Information technology — Programming languages, their environments, and system software interfaces — Floating-point extensions for C —

Part 5: Supplementary attributes

Technologies de l'information — Langages de programmation, leurs environnements et interfaces du logiciel système — Extensions à virgule flottante pour C —

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Foreword

ISO (the International Organization for Standardization) and IEC (the International Electrotechnical Commission) form the specialized system for worldwide standardization. National bodies that are members of ISO or IEC participate in the development of International Standards through technical committees established by the respective organization to deal with particular fields of technical activity. ISO and IEC technical committees collaborate in fields of mutual interest. Other international organizations, governmental and non-governmental, in liaison with ISO and IEC, also take part in the work. In the field of information technology, ISO and IEC have established a joint technical committee, ISO/IEC JTC 1.

The procedures used to develop this document and those intended for its further maintenance are described in the ISO/IEC Directives, Part 1. In particular the different approval criteria needed for the different types of document should be noted. This document was drafted in accordance with the editorial rules of the ISO/IEC Directives, Part 2 (see www.iso.org/directives).

Attention is drawn to the possibility that some of the elements of this document may be the subject of patent rights. ISO and IEC shall not be held responsible for identifying any or all such patent rights. Details of any patent rights identified during the development of the document will be in the Introduction and/or on the ISO list of patent declarations received (see www.iso.org/patents).

Any trade name used in this document is information given for the convenience of users and does not constitute an endorsement.

For an explanation on the meaning of ISO specific terms and expressions related to conformity assessment, as well as information about ISO’s adherence to the WTO principles in the Technical Barriers to Trade (TBT) see the following URL: Foreword - Supplementary information.

The committee responsible for this document is ISO/IEC JTC 1, Information technology, Subcommittee SC 22, Programming languages, their environments, and system software interfaces.

ISO/IEC TS 18661 consists of the following parts, under the general title Information technology — Programming languages, their environments, and system software interfaces — Floating-point extensions for C:

— Part 1: Binary floating-point arithmetic
— Part 2: Decimal floating-point arithmetic
— Part 3: Interchange and extended types
— Part 4: Supplementary functions
— Part 5: Supplementary attributes


Introduction

Background

IEC 60559 floating-point standard

The IEEE 754-1985 standard for binary floating-point arithmetic was motivated by an expanding diversity in floating-point data representation and arithmetic, which made writing robust programs, debugging, and moving programs between systems exceedingly difficult. Now the great majority of systems provide data formats and arithmetic operations according to this standard. The IEC 60559:1989 international standard was equivalent to the IEEE 754-1985 standard. Its stated goals were the following:

1. Facilitate movement of existing programs from diverse computers to those that adhere to this standard.
2. Enhance the capabilities and safety available to programmers who, though not expert in numerical methods, may well be attempting to produce numerically sophisticated programs. However, we recognize that utility and safety are sometimes antagonists.
3. Encourage experts to develop and distribute robust and efficient numerical programs that are portable, by way of minor editing and recompilation, onto any computer that conforms to this standard and possesses adequate capacity. When restricted to a declared subset of the standard, these programs should produce identical results on all conforming systems.
4. Provide direct support for
   a. Execution-time diagnosis of anomalies
   b. Smoother handling of exceptions
   c. Interval arithmetic at a reasonable cost
5. Provide for development of
   a. Standard elementary functions such as exp and cos
   b. Very high precision (multiword) arithmetic
   c. Coupling of numerical and symbolic algebraic computation
6. Enable rather than preclude further refinements and extensions.

To these ends, the standard specified a floating-point model comprising the following:

- **formats** – for binary floating-point data, including representations for Not-a-Number (NaN) and signed infinities and zeros
- **operations** – basic arithmetic operations (addition, multiplication, etc.) on the format data to compose a well-defined, closed arithmetic system; also specified conversions between floating-point formats and decimal character sequences, and a few auxiliary operations
- **context** – status flags for detecting exceptional conditions (invalid operation, division by zero, overflow, underflow, and inexact) and controls for choosing different rounding methods
The ISO/IEC/IEEE 60559:2011 international standard is equivalent to the IEEE 754-2008 standard for floating-point arithmetic, which is a major revision to IEEE 754-1985.

The revised standard specifies more formats, including decimal as well as binary. It adds a 128-bit binary format to its basic formats. It defines extended formats for all of its basic formats. It specifies data interchange formats (which may or may not be arithmetic), including a 16-bit binary format and an unbounded tower of wider formats. To conform to the floating-point standard, an implementation must provide at least one of the basic formats, along with the required operations.

The revised standard specifies more operations. New requirements include – among others – arithmetic operations that round their result to a narrower format than the operands (with just one rounding), more conversions with integer types, more classifications and comparisons, and more operations for managing flags and modes. New recommendations include an extensive set of mathematical functions and seven reduction functions for sums and scaled products.

The revised standard places more emphasis on reproducible results, which is reflected in its standardization of more operations. For the most part, behaviors are completely specified. The standard requires conversions between floating-point formats and decimal character sequences to be correctly rounded for at least three more decimal digits than is required to distinguish all numbers in the widest supported binary format; it fully specifies conversions involving any number of decimal digits. It recommends that transcendental functions be correctly rounded.

The revised standard requires a way to specify a constant rounding direction for a static portion of code, with details left to programming language standards. This feature potentially allows rounding control without incurring the overhead of runtime access to a global (or thread) rounding mode.

Other features recommended by the revised standard include alternate methods for exception handling, controls for expression evaluation (allowing or disallowing various optimizations), support for fully reproducible results, and support for program debugging.

The revised standard, like its predecessor, defines its model of floating-point arithmetic in the abstract. It neither defines the way in which operations are expressed (which might vary depending on the computer language or other interface being used), nor does it define the concrete representation (specific layout in storage, or in a processor’s register, for example) of data or context, except that it does define specific encodings that are to be used for the exchange of floating-point data between different implementations that conform to the specification.

IEC 60559 does not include bindings of its floating-point model for particular programming languages. However, the revised standard does include guidance for programming language standards, in recognition of the fact that features of the floating-point standard, even if well supported in the hardware, are not available to users unless the programming language provides a commensurate level of support. The implementation’s combination of both hardware and software determines conformance to the floating-point standard.

C support for IEC 60559

The C standard specifies floating-point arithmetic using an abstract model. The representation of a floating-point number is specified in an abstract form where the constituent components (sign, exponent, significand) of the representation are defined but not the internals of these components. In particular, the exponent range, significand size, and the base (or radix) are implementation-defined. This allows flexibility for an implementation to take advantage of its underlying hardware architecture. Furthermore, certain behaviors of operations are also implementation-defined, for example in the area of handling of special numbers and in exceptions.
The reason for this approach is historical. At the time when C was first standardized, before the floating-point standard was established, there were various hardware implementations of floating-point arithmetic in common use. Specifying the exact details of a representation would have made most of the existing implementations at the time not conforming.

Beginning with ISO/IEC 9899:1999 (C99), C has included an optional second level of specification for implementations supporting the floating-point standard. C99, in conditionally normative annex F, introduced nearly complete support for the IEC 60559:1989 standard for binary floating-point arithmetic. Also, C99’s informative annex G offered a specification of complex arithmetic that is compatible with IEC 60559:1989.


**Purpose**

The purpose of ISO/IEC TS 18661 is to provide a C language binding for ISO/IEC/IEEE 60559:2011, based on the C11 standard, that delivers the goals of ISO/IEC/IEEE 60559 to users and is feasible to implement. It is organized into five parts.

ISO/IEC TS 18661-1 provides changes to C11 that cover all the requirements, plus some basic recommendations, of ISO/IEC/IEEE 60559:2011 for binary floating-point arithmetic. C implementations intending to support ISO/IEC/IEEE 60559:2011 are expected to conform to conditionally normative annex F as enhanced by the changes in ISO/IEC TS 18661-1.


ISO/IEC TS 18661-3 (Interchange and extended types), ISO/IEC TS 18661-4 (Supplementary functions), and ISO/IEC TS 18661-5 (Supplementary attributes) cover recommended features of ISO/IEC/IEEE 60559:2011. C implementations intending to provide extensions for these features are expected to conform to the corresponding parts.

**Additional background on supplementary attributes**

ISO/IEC/IEEE 60559:2011 defines alternatives for certain attributes of floating-point semantics and intends that programming languages provide means by which a program can specify which of the alternative semantics apply to a given block of code. The program specification of attributes is to be constant (fixed at translation time), not dynamic (changeable at execution time).

ISO/IEC TS 18661 provides these attributes by means of standard pragmas, where the pragma parameters represent the alternative semantics.

The **FENV_ROUND** and **FENV_DEC_ROUND** pragmas, specified in ISO/IEC TS 18661-1 and ISO/IEC TS 18661-2, respectively, provide the rounding direction attributes required by ISO/IEC/IEEE 60559:2011.
ISO/IEC/IEEE 60559:2011 recommends attributes for

- preferredWidth: evaluation formats for floating-point operations
- value-changing optimizations: allow/disallow program transformations that might affect floating-point result values
- alternate exception handling: methods of handling floating-point exceptions
- reproducibility: support for getting floating-point result values and exceptions that are exactly reproducible on other systems

This part of ISO/IEC TS 18661 specifies pragmas that provide these attributes.

Note that the pragmas introduced by ISO/IEC TS 18661 are similar to the floating-point pragmas (FENV_ACCESS, FP_CONTRACT, CX_LIMITED_RANGE) that are already in ISO/IEC 9899. They all have the same general syntactic form, usage requirements, and range of effect.
Information technology — Programming languages, their environments, and system software interfaces — Floating-point extensions for C —

Part 5: Supplementary attributes

1 Scope

This part of ISO/IEC TS 18661 extends programming language C to include support for attributes specified and recommended in ISO/IEC/IEEE 60559:2011.

2 Conformance

An implementation conforms to this part of ISO/IEC TS 18661 if

a) it meets the requirements for a conforming implementation of C11 with all the changes to C11 as specified in parts 1–5 of ISO/IEC TS 18661;

b) it conforms to ISO/IEC TS 18661-1 or ISO/IEC TS 18661-2 (or both); and

c) it defines \texttt{__STDC_IEC_60559_ATTRIBS\_} to 201ymmL.

An implementation conforms to the optional specification for alternate exception handling in this part of ISO/IEC TS 18661 if, in addition to the above,

d) it defines \texttt{__STDC_IEC_60559_ATTRIB_ALTERNATE\_EXCEPTION\_HANDLING\_} to 201ymmL.

3 Normative references

The following documents, in whole or in part, are normatively referenced in this document and are indispensable for its application. For dated references, only the edition cited applies. For undated references, the latest edition of the referenced document (including any amendments) applies.


ISO/IEC TS 18661-2:2015, Information technology — Programming languages, their environments and system software interfaces — Floating-point extensions for C — Part 2: Decimal floating-point arithmetic
4 Terms and definitions


4.1 C11


5 C standard conformance

5.1 Freestanding implementations

The specification in C11 + TS18661-1 + TS18661-2 allows freestanding implementations to conform to this part of ISO/IEC TS 18661.

5.2 Predefined macros

Change to C11 + TS18661-1 + TS18661-2 + TS18661-3 + TS18661-4:

In 6.10.8.3#1, add:

```
#pragma STDC_IEC_60559_ATTRIBS__  The integer constant 201ymmL, intended to indicate support of attributes specified and recommended in IEC 60559.
```

```
#pragma STDC_IEC_60559_ATTRIB_ALTERNATE_EXCEPTION_HANDLING__  The integer constant 201ymmL, intended to indicate support of the alternate exception handling attributes specified and recommended in IEC 60559.
```

6 Standard pragmas

C11 provides standard pragmas (6.10.6) for specifying certain attributes pertaining to floating-point behavior within a compound statement or file. This part of ISO/IEC TS 16881 extends this practice by introducing additional standard pragmas to support the attributes recommended by IEC 60559.

Change to C11 + TS18661-1 + TS18661-2 + TS18661-3 + TS18661-4:

In 6.10.6#2, append to the list of standard pragmas:

```
#pragma STDC FENV_FLT_EVAL_METHOD width
#pragma STDC FENV_DEC_EVAL_METHOD width
#pragma STDC FENV_ALLOW_VALUE_CHANGING_OPTIMIZATION on-off-switch
#pragma STDC FENV_ALLOW_ASSOCIATIVE_LAW on-off-switch
#pragma STDC FENV_ALLOW_DISTRIBUTIVE_LAW on-off-switch
#pragma STDC FENV_ALLOW_MULTIPLY_BY_RECIPROCAL on-off-switch
```
7 Evaluation formats

IEC 60559 recommends attributes for specifying a preferred width for operation results. These preferred widths correspond to the evaluation formats defined in C11, though C11 does not provide means for the user to control the evaluation format. This part of ISO/IEC TS 16881 provides pragmas in `<fenv.h>` to control the evaluation format, using constants with the values of the `FLT_EVAL_METHOD` and `DEC_EVAL_METHOD` macros (5.2.4.2.2a) to represent the evaluation formats.

The evaluation methods in C99 apply to floating-point operators, but not to math functions. Hence, they do not apply to the IEC 60559 operations that are provided as library functions. This clause specifies a macro the user can define to cause the generic macros in `<tgmath.h>` to be evaluated like floating-point operators.

Changes to C11 + TS18661-1 + TS18661-2 + TS18661-3 + TS18661-4:

- After 5.2.4.2.2a#3, insert:

  > [3a] The `FLT_EVAL_METHOD` and `DEC_EVAL_METHOD` macros characterize the use of evaluation formats at the point in the program where the macro is used. Thus, the values of these macros reflect the state of any evaluation method pragmas (7.6.1c, 7.6.1d) that are in effect. These macros shall not be used in a `#if` or `#elif` expression within the scope of a corresponding evaluation method pragma.

- After 7.6.1b, insert:

  > 7.6.1c Evaluation method pragma

  Synopsis

  ```
  #define __STDC_WANT_IEC_60559_ATTRIBS_EXT__

  #include <fenv.h>
  
  #pragma STDC FENV_FLT_EVAL_METHOD width
  ```

  Description

  The `FENV_FLT_EVAL_METHOD` pragma sets the evaluation method for standard floating types and for binary interchange and extended floating types to the evaluation method represented by `width`. The parameter `width` is an expression in one of the forms

  - `decimal-constant`
  - `− decimal-constant`
or

```
        DEFAULT
```

where the value of an expression is a possible value of the \texttt{FLT_EVAL_METHOD} macro, as specified in 5.2.4.2.2a. An expression represents the evaluation method corresponding to its value (5.2.4.2.2a) and \texttt{DEFAULT} designates the implementation's default evaluation method (characterized by the \texttt{FLT_EVAL_METHOD} macro where no \texttt{FENV_FLT_EVAL_METHOD} pragma is in effect). \texttt{width} may be \texttt{-1}, \texttt{0}, or \texttt{DEFAULT}. Which, if any, other values of \texttt{width} are supported is implementation-defined. Use of unsupported values of \texttt{width} results in undefined behavior. The pragma shall occur either outside external declarations or preceding all explicit declarations and statements inside a compound statement. When outside external declarations, the pragma takes effect from its occurrence until another \texttt{FENV_FLT_EVAL_METHOD} pragma is encountered, or until the end of the translation unit. When inside a compound statement, the pragma takes effect from its occurrences until another \texttt{FENV_FLT_EVAL_METHOD} pragma is encountered (including within a nested compound statement), or until the end of the compound statement; at the end of a compound statement the state for the pragma is restored to its condition just before the compound statement.

7.6.1d Evaluation method pragma for decimal

Synopsis

```
[1] #define __STDC_WANT_IEC_60559_DFP_EXT__
#define __STDC_WANT_IEC_60559_ATTRIBS_EXT__
#include <fenv.h>
#pragma STDC FENV_DEC_EVAL_METHOD width
```

Description

[2] The \texttt{FENV_DEC_EVAL_METHOD} pragma sets the evaluation method for decimal interchange and extended floating types to the evaluation method represented by \texttt{width}. The parameter \texttt{width} is an expression in one of the forms

```
0
decimal-constant
-- decimal-constant
```

or

```
        DEFAULT
```

where the value of an expression is a possible value of the \texttt{DEC_EVAL_METHOD} macro, as specified in 5.2.4.2.2a. An expression represents the evaluation method corresponding to its value (5.2.4.2.2a) and \texttt{DEFAULT} designates the implementation's default evaluation method (characterized by the \texttt{DEC_EVAL_METHOD} macro where no \texttt{FENV_DEC_EVAL_METHOD} pragma is in effect). \texttt{width} may be \texttt{-1}, \texttt{1}, or \texttt{DEFAULT}. Which, if any, other values of \texttt{width} are supported is implementation-defined. Use of unsupported values of \texttt{width} results in undefined behavior. The pragma shall occur either outside external declarations or preceding all explicit declarations and statements inside a compound statement. When outside external declarations, the pragma takes effect from its occurrence until another \texttt{FENV_DEC_EVAL_METHOD} pragma is encountered, or until the end of the translation unit. When inside a compound statement, the pragma takes effect from its occurrences until another \texttt{FENV_DEC_EVAL_METHOD} pragma is
encountered (including within a nested compound statement), or until the end of the compound statement, at the end of a compound statement the state for the pragma is restored to its condition just before the compound statement.

At the end of 7.12#2, append:

[2] ... These types reflect the evaluation method where no evaluation method pragma (7.6.1c, 7.6.1d) is in effect.

[2a] For each of the types above, a type-like macro with the same name expands to a designation for the type whose range and precision (5.2.4.2.2a) are used for evaluating operations and constants of the corresponding standard, binary, or decimal floating type. The macro reflects the actual evaluation method, which might be determined by an evaluation method pragma (7.6.1c, 7.6.1d). Use of \#undef to remove the macro definition will ensure that the actual type will be referred to.

After 7.25#2, insert:

[2a] Except for functions that round result to a narrower type, if the macro

\[ \texttt{\_\_STDC\_TGMATH\_OPERATOR\_EVALUATION\_}\]

is defined at the point in the program where <tgmath.h> is first included, the format of the generic parameters of the function invoked by a type-generic macro is determined by the effective evaluation method (see 5.2.4.2.2 and 5.2.4.2.2a), based on the types of the arguments for generic parameters. The semantic type of the expanded type-generic macro is unchanged by the evaluation method. Neither the arguments for generic parameters nor the result are narrowed to their semantic types. Thus, if the macro \[ \texttt{\_\_STDC\_TGMATH\_OPERATOR\_EVALUATION\_}\] is appropriately defined) the evaluation method affects the operations provided by type-generic macros and floating-point operators in the same way. See EXAMPLE 2 below.

After the first bullet in 7.25#3c, insert the bullet:

If the macro

\[ \texttt{\_\_STDC\_TGMATH\_OPERATOR\_EVALUATION\_}\]

is defined at the point in the program where <tgmath.h> is first included, the format of the generic parameters of the function invoked is given by the effective evaluation method based on the type determined below. The semantic type of the expanded type-generic macro is unchanged by the evaluation method.

In 7.25#7, change “EXAMPLE” to “EXAMPLE 1”.

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After 7.25#7, append:

[7a] EXAMPLE 2 The following code uses wide evaluation to avoid overflow and underflow.

```c
#define __STDC_WANT_IEC_60559_ATTRIBS_EXT__
#define __STDC_TGMATH_OPERATOR_EVALUATION__
#include <fenv.h>
#include <tgmath.h>
{
    #pragma STDC FLT_EVAL_METHOD 1
    float x, y, z;
    ...
    z = sqrt(x * x + y * y);
}
```

Because of the use of the evaluation method pragma, the sum of squares, whose semantic type is `float`, is evaluated with the range and precision of `double`, hence does not overflow or underflow. The expanded `<tgmath.h>` macro `sqrt` acquires the semantic type of its argument: `float`. However, because the macro `__STDC_TGMATH_OPERATOR_EVALUATION__` is defined before the inclusion of `<tgmath.h>`, the `sqrt` macro behaves like an operator with respect to the evaluation method and does not narrow its argument to its semantic type. Without the definition of the macro `__STDC_TGMATH_OPERATOR_EVALUATION__`, the `sqrt` macro would expand to `sqrtf` and its evaluated argument would be converted to `float`, which might overflow or underflow.

### 8 Optimization controls

IEC 60559 recommends attributes to allow and disallow value-changing optimizations, individually and collectively. C11 Annex F disallows value-changing optimizations, except for contractions (which can be controlled as a group with the `FP_CONTRACT` pragma). This part of ISO/IEC TS 18661 provides pragmas to allow or disallow certain value-changing optimizations, including those mentioned in IEC 60559. These pragmas apply to all floating types, not just the real floating types (which provide the IEC 60559 formats).

Change to C11 + TS18661-1 + TS18661-2 + TS18661-3 + TS18661-4:

After 7.6.1d, insert:

#### 7.6.1e Optimization control pragmas

[1] The pragmas in this subclause can be used to allow the implementation to do certain floating-point optimizations that are generally disallowed because the optimization might change values of floating-point expressions. These pragmas apply to all floating types. It is unspecified whether optimizations allowed by these pragmas occur consistently, or at all. These pragmas (among other standard pragmas) apply to user code. They do not apply to code for operators or library functions that might be placed inline by the implementation.

[2] Some of the pragmas allow optimizations based on identities of real number arithmetic that are not valid for (IEC 60559) floating-point arithmetic (5.1.2.3, F.9.2). Optimizations based on identities that are valid for floating-point arithmetic are always allowed. Optimizations based on identities derived from identities whose use is allowed (either by a standard pragma or by virtue of being valid for floating-point arithmetic) may also be done.

[3] These pragmas do not affect the requirements on volatile or atomic variables.
[4] Each pragma shall occur either outside external declarations or preceding all explicit declarations and statements inside a compound statement. When outside external declarations, the pragma takes effect, on each optimization it controls, from its occurrence until another pragma that affects the same optimization is encountered, or until the end of the translation unit. When inside a compound statement, the pragma takes effect, on each optimization it controls, from its occurrence until another pragma that affects the same optimization is encountered (including within a nested compound statement), or until the end of the compound statement: at the end of a compound statement the state for allowing each optimization controlled by the pragma is restored to its condition just before the compound statement.

7.6.1 e.1 The FENV_ALLOW_VALUE_CHANGING_OPTIMIZATION pragma

Synopsis

[1] #define __STDC_WANT_IEC_60559_ATTRIBS_EXT__
#include <fenv.h>
#pragma STDC FENV_ALLOW_VALUE_CHANGING_OPTIMIZATION on-off-switch

Description

[1] This pragma is equivalent to all the optimization pragmas specified below, with the same value of on-off-switch (ON, OFF, or DEFAULT).

[2] NOTE The FENV_ALLOW_VALUE_CHANGING_OPTIMIZATION pragma does not affect the evaluation methods. Nevertheless, an evaluation method characterized by a negative value of width (5.2.4.2.2a) might allow for indeterminable evaluation formats, hence unspecified result values.

7.6.1 e.2 The FENV_ALLOW_ASSOCIATIVE_LAW pragma

Synopsis

[1] #define __STDC_WANT_IEC_60559_ATTRIBS_EXT__
#include <fenv.h>
#pragma STDC FENV_ALLOW_ASSOCIATIVE_LAW on-off-switch

Description

[2] This pragma allows or disallows optimizations based on the associative laws for addition and multiplication

\[ x + (y + z) = (x + y) + z \]
\[ x \times (y \times z) = (x \times y) \times z \]

where on-off-switch is one of

ON – allow application of the associative laws
OFF – do not allow application of the associative laws
DEFAULT – “off”
[3] Note that this pragma allows optimizations based on similar mathematical identities involving subtraction and division. For example, since the identity \( x - y = x + (-y) \) is valid for IEC 60559 floating-point arithmetic (F.9.2), this pragma also allows optimizations based on

\[
x + (y - z) = (x + y) - z
\]

Similarly, if the states for this pragma and the FENV_ALLOW_MULTIPLY_BY_RECIPROCAL pragma (7.6.1e.4) are both "on", then optimizations based on the following are allowed:

\[
x \times (y/z) = (x \times y) / z
\]

Note also that, since the commutative laws

\[
x + y = y + x \\
x \times y = y \times x
\]

are valid for IEC 60559 floating-point arithmetic, the pragma allows optimizations based on identities derived from the associative and commutative laws, such as

\[
x + (z + y) = (x + y) + z
\]

7.6.1e.3 The FENV_ALLOW_DISTRIBUTIVE_LAW pragma

Synopsis

[1] #define __STDC_WANT_IEC_60559_ATTRIBS_EXT__
#include <fenv.h>
#pragma STDC FENV_ALLOW_DISTRIBUTIVE_LAW on-off-switch

Description

[2] This pragma allows or disallows optimizations based on the distributive laws for multiplication and division

\[
x \times (y + z) = (x \times y) + (x \times z) \\
x \times (y - z) = (x \times y) - (x \times z) \\
(x + y) / z = (x / z) + (y / z) \\
(x - y) / z = (x / z) - (y / z)
\]

where on-off-switch is one of

ON – allow application of the distributive laws
OFF – do not allow application of the distributive laws
DEFAULT – “off”

7.6.1e.4 The FENV_ALLOW_MULTIPLY_BY_RECIPROCAL pragma

Synopsis

[1] #define __STDC_WANT_IEC_60559_ATTRIBS_EXT__
#include <fenv.h>
#pragma STDC FENV_ALLOW_MULTIPLY_BY_RECIPROCAL on-off-switch
Description

[2] This pragma allows or disallows optimizations based on the mathematical equivalence of division and multiplication by the reciprocal of the denominator

\[ \frac{x}{y} = x \times \left( \frac{1}{y} \right) \]

where on-off-switch is one of

- **ON** – allow multiply by reciprocal
- **OFF** – do not allow multiply by reciprocal
- **DEFAULT** – “off”

### 7.6.1e.5 The FENV_ALLOW_ZERO_SUBNORMAL pragma

#### Synopsis

[1] #define __STDC_WANT_IEC_60559_ATTRIBS_EXT__
#include <fenv.h>
#pragma STDC FENV_ALLOW_ZERO_SUBNORMAL on-off-switch

Description

[2] This pragma allows or disallows replacement of subnormal operands and results by zero, where on-off-switch is one of

- **ON** – allow replacement of subnormals with zero
- **OFF** – do not allow replacement of subnormals with zero
- **DEFAULT** – “off”

[3] Within the scope of this pragma, the floating-point operations affected by the pragma are the same operations as would be affected by the FENV_ROUND and FENV_DEC_ROUND pragmas (7.6.1a, 7.6.1b). Thus, subnormal operands and results of affected operations may be replaced by zero. Functions not affected by the pragma behave as though no FENV_ALLOW_ZERO_SUBNORMAL pragma were in effect at the site of the call.

#### 7.6.1e.6 The FENV_ALLOW_CONTRACT_FMA pragma

#### Synopsis

[1] #define __STDC_WANT_IEC_60559_ATTRIBS_EXT__
#include <fenv.h>
#pragma STDC FENV_ALLOW_CONTRACT_FMA on-off-switch
Description

[2] This pragma allows or disallows contraction (6.5) of floating-point multiply and add or subtract (with the result of the multiply)

\[
x \times y + z \\
x \times y - z \\
x + y \times z \\
x - y \times z
\]

where on-off-switch is one of

- **ON** – allow contraction for floating-point multiply-add
- **OFF** – do not allow contraction for floating-point multiply-add
- **DEFAULT** – implementation defined whether “on” or “off”

[3] **NOTE** IEC 60559 uses the term *synthesize* instead of *contract*.

7.6.1e.7 The **FENV_ALLOW_CONTRACT_OPERATION_CONVERSION** pragma

Synopsis

[1] #define __STDC_WANT_IEC_60559_ATTRIBS_EXT__
    #include <fenv.h>
    #pragma STDC FENV_ALLOW_CONTRACT_OPERATION_CONVERSION on-off-switch

Description

[2] This pragma allows or disallows contraction (6.5) of a floating-point operation and a conversion (of the result of the operation), where on-off-switch is one of

- **ON** – allow contraction for floating-point operation-conversion
- **OFF** – do not allow contraction for floating-point operation-conversion
- **DEFAULT** – implementation defined whether “on” or “off”

[3] Within the scope of this pragma, the floating-point operations affected by the pragma are the same operations as would be affected by the **FENV_ROUND** and **FENV_DEC_ROUND** pragmas (7.6.1a, 7.6.1b). Thus, an affected operation may be contracted with a conversion of its result. Functions not affected by the pragma behave as though no **FENV_ALLOW_CONTRACT_OPERATION_CONVERSION** pragma were in effect at the site of the call.
EXAMPLE

For the code sequence

```c
#define __STDC_WANT_IEC_60559_ATTRIBS_EXT__
#include <fenv.h>
#pragma STDC FENV_ALLOW_CONTRACT_OPERATION_CONVERSION ON
float f1, f2;
double d1, d2;
...
f1 = d1 * d2;
f2 = sqrt(d1);
```

the multiply (operation) and assignment (conversion) are allowed to be evaluated with just one rounding (to the range and precision of `float`). If the on-off-switch for the pragma were `OFF`, then the multiply would have to be rounded according to the evaluation method and the assignment would require a second rounding. With the given code, the `sqrt` function may be replaced by `fsqrt`, avoiding the need for a separate operation to convert the `double` result of `sqrt` to `float`.

7.6.1e.8 The FENV_ALLOW_CONTRACT pragma

Synopsis

```c
#define __STDC_WANT_IEC_60559_ATTRIBS_EXT__
#include <fenv.h>
#pragma STDC FENV_ALLOW_CONTRACT on-off-switch
```

Description

This pragma allows or disallows contraction (6.5) for floating-point operations, where `on-off-switch` is one of

- **ON** – allow contraction for floating-point operations
- **OFF** – do not allow contraction for floating-point operations
- **DEFAULT** – implementation defined whether “on” or “off”

The optimizations controlled by this pragma include those controlled by the `FENV_ALLOW_CONTRACT_FMA` and `FENV_ALLOW_CONTRACT_OPERATION_CONVERSION` pragmas.

This pragma is equivalent to the `FP_CONTRACT` pragma in `<math.h>`: the two pragmas may be used interchangeably, provided the appropriate header is included and the implementation defines `__STDC_WANT_IEC_60559_ATTRIBS_EXT__`.

Reproducibility

IEC 60559 recommends an attribute to facilitate writing programs whose floating-point results and exception flags will be reproducible on any implementation that supports the language and library features used by the program. Such code must use only those features of the language and library that support reproducible results. These features include ones with a well-defined binding to reproducible features of IEC 60559, so that no unspecified or implementation-defined behavior is admitted. This part of ISO/IEC TS 18661 provides a pragma to support the IEC 60559 attribute for reproducible results and gives requirements for programs to have reproducible results.
Change to C11 + TS18661-1 + TS18661-2 + TS18661-3 + TS18661-4:

After 7.6.1e, insert:

7.6.1f  Reproducible results

[1] The pragma in this subclause supports the reproducible results attribute recommended in IEC 60559. Where the state of the pragma is "on", floating-point numerical results and exception flags are reproducible on implementations that define __STDC_IEC_60559_ATTRIBS__ and that support the language and library features used by the source code, provided the source code uses a limited set of features as described below (7.6.1f.2).

[2] An implementation that defines __STDC_IEC_60559_ATTRIBS__ also defines either __STDC_IEC_60559_BFP__ or __STDC_IEC_60559_DFP__, or both. If the implementation defines __STDC_IEC_60559_BFP__, it supports reproducible results for binary floating-point arithmetic. If the implementation defines __STDC_IEC_60559_DFP__, it supports reproducible results for decimal floating-point arithmetic. If the implementation defines __STDC_IEC_60559_TYPES__, then it supports reproducible results for code using its interchange floating types. If the implementation defines __STDC_IEC_60559_FUNCS__ and it provides a set of correctly rounded math functions (7.31.6a), then it supports reproducible results for code using correctly rounded math functions from that set.

7.6.1f.1 The FENV_REPRODUCIBLE pragma

Synopsis

[1] #define __STDC_WANT_IEC_60559_ATTRIBS_EXT__
#include <fenv.h>
#pragma STDC FENV_REPRODUCIBLE on-off-switch

Description

[2] This pragma enables or disables support for reproducible results. The pragma shall occur either outside external declarations or preceding all explicit declarations and statements inside a compound statement. When outside external declarations, the pragma takes effect from its occurrence until another FENV_REPRODUCIBLE pragma is encountered, or until the end of the translation unit. When inside a compound statement, the pragma takes effect from its occurrence until another FENV_REPRODUCIBLE pragma is encountered (including within a nested compound statement), or until the end of the compound statement; at the end of a compound statement the state for the pragma is restored to its condition just before the compound statement.

[3] If the state of the pragma is "on", then the effects of the following are implied

    #pragma STDC FENV_ACCESS ON
    #pragma STDC FENV_ALLOW_VALUE_CHANGING_OPTIMIZATION OFF

and if __STDC_IEC_60559_BFP__ is defined

    #pragma STDC FENV_FLT_EVAL_METHOD 0
and if __STDC_IEC_60559_DFP__ is defined

    #pragma STDC FENV_DEC_EVAL_METHOD 1

[4] If the pragma appears with the on-off-switch OFF under the effect of a pragma with on-off-switch ON, then the states of the FENV_ACCESS pragma, the value-changing optimization pragmas, and the evaluation method pragmas (even an evaluation method pragma whose state was explicitly changed under the effect of the pragma with on-off-switch ON) revert to their states prior to the pragma with on-off-switch ON. The pragma with on-off-switch OFF has no effect if it occurs where the state of the pragma is “off”.

[5] The “default” state of the pragma is “off”.

[6] The implementation should produce a diagnostic message if, where the state of the FENV_REPRODUCIBLE pragma is “on”, the source code uses a language or library feature whose results may not be reproducible.

7.6.1f.2 Reproducible code

[1] Following are requirements for a code sequence in order that its results will be reproducible. If the code uses optional features noted below, then results are reproducible only on implementations that support those features.

— The code translates into a sequence of floating-point operations that are bound to IEC 60559 operations, as described in F.3 in the table entitled “Operation binding”.

— The code does not contain any use that may result in undefined behavior. The code does not depend on any behavior that is unspecified, implementation-defined, or locale-specific. See Annex J.

— The code is under the effect of the FENV_REPRODUCIBLE pragma (with state “on”).

— The code does not set the state of any pragma that allows value-changing optimizations to “on” or “default”.

— The code does not set the state of the FENV_ACCESS pragma to “off” or “default”.

— The code does not use the FENV_FLT_EVAL_METHOD pragma with any width except 0 or 1. Support for width equal to 1 is an optional feature.

— The code does not use the FENV_DEC_EVAL_METHOD pragma with any width except 1 or 2. Support for width equal to 2 is an optional feature.

— Use of an FENV_EXCEPT pragma with an except-list that includes sub-exceptions is an optional feature.

— The code does not use an FENV_EXCEPT pragma with an action BREAK, TRY, or CATCH.

— The code does not use the long double type.

— If __STDC_IEC_60559_BFP__ is not defined by the implementation, the code does not use any standard floating types.
Even if __STDC_IEC_60559_TYPES__ is defined, the code does not use extended floating types. Even if __STDC_IEC_60559_TYPES__ is defined, some interchange floating types are optional features.

The code does not use complex or imaginary types.

The code does not use signaling NaNs.

The code does not use the remquo functions.

The code does not depend on the sign of a zero result or the quantum of a decimal result for the fmin, fmax, fminmag, and fmaxmag functions when the arguments are equal.

The code does not depend on the payloads (F.10.13) or sign bits of quiet NaNs.

The code does not depend on the “underflow” or “inexact” floating-point exceptions or flags.

The code does not use extended integer types.

The code does not depend on conversions between binary floating types and character sequences with more than \( M + 3 \) significant decimal digits, where \( M \) is 17 if __STDC_IEC_60559_TYPES__ is not defined (by the implementation), and \( M \) is \( 1 + \lceil p \times \log_{10}(2) \rceil \), where \( p \) is the precision of the widest supported binary interchange floating type, if __STDC_IEC_60559_TYPES__ is defined. Even if __STDC_IEC_60559_TYPES__ is defined, support for interchange floating types wider than binary64 is an optional feature. (This specification differs from IEC 60559 which requires that an implementation supporting reproducibility not limit the number of significant decimal digits for correct rounding.)

The code does not depend on the actual character sequence in printf results with style a (or A), nor does it depend on numerical values of such results when the precision is not sufficient for an exact representation.

## 10 Alternate exception handling

IEC 60559 arithmetic raises floating-point exceptions as a way to inform the program when an operation encounters problematic inputs, such that no one result would be suitable for all situations. The default exception handling in IEC 60559 is intended to be more useful in more situations than other schemes, or at least predictable. However, other exception handling is more useful in certain situations. Thus, IEC 60559 describes alternate exception handling and recommends that programming languages provide a means for the program to specify which exception handling will be done.

**Changes to C11 + TS18661-1 + TS18661-2 + TS18661-3 + TS18661-4:**

After 7.6.1f.2, insert:

### 7.6.1g Alternate exception handling

[1] When a floating-point exception is raised, the IEC 60559 default exception handling sets the appropriate exception flag(s), returns a specified result, and continues execution. IEC 60559 also prescribes alternate exception handling. The pragma in this subclause provides a means for the program to choose the method of exception handling. The pragma applies to operations on all floating-point types.
[2] For the “underflow” exception, the chosen exception handling occurs if the exception is raised, whether the default result would be exact or inexact, unless stated otherwise.

[3] Alternate exception handling is an optional feature for implementations that support IEC 60559 attributes. Implementations that provide the feature define the macro

```c
__STDC_IEC_60559_ATTRIB_ALTERNATE_EXCEPTION_HANDLING__
```

as well as

```c
__STDC_IEC_60559_ATTRIBS__
```

See 6.10.8.3.

7.6.1g.1 The FENV_EXCEPT pragma

**Synopsis**

[1] #define __STDC_WANT_IEC_60559_ATTRIBS_EXT__

```c
#include <fenv.h>
#pragma STDC FENV_EXCEPT action except-list
```

**Description**

[2] The **FENV_EXCEPT** pragma sets the method specified by `action` for handling the exceptions represented by `except-list`.

[3] The pragma shall occur either outside external declarations, unless stated otherwise, or preceding all explicit declarations and statements inside a compound statement, which then is the compound statement associated with the pragma. When outside external declarations, the pragma `action` for a designated exception takes effect from the occurrence of the pragma until another **FENV_EXCEPT** pragma designating the same exception is encountered, or until the end of the translation unit. When inside a compound statement, the pragma action for a designated exception takes effect from the occurrence of the pragma until another **FENV_EXCEPT** pragma designating the same exception is encountered (including within a nested compound statement), or until the end of the compound statement; at the end of a compound statement the state for handling each exceptions in `except-list` is restored to its condition just before the compound statement.

[4] Within the scope of an **FENV_EXCEPT** pragma, the floating-point operations affected by the pragma are the same operations as would be affected by the **FENV_ROUND** and **FENV_DEC_ROUND** pragmas (7.6.1a, 7.6.1b). Thus, exceptions raised by affected operations are handled according to the specified `action`. Functions not affected by the pragma behave as though no **FENV_EXCEPT** pragma were in effect at the site of the call.

[5] Within the scope of an **FENV_EXCEPT** pragma, explicitly accessing the flag of a designated exception results in undefined behavior.
[6] except-list shall be a comma-separated list of supported exception designations (or one supported exception designation). The supported exception designations shall include the exception macro identifiers (7.6)

```
FE_DIVBYZERO
FE_INEXACT
FE_INVALID
FE_OVERFLOW
FE_UNDERFLOW
FE_ALL_EXCEPT
```

and should include designations for the sub-exceptions

- "invalid" floating-point exceptions from add and subtract operators and functions that add or subtract (7.12.13a.1, 7.12.13a.2), not caused by signaling NaN input
  ```
  FE_INVALID_ADD
  ```

- "invalid" floating-point exceptions from divide operators and functions that divide (7.12.13a.4), not caused by signaling NaN input
  ```
  FE_INVALID_DIV
  ```

- "invalid" floating-point exceptions from functions that compute multiply-add (7.12.13.1, 7.12.13a.5) and from contracted multiply and add operators, not caused by signaling NaN input
  ```
  FE_INVALID_FMA
  ```

- "invalid" floating-point exceptions from conversions from floating to integer types, not caused by signaling NaN input
  ```
  FE_INVALID_INT
  ```

- "invalid" floating-point exceptions from `ilogb` and `llogb` functions, not caused by signaling NaN input
  ```
  FE_INVALID_ILOGB
  ```

- "invalid" floating-point exceptions from multiply operators and functions that multiply (7.12.13a.3), not caused by signaling NaN input
  ```
  FE_INVALID_MUL
  ```

- "invalid" floating-point exceptions from the `quantize` functions, not caused by signaling NaN input
  ```
  FE_INVALID_QUANTIZE
  ```

- "invalid" floating-point exceptions from the `remainder` and `remquo` functions, not caused by signaling NaN input
  ```
  FE_INVALID_REM
  ```

- "invalid" floating-point exceptions from functions that compute square root (7.12.7.5, 7.12.13a.6), not caused by signaling NaN input
  ```
  FE_INVALID_SQRT
  ```

- "invalid" floating-point exceptions caused by signaling NaN input
  ```
  FE_INVALID_SNAN
  ```
— “invalid” floating-point exceptions from relational operators and comparison macros, not caused by signaling NaN input
    FE_INVALID_UNORDERED

— “divide-by-zero” floating-point exceptions from divide operators and functions that divide
    FE_DIVBYZERO_ZERO

— “divide-by-zero” floating-point exceptions from logarithm and logb functions
    FE_DIVBYZERO_LOG

[7] action shall be a designation of a supported exception handling method, which shall include

— default exception handling (as specified in IEC 60559).
    DEFAULT

— default exception handling, but without setting the flag.
    NO_FLAG

— default exception handling, but whether the flag is set (as with default exception handling), and for which operations and their occurrences, is unspecified.
    OPTIONAL_FLAG

— abrupt underflow. If an “underflow” floating-point exception occurs (see IEC 60559), the operation delivers a result with magnitude zero or the minimum normal magnitude (for the result format) and with the same sign as the default result, sets the “underflow” floating-point exception flag, and raises the “inexact” floating-point exception. When

        rounding to nearest, ties to even
        rounding to nearest, ties away from zero

or

        rounding toward zero

the result magnitude is zero. When rounding toward positive infinity, the result magnitude is the minimum normal magnitude if the result sign is positive, and zero if the result sign is negative. When rounding toward negative infinity, the result magnitude is the minimum normal magnitude if the result sign is negative, and zero if the result sign is positive. Abrupt underflow has no effect on the interpretation of subnormal operands. The action has no effect if FE_UNDERFLOW is not included in except-list.
    ABRUPT_UNDERFLOW

— break. Terminate execution of the compound statement associated with the pragma. Then, continue execution after the associated compound statement. When termination occurs, the following apply: if the execution to completion of the associated compound statement (without the break) would at any point modify an object, the value of the object is indeterminate; if the execution would modify the state of the dynamic rounding mode or any state maintained by the standard library (e.g., in the I/O system), the state is unspecified; the values of flags for the designated exceptions are unspecified. (Thus, termination may occur as soon as possible after the exception is raised, to maximize performance.) The pragma with this action shall appear only in a compound statement.
    BREAK
The following two actions work together. The pragmas with these actions shall appear only in compound statements. A compound statement associated with a try action may be paired with one or more compound statements each associated with a catch action. Compound statements associated with catch actions shall appear contiguously immediately below the compound statement associated with the try action, except for white space (including comments).

— try. The designated exceptions may be handled by a catch action. It is unspecified whether flags for designated exceptions that are set in the execution of the associated compound statement are restored to their states before the associated compound statement. The associated compound statement shall not be the statement of a selection (6.8.4) or iteration (6.8.5) statement. There shall be no jumps into or out of the associated compound statement, other than to handle an exception, as specified below.

TRY

— catch. If any exception with the same designation for both the try action and a catch action occurs in the execution of the compound statement associated with the try action, jump to a compound statement associated with some occurring exception with the same designation for the try action and a catch action. Upon completion of the associated compound statement, continue execution after the last of the compound statements associated with catch actions. Each exception designation shall be listed in at most one of the pragmas with a catch action. The jump target should be a compound statement associated with the first occurring designated exception. When the jump occurs, the following apply: if the execution of the associated compound statement to completion (without the jump) would at any point modify an object, the value of the object is indeterminate; if the execution would modify the state of the dynamic rounding mode or any state maintained by the standard library (e.g., in the I/O system), the state is unspecified. (Thus, the jump may occur as soon as possible after the exception is raised, to maximize performance). The compound statement associated with a catch action is executed only to handle an exception occurring in the compound statement associated with the try action. There shall be no other jumps into or out of the compound statement associated with a catch action.

CATCH

The following two actions work together. The pragmas with these actions shall appear only in compound statements. A compound statement associated with a delayed-try action may be paired with one or more compound statements each associated with a delayed-catch action. Compound statements associated with delayed-catch actions shall appear contiguously immediately below the compound statement associated with the delayed-try action, except for white space (and comments). For (optional) support of sub-expressions, the exceptions, flags, and functions in the specification below are assumed to be extended for sub-expressions.

— delayed-try. The designated exceptions may be handled by a delayed-catch action. Before executing the compound statement associated with the delayed-try action, save (as by fegetexceptflag) the states of the flags for the designated exceptions, and then clear (as by feclearexcept) the designated exceptions. After normal completion of the associated compound statement, re-save the states of the designated exceptions. Then restore (as by fegetexceptflag) the designated exception flag states before the associated compound statement. The associated compound statement shall not be the statement of a selection (6.8.4) or iteration (6.8.5) statement. There shall be no jumps into or out of the associated compound statement.

DELAYED_TRY

— delayed-catch. Test (as by fetestexceptflag) the exception flag states saved after completion of the compound statement associated with the delayed-try action. If any
exception with the same designation for the delayed-try action and a delayed-catch action occurred (as determined by flag state tests), jump to the first compound statement associated with an occurring exception with the same designation for the delayed-try action and a delayed-catch action. Upon completion of the associated compound statement, continue execution after the last of the compound statements associated with delayed-catch actions. Each exception designation shall be listed in at most one of the pragmas with a delayed-catch action. The compound statement associated with a delayed-catch action is executed only to handle an exception occurring in the compound statement associated with the delayed-try action. There shall be no other jumps into or out of the compound statement associated with a delayed-catch action.

EXAMPLE This example illustrates behavioral differences between try and catch actions and delayed-try and delayed-catch actions.

Code sequence with try and catch actions:

```c
#include <fenv.h>
double d[n];
float f[n];

#pragma STDC FENV_EXCEPT TRY FE_DIVBYZERO, FE_OVERFLOW
for (i=0; i<n; i++) {
    f[i] = 1.0 / d[i];
}

#pragma STDC FENV_EXCEPT CATCH FE_DIVBYZERO
printf("divide-by-zero\n");

#pragma STDC FENV_EXCEPT CATCH FE_OVERFLOW
printf("overflow\n");
```

The same code but with delayed-try and delayed-catch actions:
#pragma STDC FENV_ACCESS ON
#include <fenv.h>
double d[n];
float f[n];
...
{
    #pragma STDC FENV_EXCEPT DELAYED_TRY \ 
    FE_DIVBYZERO, FE_OVERFLOW
    for (i=0; i<n; i++){
        f[i] = 1.0 / d[i];
    }
}
{
    #pragma STDC FENV_EXCEPT DELAYED_CATCH FE_DIVBYZERO
    printf("divide-by-zero\n");
}
{
    #pragma STDC FENV_EXCEPT DELAYED_CATCH FE_OVERFLOW
    printf("overflow\n");
}
...

The following table shows examples of inputs and results for the two code sequences above.

<table>
<thead>
<tr>
<th>Input (d)</th>
<th>try - catch</th>
<th>delayed-try – delayed-catch</th>
</tr>
</thead>
<tbody>
<tr>
<td>0.5, 0.0</td>
<td>(f = 1 / d)</td>
<td>indeterminate, indeterminate</td>
</tr>
<tr>
<td></td>
<td>output</td>
<td>“divide-by-zero”</td>
</tr>
<tr>
<td></td>
<td>“divide-by-zero” flag</td>
<td>unspecified (set or restored)</td>
</tr>
<tr>
<td></td>
<td>“overflow” flag</td>
<td>unchanged</td>
</tr>
</tbody>
</table>

<table>
<thead>
<tr>
<th>Input (d)</th>
<th>0.5, (-1e100)</th>
</tr>
</thead>
<tbody>
<tr>
<td></td>
<td>(f = 1 / d)</td>
</tr>
<tr>
<td></td>
<td>output</td>
</tr>
<tr>
<td></td>
<td>“divide-by-zero” flag</td>
</tr>
<tr>
<td></td>
<td>“overflow” flag</td>
</tr>
</tbody>
</table>

<table>
<thead>
<tr>
<th>Input (d)</th>
<th>(-1e100, 0.0)</th>
</tr>
</thead>
<tbody>
<tr>
<td></td>
<td>(f = 1 / d)</td>
</tr>
<tr>
<td></td>
<td>output</td>
</tr>
<tr>
<td></td>
<td>“divide-by-zero” flag</td>
</tr>
<tr>
<td></td>
<td>“overflow” flag</td>
</tr>
</tbody>
</table>
NOTE The delayed-try and delayed-catch actions are deterministic. They can be implemented with the floating-point exception flags. The following code sequence is equivalent to the code sequence using delayed-try and delayed-catch in the example above.

```
#pragma STDC FENV_ACCESS ON
#include <fenv.h>
double d[n];
float f[n];
...
{
    fexcept_t old_except, new_except;
    fegetexceptflag(&old_except, FE_DIVBYZERO | FE_OVERFLOW);
    feclearexcept(FE_DIVBYZERO | FE_OVERFLOW);
    {
        for (i=0; i<n; i++) {
            f[i] = 1.0 / d[i];
        }
    }
    fegetexceptflag(&new_except, FE_DIVBYZERO | FE_OVERFLOW);
    fesetexceptflag(&old_except, FE_DIVBYZERO | FE_OVERFLOW);
    if (fetestexceptflag(&new_except, FE_DIVBYZERO)) {
        printf("divide-by-zero\n");
    } else if (fetestexceptflag(&new_except, FE_OVERFLOW)) {
        printf("overflow\n");
    }
}
...
```

NOTE The try and catch actions are not deterministic (see example above), which allows more implementation flexibility for better performance.

In most cases, the try and catch actions can be implemented like delayed-try and delayed-catch actions, though not for the “underflow” exception (which occurs without causing the “underflow” flag to be set, in cases of exact results near the underflow threshold). Such an implementation would not always handle the first occurring designated exception, as recommended.

An implementation of try and catch actions using floating-point exception traps might well be able to handle the first occurring designated exception (including “underflow”), as recommended, and achieve better performance.

In 7.6.1g.1, in the try bullet, attach a footnote to the wording:

The associated compound statement shall not be the statement of a selection (6.8.4) or iteration (6.8.5) statement.

where the footnote is:

*) The compound statements associated with a try action and its catch actions, together enclosed in braces, may be the statement of a selection or iteration statement. For example, the following code segment is permitted:
for (i = 0; i < n; i++) {
    
    #pragma STDC FENV_EXCEPT TRY FE_OVERFLOW
    y[i] = x[i] * x[i];

    #pragma STDC FENV_EXCEPT CATCH FE_OVERFLOW
    y[i] = DBL_MAX;
}

In 6.5.16.2, in footnote 113, change:

For example, if annex F is in effect, the floating types involved have IEC 60559 formats, and FLT_EVAL_METHOD is 0, the equivalent code would be:

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For example, if annex F is supported, the floating types involved have IEC 60559 formats, FLT_EVAL_METHOD is 0, and no FENV_EXCEPT pragma is in effect, the equivalent code would be:
Bibliography


