type specifiers `int unsigned` determine the type “`unsigned int`”, and the type specifier `const auto` determines the type scheme “`const auto`” ([dcl.type.simple].)]

Section 8.3.1 Pointers [dcl.ptr]

Change the first paragraph to:

In a declaration `T D` where `D` has the form

```
* cv-qualifier-seqopt D1
```

and the type, or type scheme, of the identifier in the declaration `T D1` is *derived-declarator-type-list* `T`, then the type, or type scheme, of the identifier of `D` is *derived-declarator-type-list cv-qualifier-seq* pointer to `T`. The cv-qualifiers apply to the pointer and not to the object pointed to.

The change to this paragraph was the addition of the “or type scheme” in two places. Similar changes are needed to Sections 8.3.2–5 discuss how references, arrays, and function types in the declarator propagate to the type of the *declarator-id*. Details not shown.

New subsection: Auto [dcl.auto]

The section should be a subsection of Section 8.3 ([dcl.meaning]). The text of the new subsection:

Once the type scheme of a *declarator-id* has been determined, the type of the declared variable using the *declarator-id* is determined from the type of its initializer using the rules for template argument deduction ([temp.deduct]). Let `T` be the type scheme that has been determined for a variable identifier `d`, and `e` be the initializer expression for `d`. Obtain `U` from `T` by replacing the occurrence of `auto` with a new invented type template parameter `t`. Define a function template as follows:

```
template <class t>
void __f(U __d) {}
```

The type deduced for the variable `d` is then the type that would be deduced for the parameter `__d` in a call to `__f` with `e` as its actual argument. If the template argument deduction would fail, the declaration is ill-formed.

If the list of declarators contains more than one declarator, the type of each declared variable is determined as described above. If the type deduced for the template parameter `t` is not the same in each deduction, the program is ill-formed.

[*Example:*

```
const auto &i = expr;
```

The type scheme is `const auto&`, and the type of `i` is the deduced type of the argument `i` in the call `__f(expr)` of the following function template:

```
template <class t> void __f(const t& i);
```

— *end example*]

Section 8.5 Initializers [dcl.init]

To paragraph 14 add a case:

If the destination type contains the `auto` specifier, see section [dcl.auto].

Section 5.3.4 New [expr.new]

Paragraph 1 specifies the valid forms of new expressions. Add the following form for *new-type-id* to the grammar:

new-type-id:

...

`cv auto direct-new-declaratoropt`

And the text:

If *new-type-id* is of the form “`cv auto direct-new-declaratoropt`”, *new-initializer* with exactly one initializer argument must follow *new-type-id*, or the program is ill-formed. The allocated type is deduced from the type of this initializer argument as follows: Let (`e`) be the *new-initializer*, then the allocated type is the type deduced for the variable `x` in the declaration ([dcl.auto]):

```
cv auto x = e
```

Once the allocated type has been deduced, the semantics of the *new-expression* is as if the form “`cv auto direct-new-declaratoropt`” was written “`T direct-new-declaratoropt`”, where `T` is the type deduced for the allocated type. [*Example:*

```
new auto(1);           // allocated type is int
double& foo();
new const auto[10](foo()); // allocated type is const double
auto x = new auto('a'); // allocated type is char, x is of type char*
```

— *end example*]

4 Acknowledgments

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